

Applications of Program analysis in Model-Based Design

Prahlad Sampath (Prahlad.Sampath@mathworks.com)

“©2018 by The MathWorks, Inc., MATLAB, Simulink, Stateflow, are registered trademarks of The MathWorks, Inc. Other product or brand names are trademarks or registered trademarks of their respective holders.”

Outline

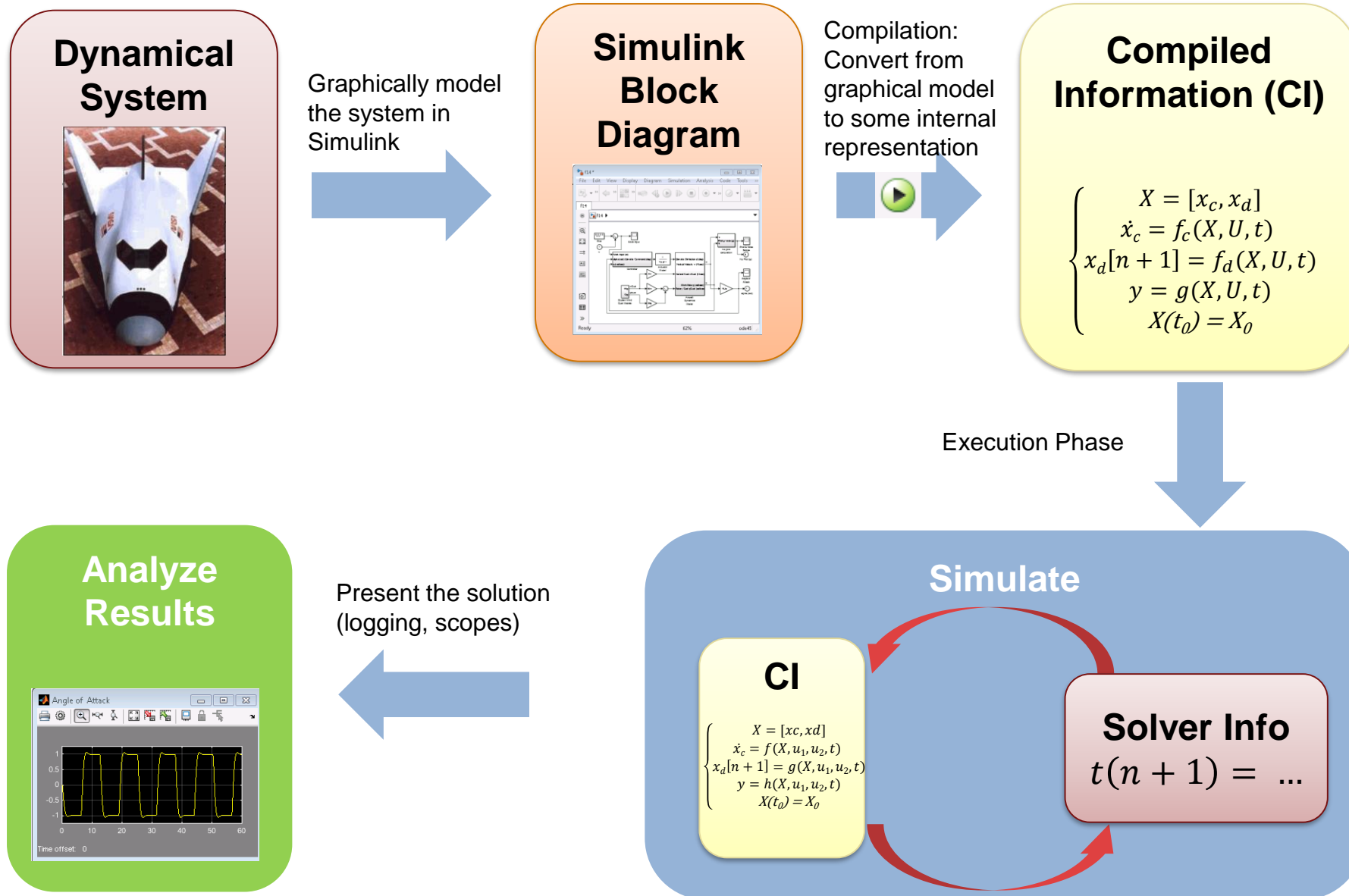
- Model-Based Design overview and role of program analysis (20 mins)
- Data flow analysis (20 mins)
- Abstract interpretation (20 mins)
- Software model checking (20 mins)

Model-Based Design Overview

Model-Based Design

- Design activities
 - Specification, prototyping (at different levels of fidelity), design space exploration, code-generation, verification
- Use of appropriate abstractions for modelling – from algorithms to systems
 - Graphical languages where appropriate
 - Textual languages where appropriate
 - Provide a framework where these can be combined and analysed
- Analysis is the corner stone of Model-Based Design
 - Different levels of adoption of Model-Based Design will use different kinds of analysis

Simulink Workflow



Modelling Formalisms

- Linear systems (ODEs)

$$\frac{dv}{dt} = -g,$$

$$\frac{dx}{dt} = v,$$

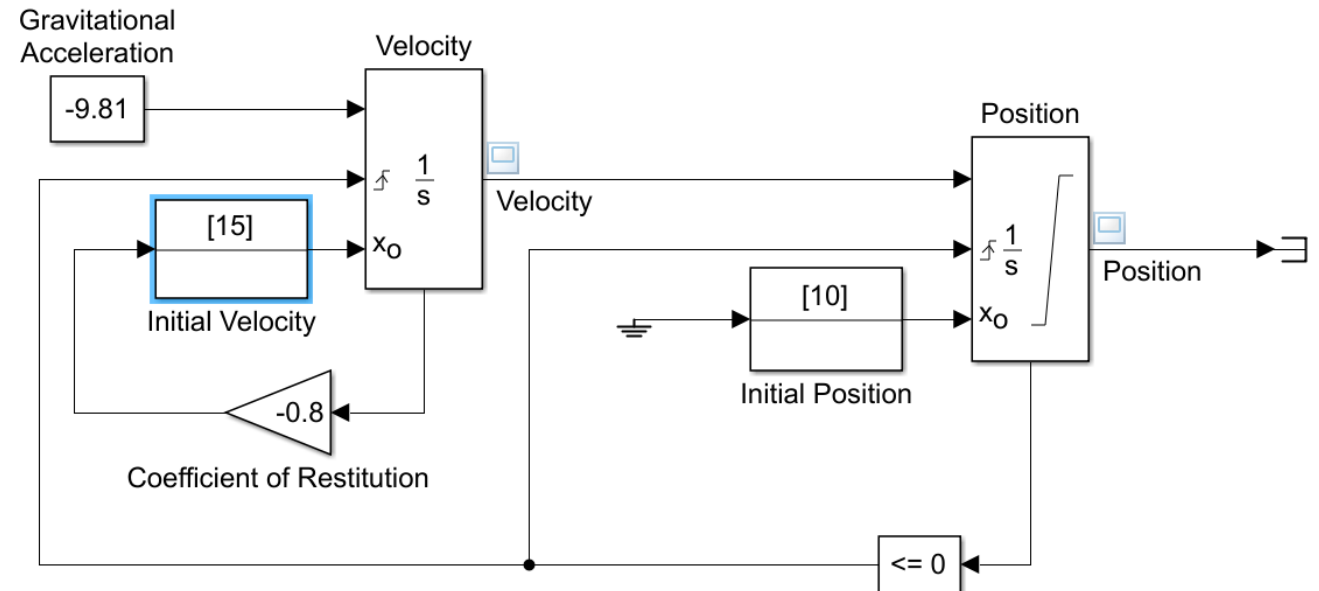
$$v_0 = 15 \text{ m/s} \quad x_0 = 10 \text{ m}$$

$$v^+ = -\kappa v^-, \quad x = 0$$

```
>>sldemo_bounce_two_integrators
```

Modelling Formalisms

- Linear systems (ODEs)

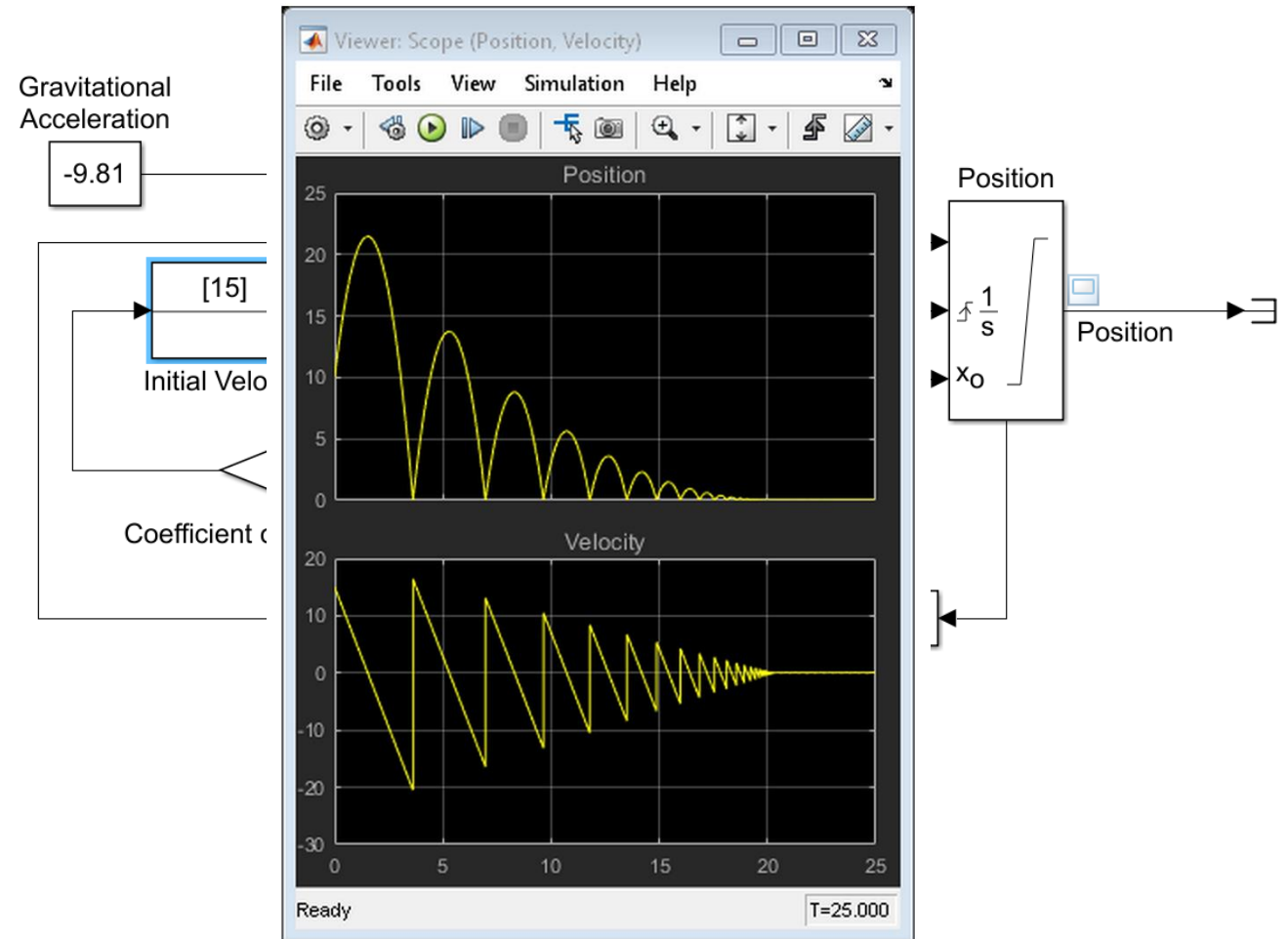


Copyright 1990-2013 The MathWorks, Inc.

```
>>sldemo_bounce_two_integrators
```

Modelling Formalisms

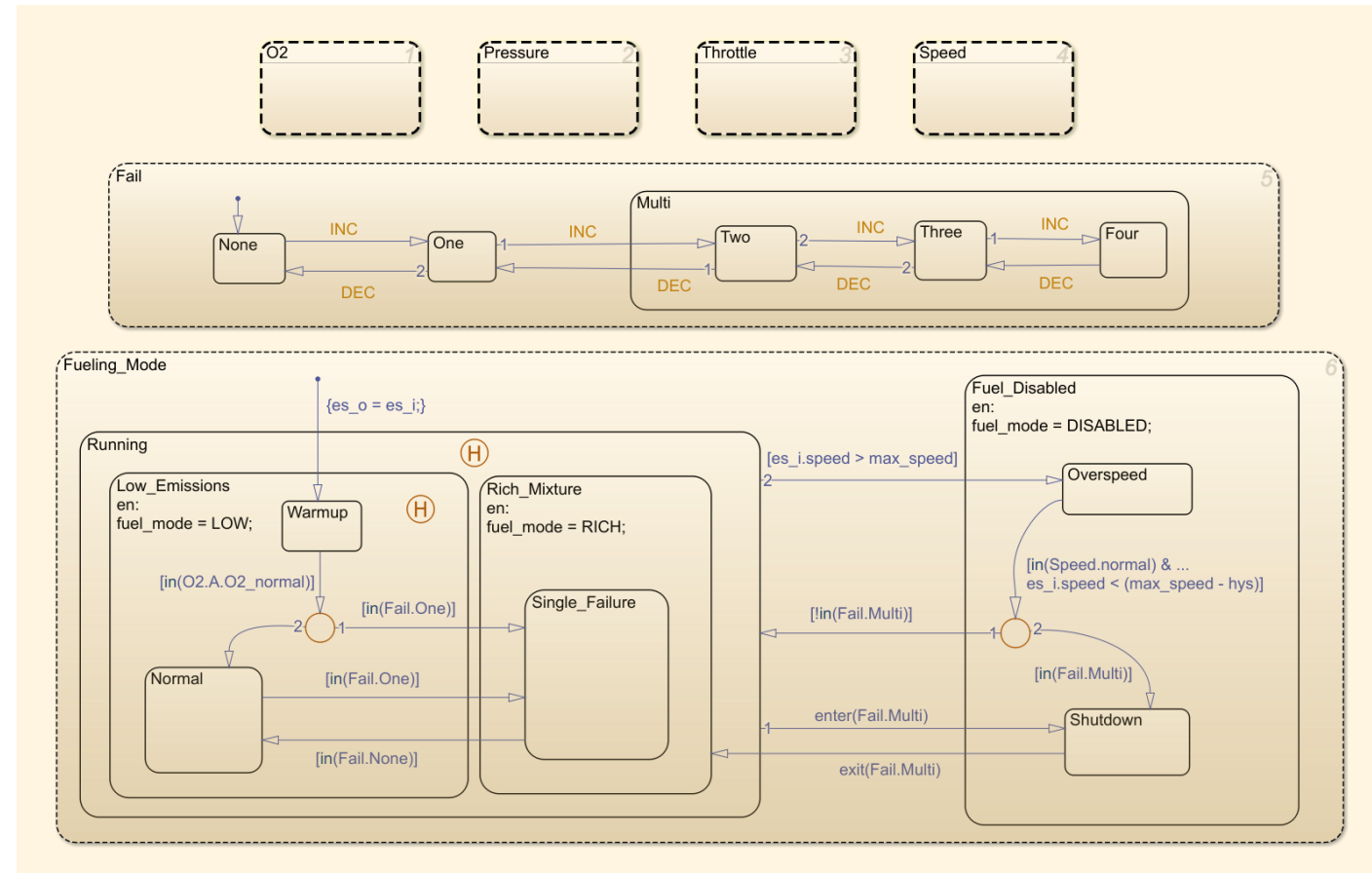
- Linear systems (ODEs)



```
>>sldemo_bounce_two_integrators
```


Modelling Formalisms

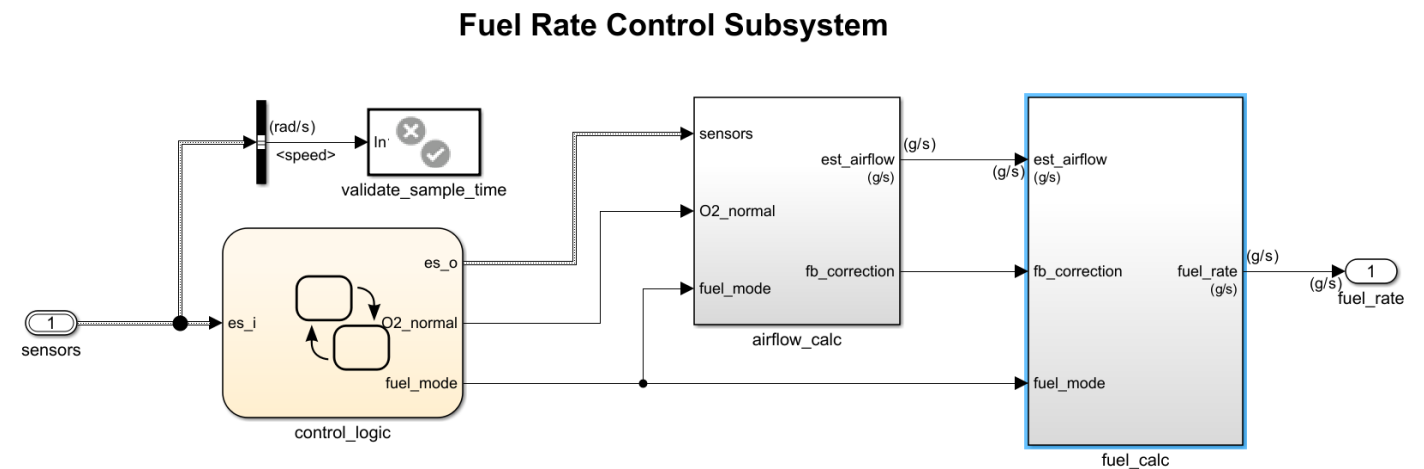
- Linear systems (ODEs)
- Hierarchical state machines



>>sldemo_fuelsys

Modelling Formalisms

- Linear systems (ODEs)
- Hierarchical state machines



>>sldemo_fuelsys

Modelling Formalisms

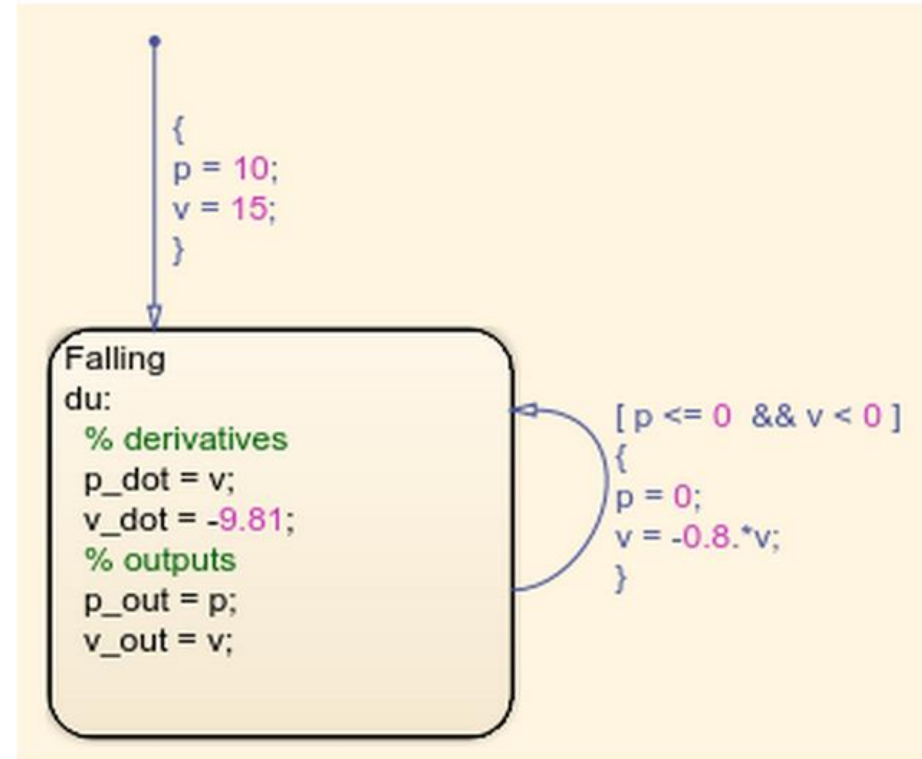
- Linear systems (ODEs)
- Hierarchical state machines
- Tabular specifications

STATES	TRANSITIONS	
	IF	ELSE-IF(2)
Normal	[ALARM]	
	Alarm	
Off entry: boiler_cmd = 0; doneWarmup = false;	[temp <= reference_low]	
	Warmup	
Warmup entry: boiler_cmd = 2;	[doneWarmup]	[after(10, sec)] {doneWarmup = true;}
	On	On
On entry: boiler_cmd = 1;	[temp >= reference_high]	
	Off	
Alarm entry: boiler_cmd = 0;	[CLEAR]	
	Normal	

>>ex_stt_boiler

Modelling Formalisms

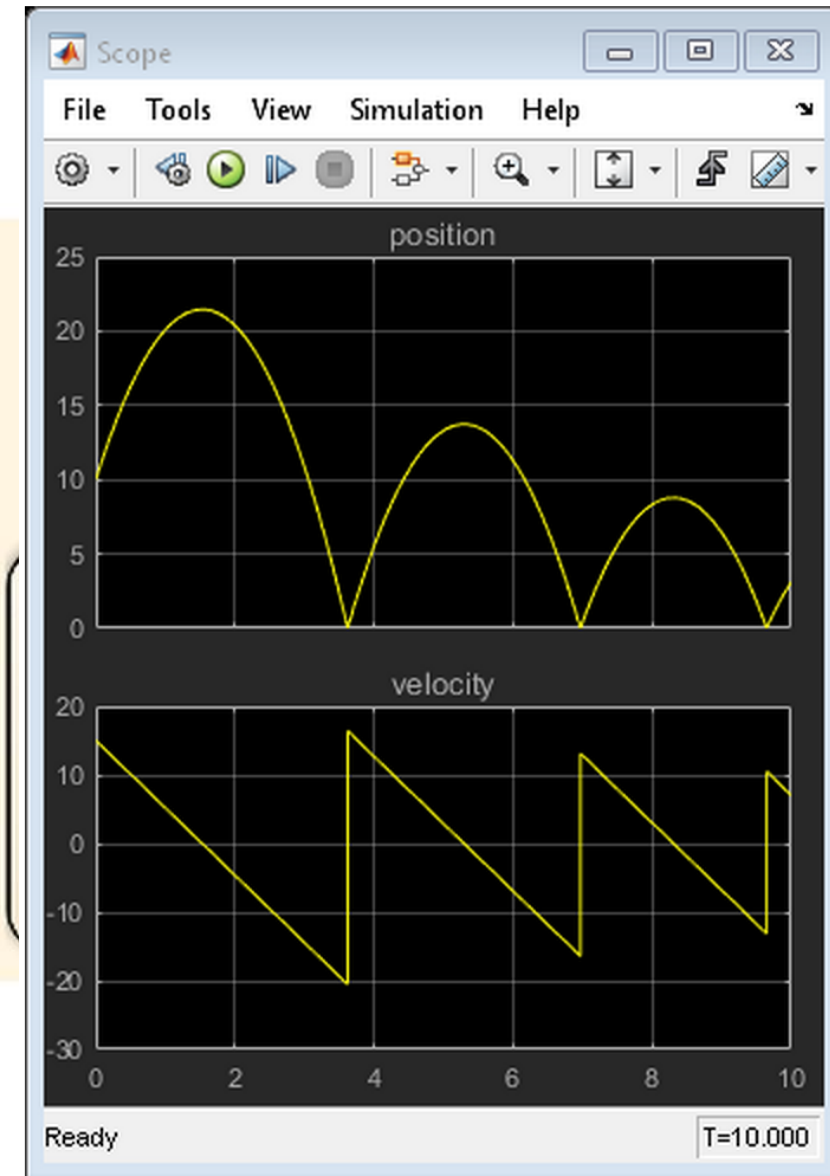
- Linear systems (ODEs)
- Hierarchical state machines
- Tabular specifications
- Hybrid automata



>>sf_bounce

Modelling Formalisms

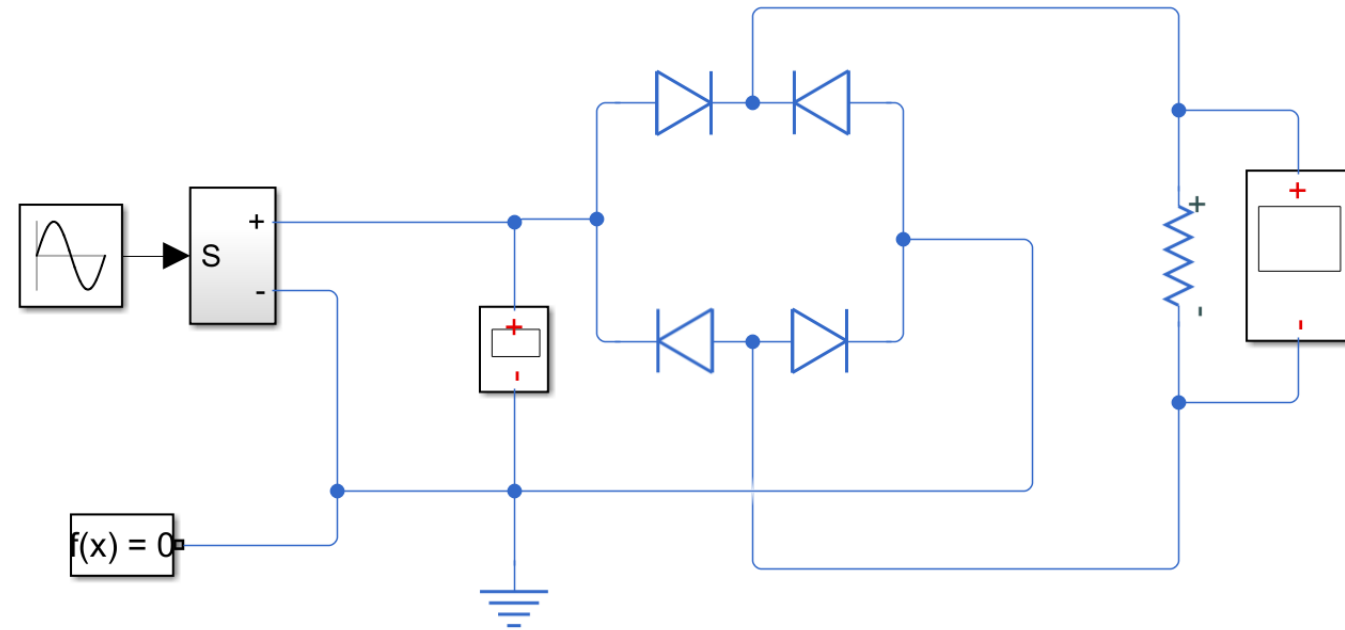
- Linear systems (ODEs)
- Hierarchical state machines
- Tabular specifications
- Hybrid automata



```
>>sf_bounce
```

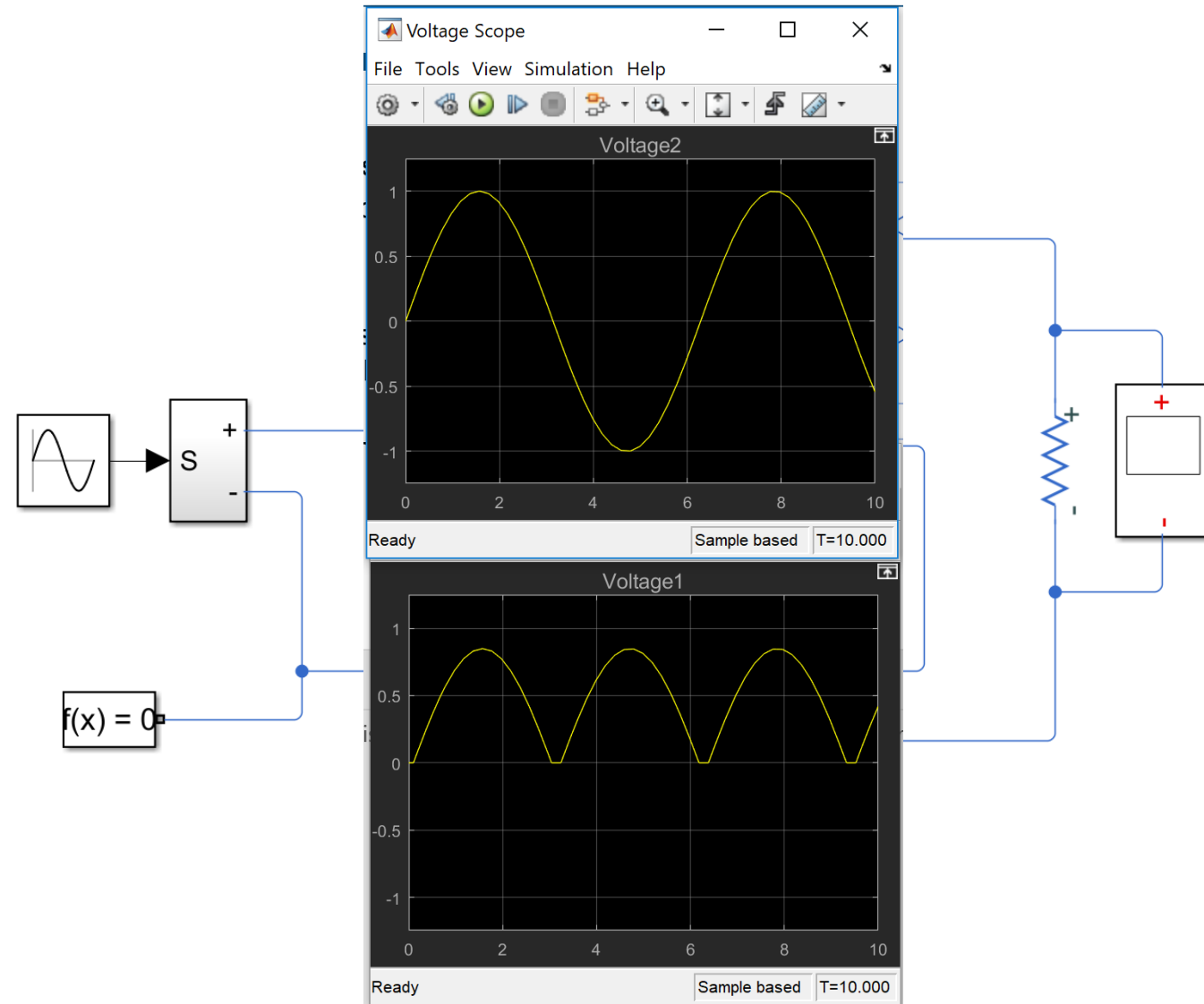
Modelling Formalisms

- Linear systems (ODEs)
- Hierarchical state machines
- Tabular specifications
- Hybrid automata
- Differential algebraic system (acausal models for physical systems)



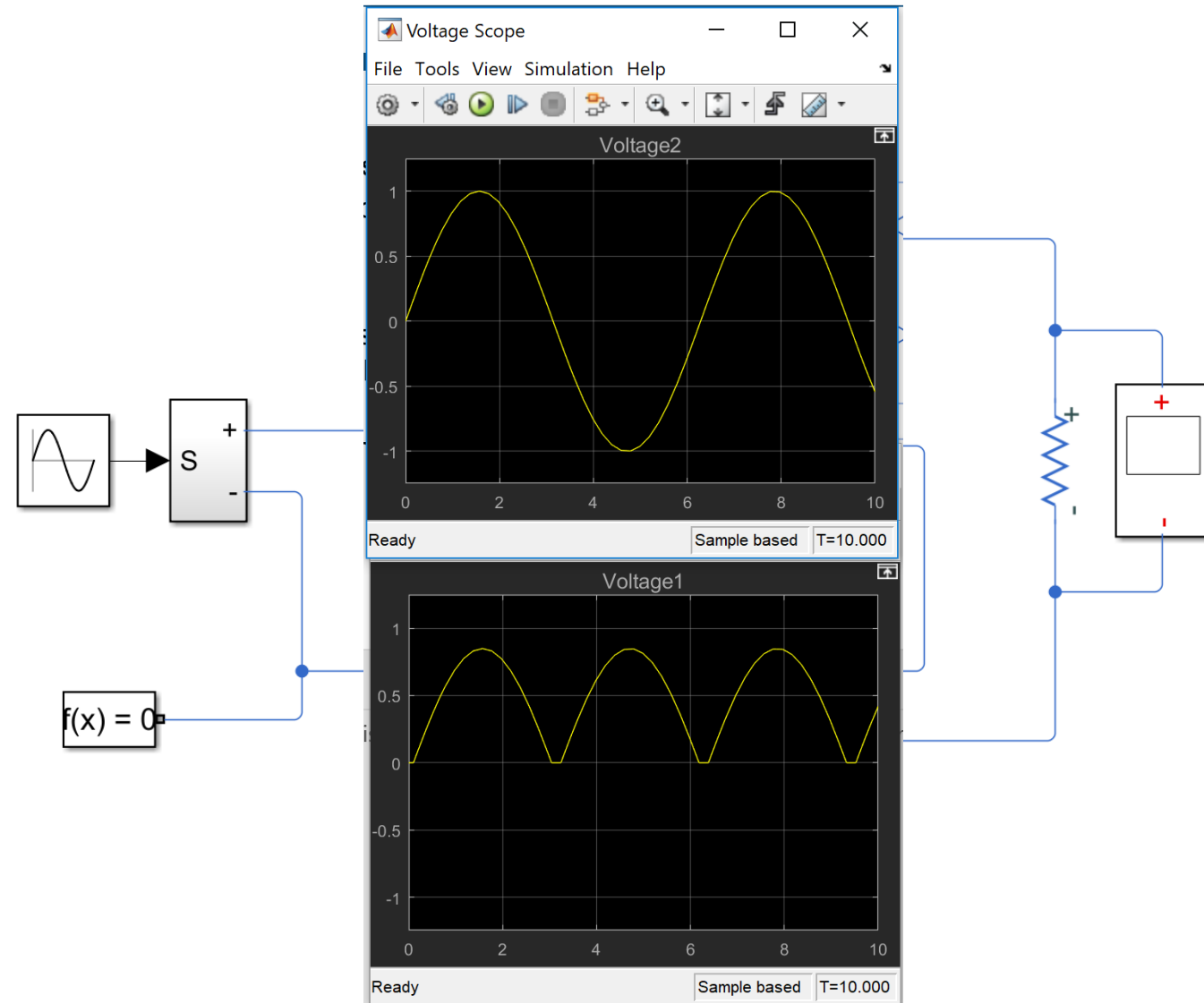
Modelling Formalisms

- Linear systems (ODEs)
- Hierarchical state machines
- Tabular specifications
- Hybrid automata
- Differential algebraic system (acausal models for physical systems)



Modelling Formalisms

- Linear systems (ODEs)
- Hierarchical state machines
- Tabular specifications
- Hybrid automata
- Differential algebraic system (acausal models for physical systems)
- Discrete event systems
- Procedural code



Model-Based Design – Hardware

- Hardware connectivity
- Real time simulation

Model-Based Design – Hardware

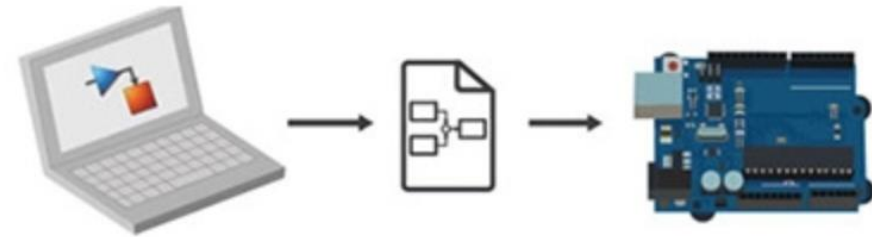
- Hardware connectivity
- Real time simulation



With MATLAB support package for Arduino, the Arduino is connected to a computer running MATLAB. Processing is done on the computer with MATLAB.

Model-Based Design – Hardware

- Hardware connectivity
- Real time simulation

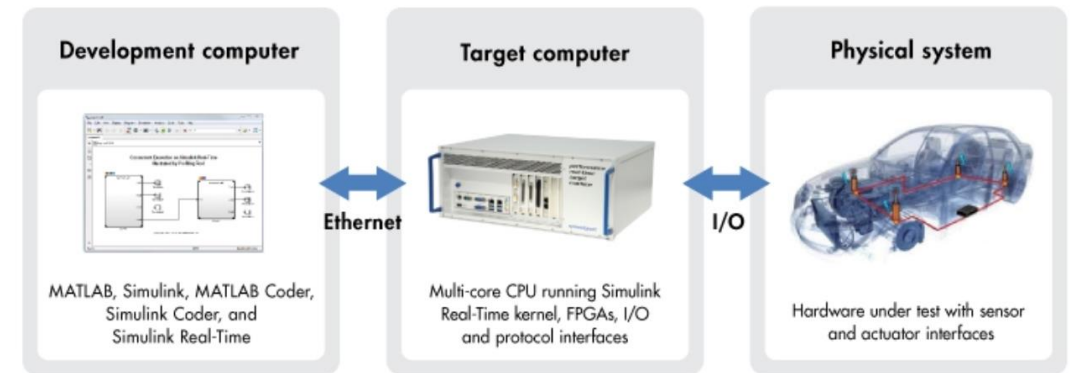


With Simulink support package for Arduino, you develop the algorithm in Simulink and deploy to the Arduino using automatic code generation. Processing is then done on the Arduino.

Model-Based Design – Hardware

- Hardware connectivity
- Real time simulation

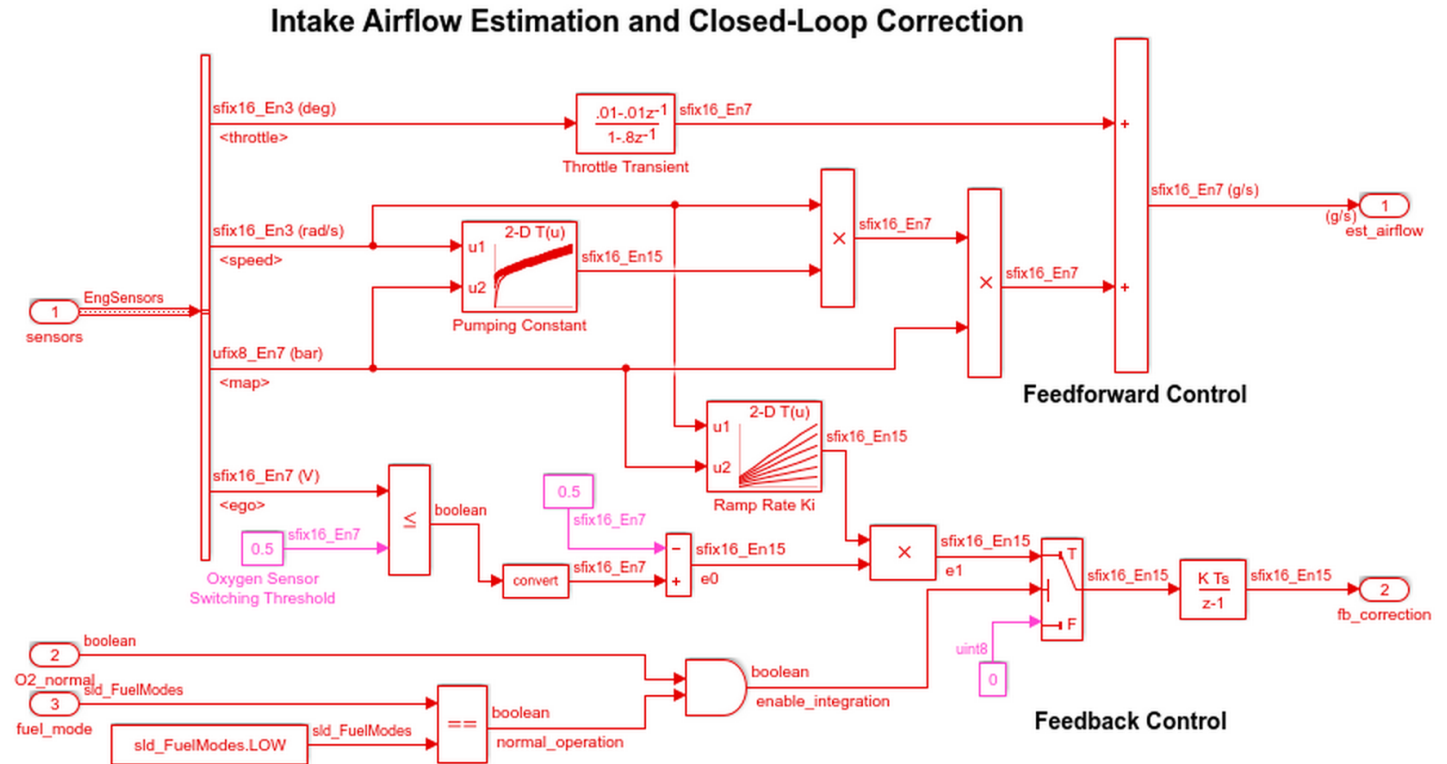
Typical real-time simulation and testing environment using Simulink Real-Time.



Typical real-time simulation and testing environment using Simulink Real-Time.

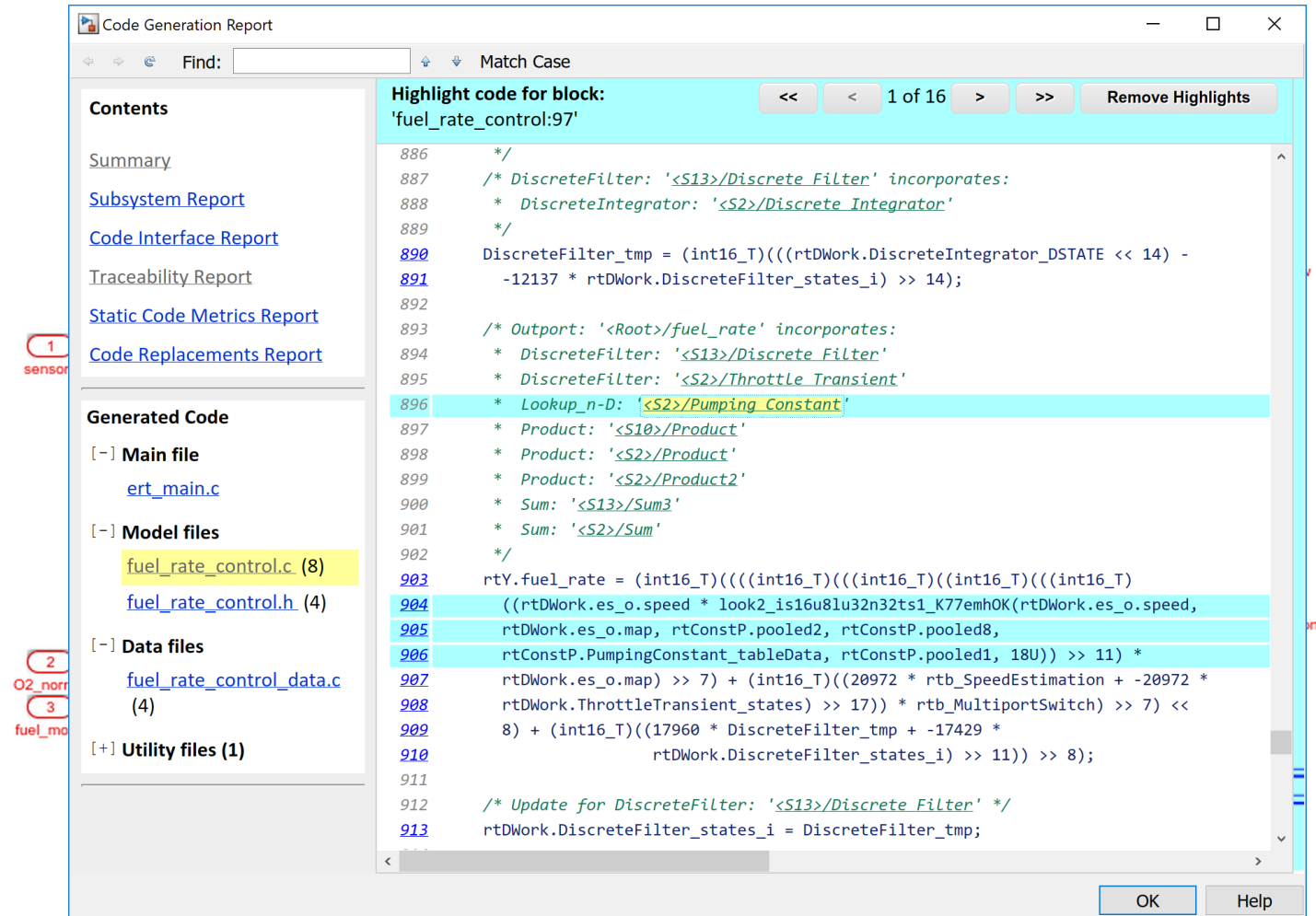
Model-Based Design – Code generation

- Programming
 - Code generation
 - Design optimizations
 - Verification



Model-Based Design – Code generation

- Programming
 - Code generation
 - Design optimizations
 - Verification



Code Generation Report

Find: Match Case

Highlight code for block: 'fuel_rate_control:97'

1 of 16

Remove Highlights

Contents

- Summary
- Subsystem Report
- Code Interface Report
- Traceability Report
- Static Code Metrics Report
- Code Replacements Report

Generated Code

- [-] Main file
 - ert_main.c
- [-] Model files
 - fuel_rate_control.c (8)
 - fuel_rate_control.h (4)
- [-] Data files
 - fuel_rate_control_data.c (4)
- [+] Utility files (1)

Code Snippet:

```

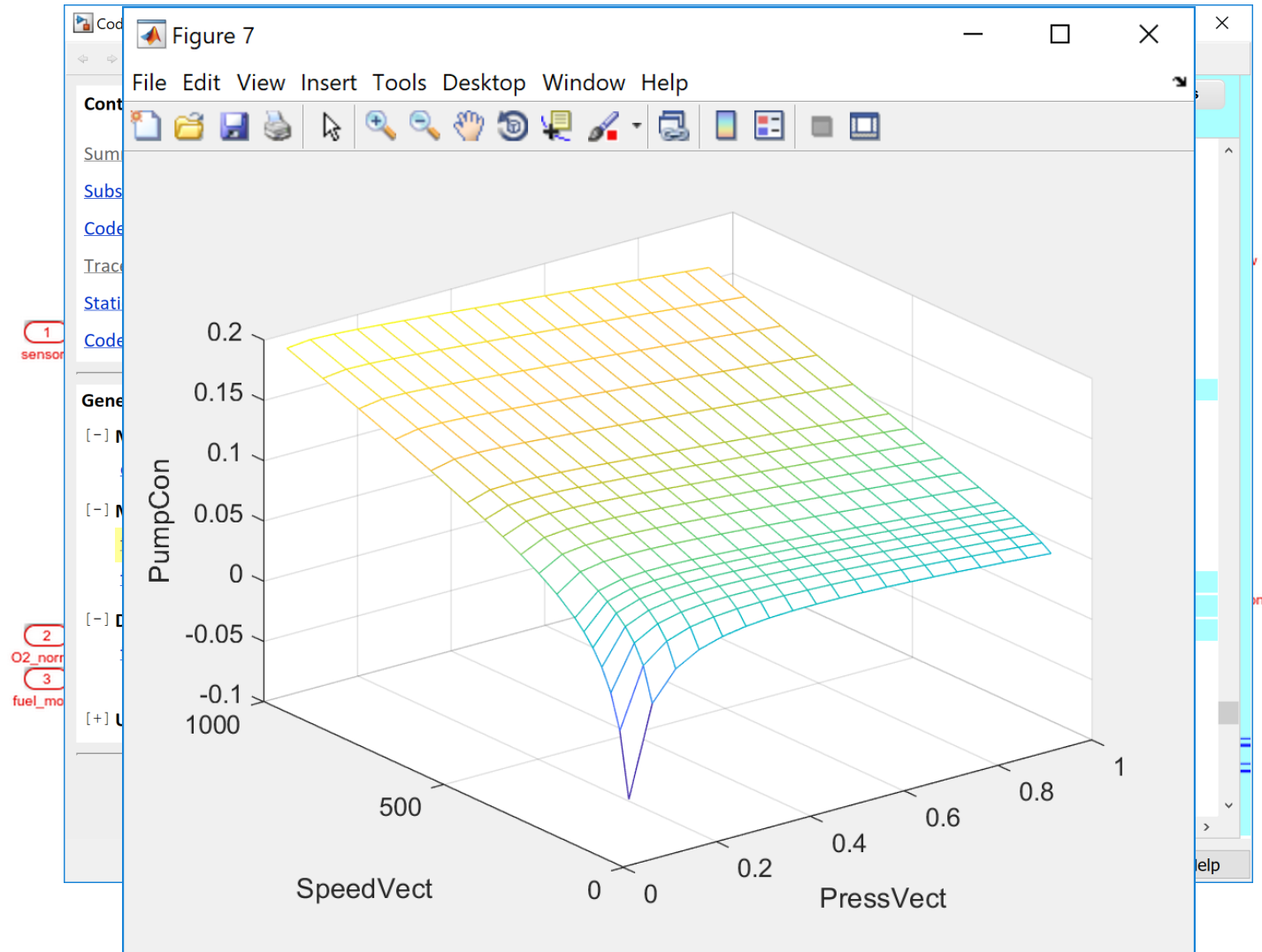
886  */
887  /* DiscreteFilter: '<S13>/Discrete_Filter' incorporates:
888  * DiscreteIntegrator: '<S2>/Discrete_Integrator'
889  */
890  DiscreteFilter_tmp = (int16_T)((((rtDWork.DiscreteIntegrator_DSTATE << 14) -
891  -12137 * rtDWork.DiscreteFilter_states_i) >> 14));
892
893  /* Output: '<Root>/fuel_rate' incorporates:
894  * DiscreteFilter: '<S13>/Discrete_Filter'
895  * DiscreteFilter: '<S2>/Throttle_Transient'
896  * Lookup_n-D: '<S2>/Pumping_Constant'
897  * Product: '<S10>/Product'
898  * Product: '<S2>/Product'
899  * Product: '<S2>/Product2'
900  * Sum: '<S13>/Sum3'
901  * Sum: '<S2>/Sum'
902  */
903  rtY.fuel_rate = (int16_T)((((int16_T)((((int16_T)((((int16_T)((((int16_T)
904  ((rtDWork.es_o.speed * look2_is16u8lu32n32ts1_K77emhOK(rtDWork.es_o.speed,
905  rtDWork.es_o.map, rtConstP.pooled2, rtConstP.pooled8,
906  rtConstP.PumpingConstant_tableData, rtConstP.pooled1, 18U)) >> 11) *
907  rtDWork.es_o.map) >> 7) + (int16_T)((20972 * rtb_SpeedEstimation + -20972 *
908  rtDWork.ThrottleTransient_states) >> 17)) * rtb_MultiportSwitch) >> 7) <<
909  8) + (int16_T)((17960 * DiscreteFilter_tmp + -17429 *
910  rtDWork.DiscreteFilter_states_i) >> 11)) >> 8));
911
912  /* Update for DiscreteFilter: '<S13>/Discrete_Filter' */
913  rtDWork.DiscreteFilter_states_i = DiscreteFilter_tmp;

```

OK Help

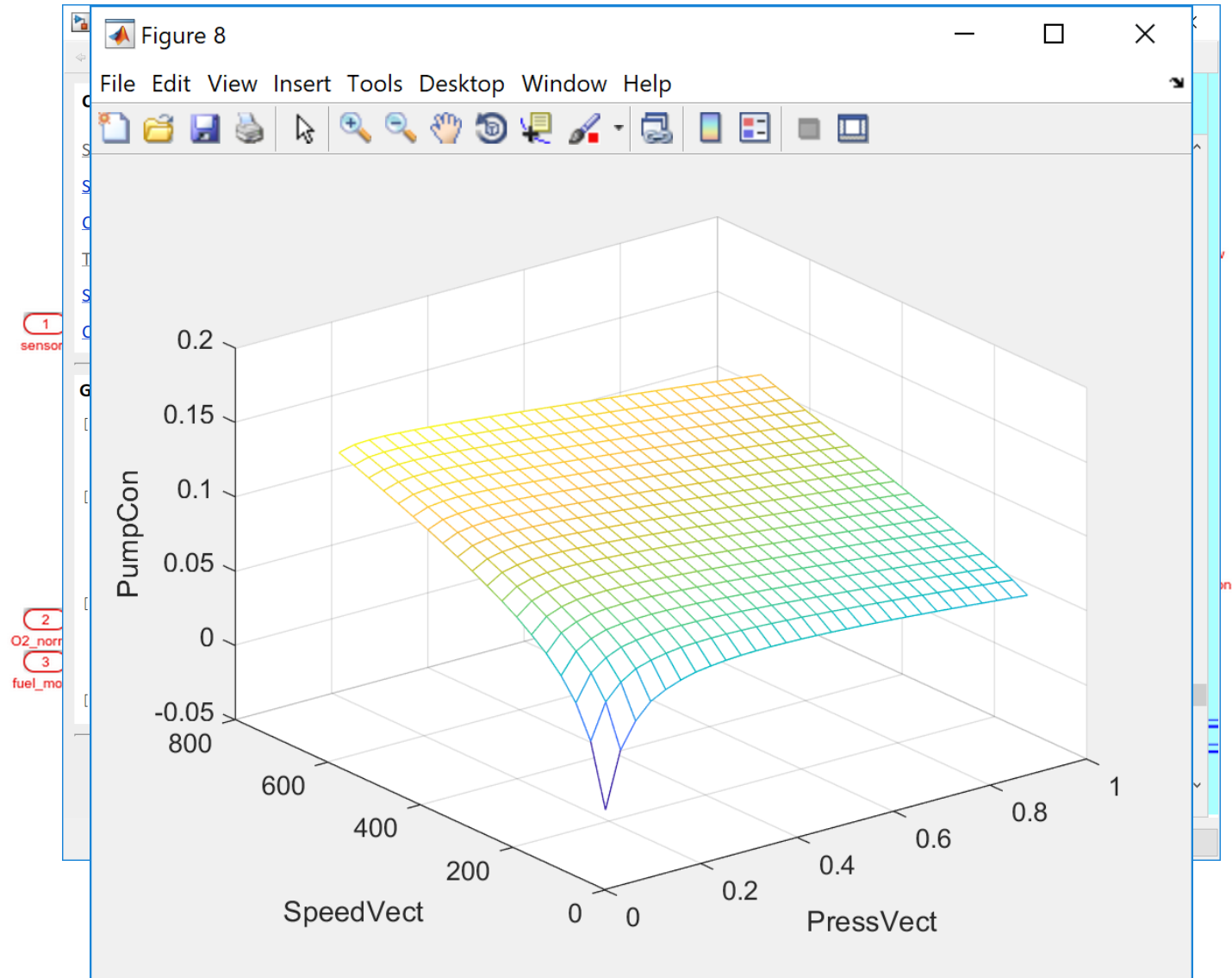
Model-Based Design – Code generation

- Programming
 - Code generation
 - Design optimizations
 - Verification



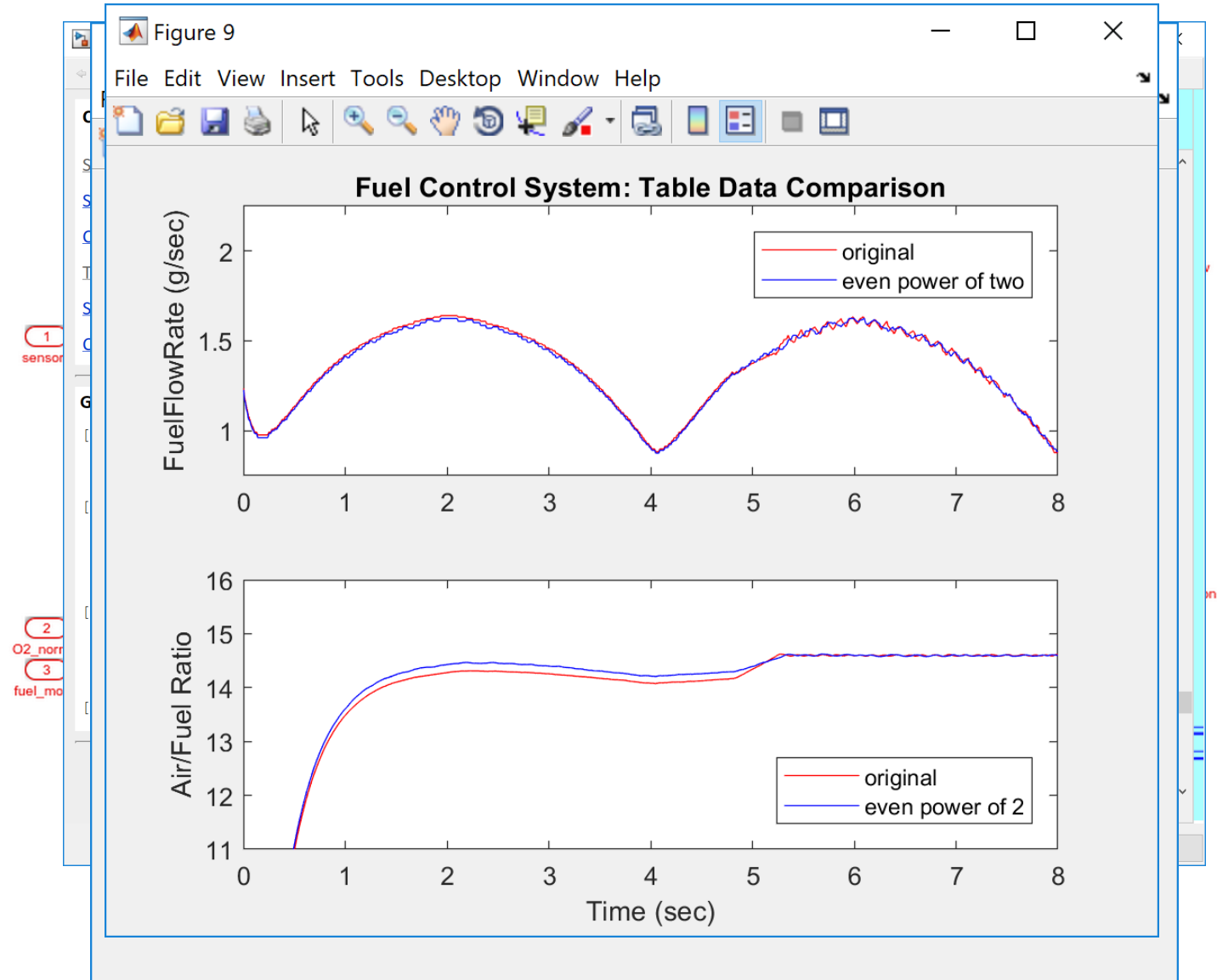
Model-Based Design – Code generation

- Programming
 - Code generation
 - Design optimizations
 - Verification



Model-Based Design – Code generation

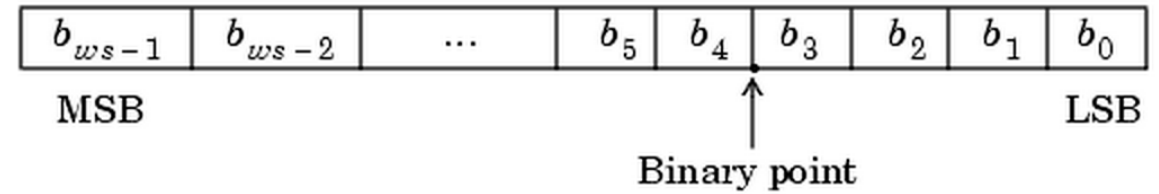
- Programming
 - Code generation
 - Design optimizations
 - Verification



Model-Based Design – Code generation

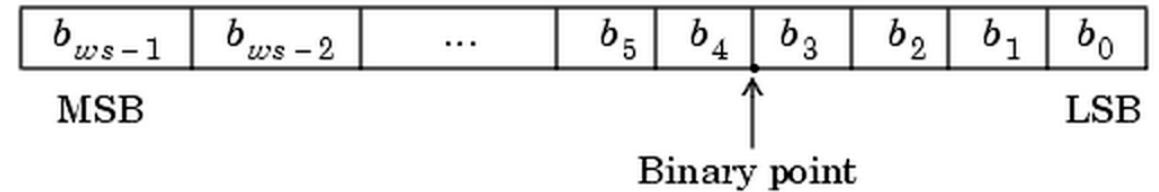
- Programming
 - Code generation
 - Design optimizations
 - Verification

- Data-type design



Model-Based Design – Code generation

- Programming
 - Code generation
 - Design optimizations
 - Verification
- Data-type design
- Systematic verification
 - Coverage analysis
 - Detect run-time error
 - Property proving



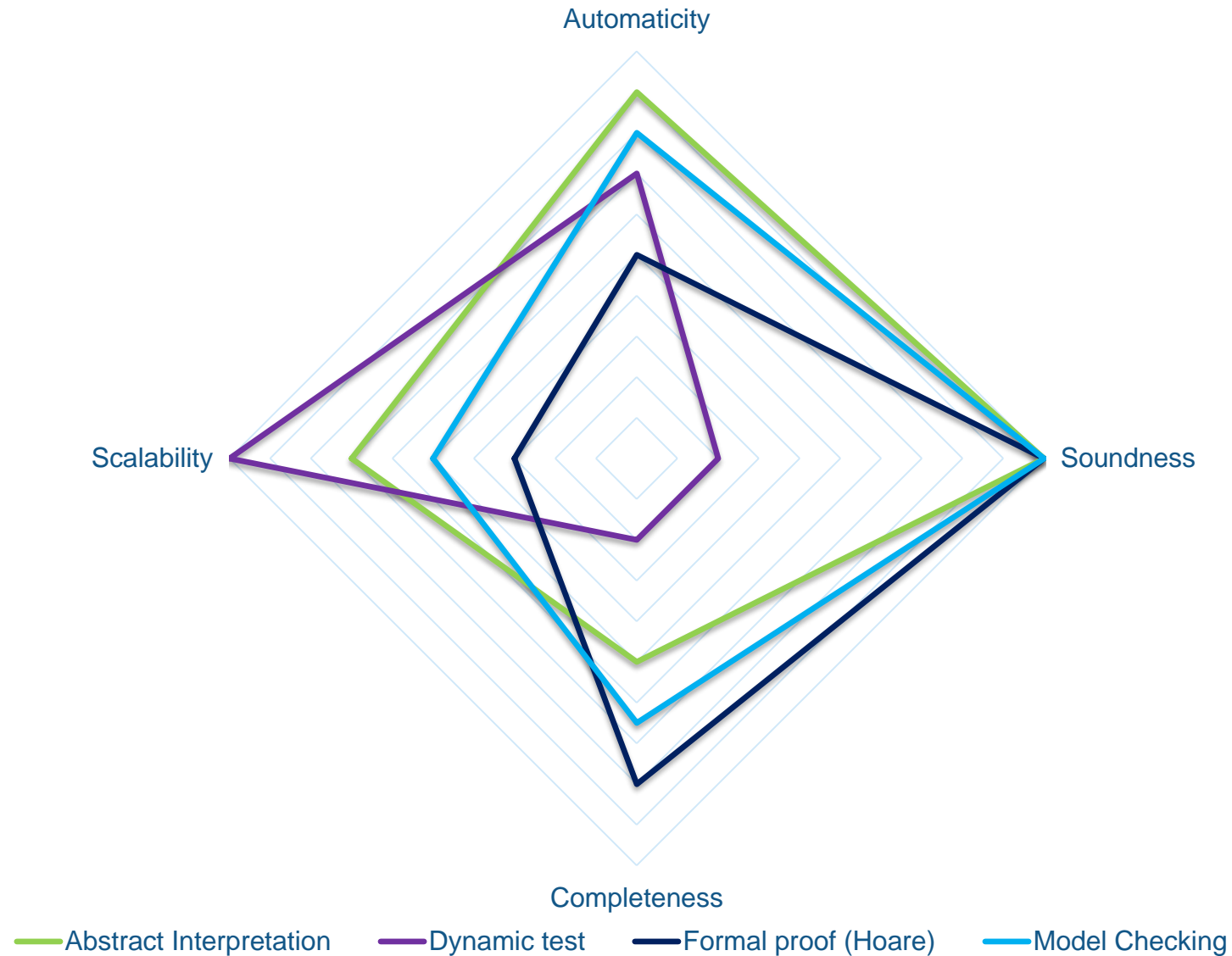
Where does Program Analysis fit into this picture?

- Techniques for reasoning about behaviour
 - Spanning many programming paradigms
- Well developed abstraction techniques
 - Abstraction vs. approximation!
- Scalable algorithms
 - Algorithms that scale to millions of lines of code
- Active area of research
 - New techniques and algorithms

Four properties for a tool

- Automaticity
 - Is the tool press-button or does it require user interaction ?
- Soundness
 - Are results trustable ?
- Completeness
 - Will the tool find all my bugs ?
- Scalability
 - How does the tool react to large programs ?

Comparison of formal methods properties

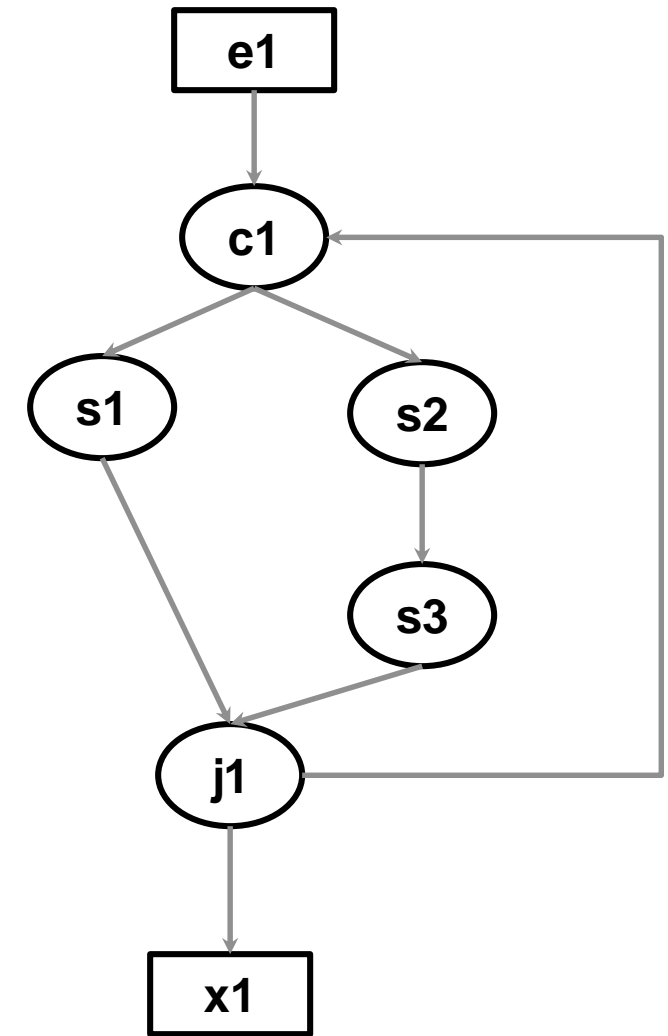


It's all about the Workflow!

Data Flow Analysis

A Refresher

- Infer “flows” in programs
 - Reason about program statements
 - System of equations that capture the “flow” across statements in a program
 - Algorithms well suited for program graphs
 - Basis for compiler optimizations



A Refresher

- Infer “flows” in programs
 - Reason about program statements
 - System of equations that capture the “flow” across statements in a program
 - Algorithms well suited for program graphs
 - Basis for compiler optimizations
- Generate a set of equations
 - Over a lattice domain
 - Solution calculated by using fixpoint iteration

$$entry(s1) = exit(e1)$$

...

$$entry(c1) = entry(e1) \cup exit(j1)$$

$$exit(c1) = entry(c1) - kill(c1) \cup gen(c1)$$

...

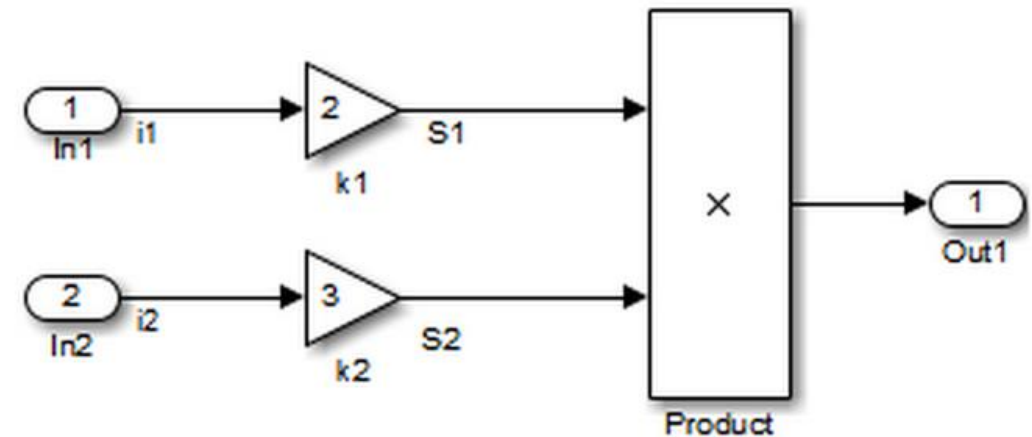
$$entry(j1) = exit(s3) \cup exit(s3)$$

A Refresher

- Infer “flows” in programs
 - Reason about program statements
 - System of equations that capture the “flow” across statements in a program
 - Algorithms well suited for program graphs
 - Basis for compiler optimizations
- Generate a set of equations
 - Over a lattice domain
 - Solution calculated by using fixpoint iteration
- Forward vs. Backward analysis
- Context sensitivity
- Flow sensitivity
- Inter vs. Intra procedural analysis
- ...

Applications of data-flow analysis

- Compiler optimizations
 - Expression folding
 - Variable reuse (RAM)
 - Constant folding
 - Strength reduction



Applications of data-flow analysis

- Compiler optimizations
 - Expression folding
 - Variable reuse (RAM)
 - Constant folding
 - Strength reduction

```
/* Model step function */
void exprfld_step(void)
{
    /* Gain: '<Root>/Gain' incorporates:
     *   Inport: '<Root>/In1'
     */
    exprfld_B.S1 = exprfld_P.Gain_Gain * exprfld_U.i1;

    /* Gain: '<Root>/Gain1' incorporates:
     *   Inport: '<Root>/In2'
     */
    exprfld_B.S2 = exprfld_P.Gain1_Gain * exprfld_U.i2;

    /* Outport: '<Root>/Out1' incorporates:
     *   Product: '<Root>/Product'
     */
    exprfld_Y.Out1 = exprfld_B.S1 * exprfld_B.S2;
}
```

Applications of data-flow analysis

- Compiler optimizations
 - Expression folding
 - Variable reuse (RAM)
 - Constant folding
 - Strength reduction

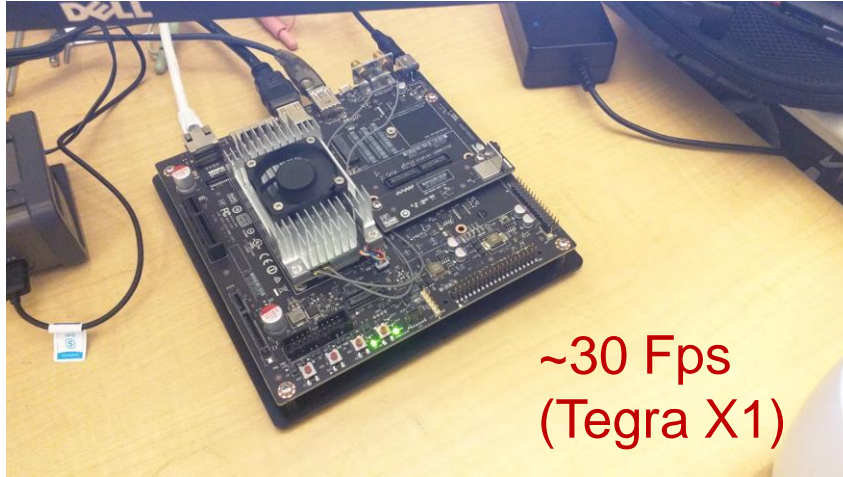
```
/* Model step function */  
void exprfld_step(void)  
{  
    /* Outport: '<Root>/Out1' incorporates:  
     * Gain: '<Root>/Gain'  
     * Gain: '<Root>/Gain1'  
     * Inport: '<Root>/In1'  
     * Inport: '<Root>/In2'  
     * Product: '<Root>/Product'  
     */  
    exprfld_Y.Out1 =  
        exprfld_P.Gain_Gain *  
        exprfld_U.i1 *  
        (exprfld_P.Gain1_Gain * exprfld_U.i2);  
}
```

Applications of data-flow analysis

- Compiler optimizations
 - Expression folding
 - Variable reuse (RAM)
 - Constant folding
 - Strength reduction
- Parallelization
 - GPU Coder™

Deploy with GPU Coder™

Alexnet



Vehicle
Detection



People detection



Lane detection



Debugging Workflow

- Problem:

Debugging Workflow

- Problem:
 - You have a large model, with an error. How should you go about localizing the error?

Debugging Workflow

- Problem:
 - You have a large model, with an error. How should you go about localizing the error?
 - Localize the behaviour over time
 - Simulate the model
 - Insert break-points
 - Visualize signals/outputs

Debugging Workflow

- Problem:
 - You have a large model, with an error. How should you go about localizing the error?
 - Localize the behaviour over time
 - Simulate the model
 - Insert break-points
 - Visualize signals/outputs
 - Localize the behaviour over space
 - Dependency analysis of model

Debugging Workflow

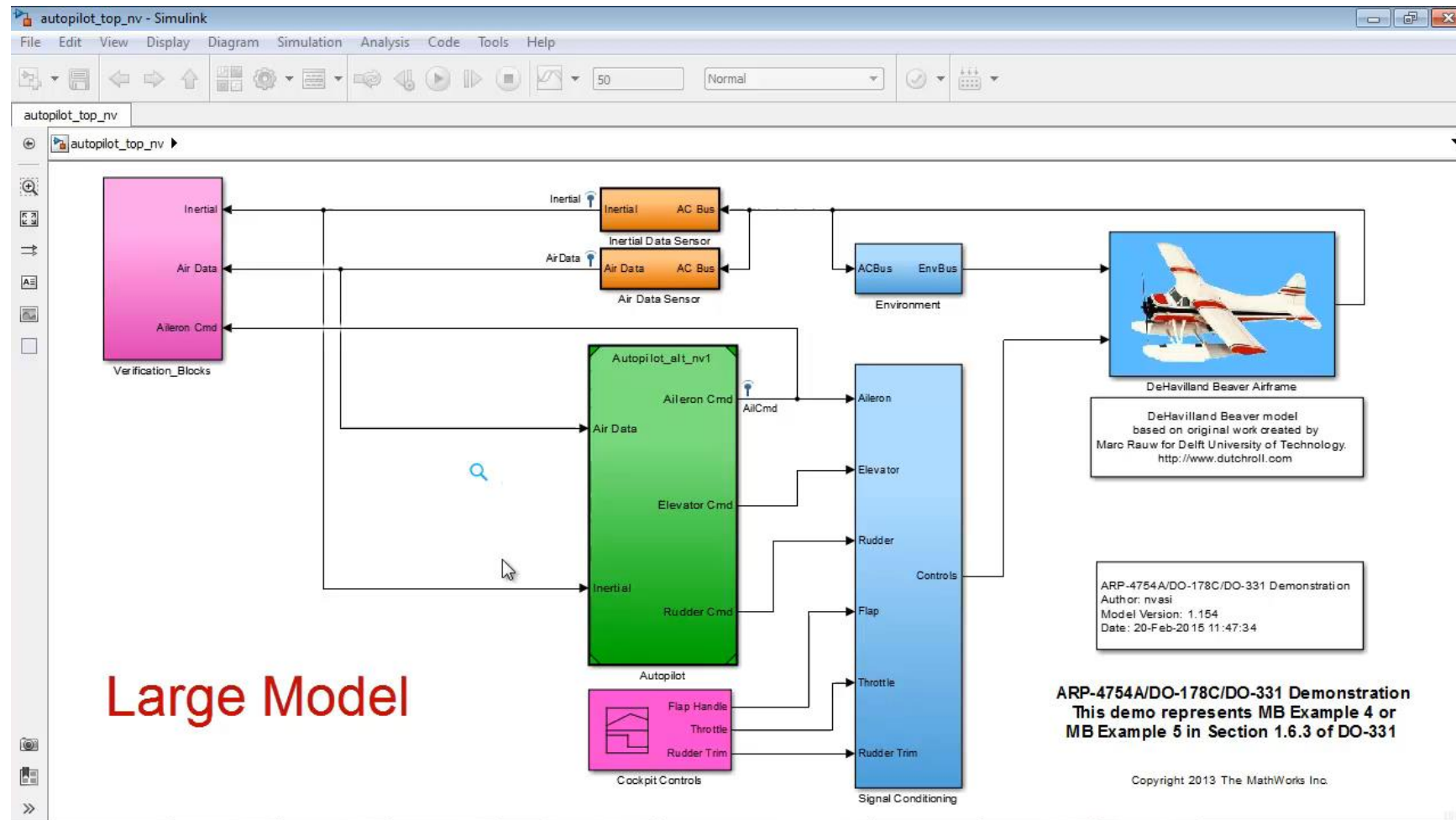
- Problem:
 - You have a large model, with an error. How should you go about localizing the error?
 - Localize the behaviour over time
 - Simulate the model
 - Insert break-points
 - Visualize signals/outputs
 - Localize the behaviour over space
 - Dependency analysis of model
 - Extract and simplify

Debugging Workflow

- Problem:
 - You have a large model, with an error. How should you go about localizing the error?
 - Localize the behaviour over time
 - Simulate the model
 - Insert break-points
 - Visualize signals/outputs
 - Localize the behaviour over space
 - Dependency analysis of model
 - Extract and simplify

Google search: “simulink model slicer youtube”
<https://www.youtube.com/watch?v=48EAr508mmw>

Slicing Models to Reduce Complexity



Abstract Interpretation

Prove that no division by zero occurs here

```
1 // x is an unknown parameter
2 float runTheLoop(int x)
3 {
4     if (x<0 || x>100)
5         return 1.0;
6     float in = input();
7     float current_out = 0.0;
8     float prev_out = 0.0;
9     int i;
10    for (i=0;i<=x;i++)
11    {
12        current_out = in*0.1 + prev_out*0.9;
13        prev_out = current_out;
14        if (random())
15            prev_out = 0.0;
16        in = input();
17    }
18
19    current_out = current_out + 3.0;
20
21    return 1/current_out;
22 }
```

Prove that no division by zero occurs here

```
1 // x is an unknown parameter
2 float runTheLoop(int x)
3 {
4     if (x<0 || x>100)
5         return 1.0;
6     float in = input();
7     float current_out = 0.0;
8     float prev_out = 0.0;
9     int i;
10    for (i=0;i<=x;i++)
11    {
12        current_out = in*0.1 + prev_out*0.9;
13        prev_out = current_out;
14        if (random())
15            prev_out = 0.0;
16        in = input();
17    }
18
19    current_out = current_out + 3.0;
20
21    return 1/current_out;
22 }
```

Clearly testing is not practicable

- Too many unknown

Prove that no division by zero occurs here

```
1 // x is an unknown parameter
2 float runTheLoop(int x)
3 {
4     if (x<0 || x>100)
5         return 1.0;
6     float in = input();
7     float current_out = 0.0;
8     float prev_out = 0.0;
9     int i;
10    for (i=0;i<=x;i++)
11    {
12        current_out = in*0.1 + prev_out*0.9;
13        prev_out = current_out;
14        if (random())
15            prev_out = 0.0;
16        in = input();
17    }
18
19    current_out = current_out + 3.0;
20
21    return 1/current_out;
22 }
```

Clearly testing is not practicable

- Too many unknown

Philosophy of AI to prove correctness:

- You don't need the exact value of `current_out` to prove absence of ZDIV
- Having a *less precise* information
e.g. `current_out >= 0.2`
is sufficient

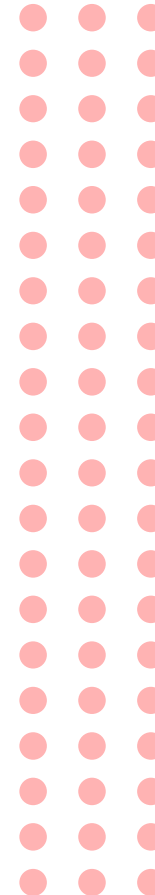
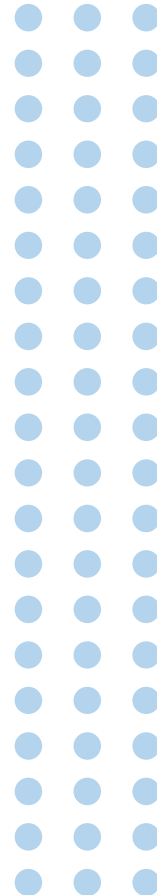
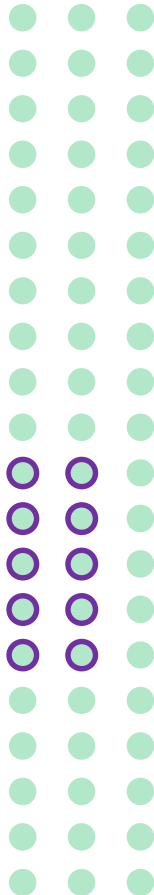
Formalization of the **concrete** semantics of the program

States of the program

```

1  // x is an unknown parameter
2  float runTheLoop(int x)
3  {
4    if (x<0 || x>100)
5      return 1.0;
6    float in = input();
7    float current_out = 0.0;
8    float prev_out = 0.0;
9    int i;
10   for (i=0;i<=x;i++)
11   {
12     current_out = in*0.1 + prev_out*0.9;
13     prev_out = current_out;
14     if (random())
15       prev_out = 0.0;
16     in = input();
17   }
18   current_out = current_out + 3.0;
19   return 1/current_out;
20 }

```



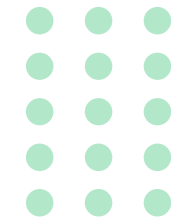
Formalization of the **concrete** semantics of the program

States of the program

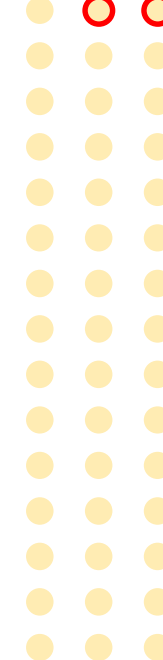
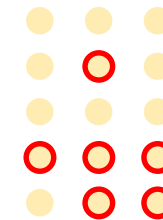
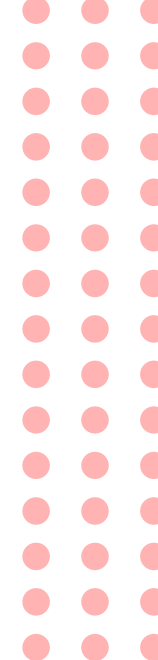
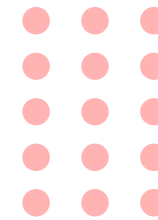
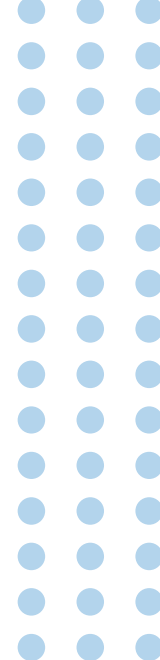
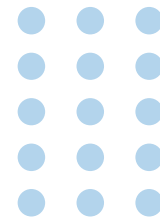
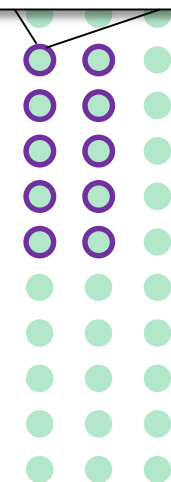
```

1 // x is an unknown parameter
2 float runTheLoop(int x)
3 {
4   if (x<0 || x>100)
5     return 1.0;
6   float in = input();
7   float current_out = 0.0;
8   float prev_out = 0.0;
9   int i;
10  for (i=0;i<=x;i++)
11  {
12    current_out = in*0.1 + prev_out*0.9;
13    prev_out = current_out;
14    if (random())
15      prev_out = 0.0;
16    in = input();
17  }
18  current_out = current_out + 3.0;
19
20
21  return 1/current_out;
22 }

```



x=1
current_out=?
in=?
i=?



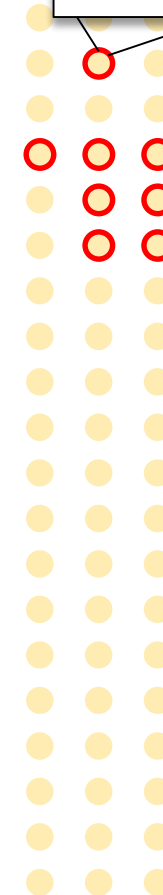
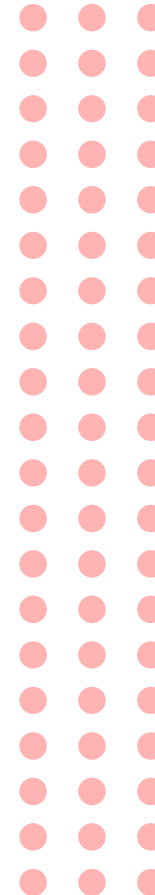
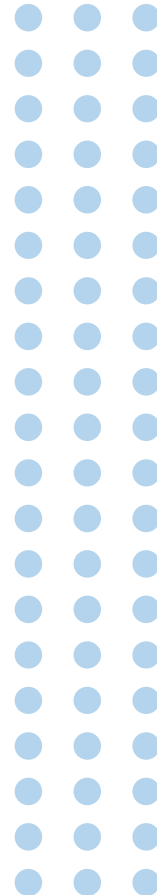
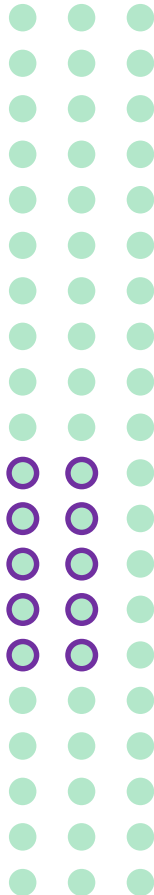
Formalization of the **concrete** semantics of the program

States of the program

```

1 // x is an unknown parameter
2 float runTheLoop(int x)
3 {
4   if (x<0 || x>100)
5     return 1.0;
6   float in = input();
7   float current_out = 0.0;
8   float prev_out = 0.0;
9   int i;
10  for (i=0;i<=x;i++)
11  {
12    current_out = in*0.1 + prev_out*0.9;
13    prev_out = current_out;
14    if (random())
15      prev_out = 0.0;
16    in = input();
17  }
18  current_out = current_out + 3.0;
19  return 1/current_out;
20 }

```



x=1
 current_out=0
 in=1.6
 i=17

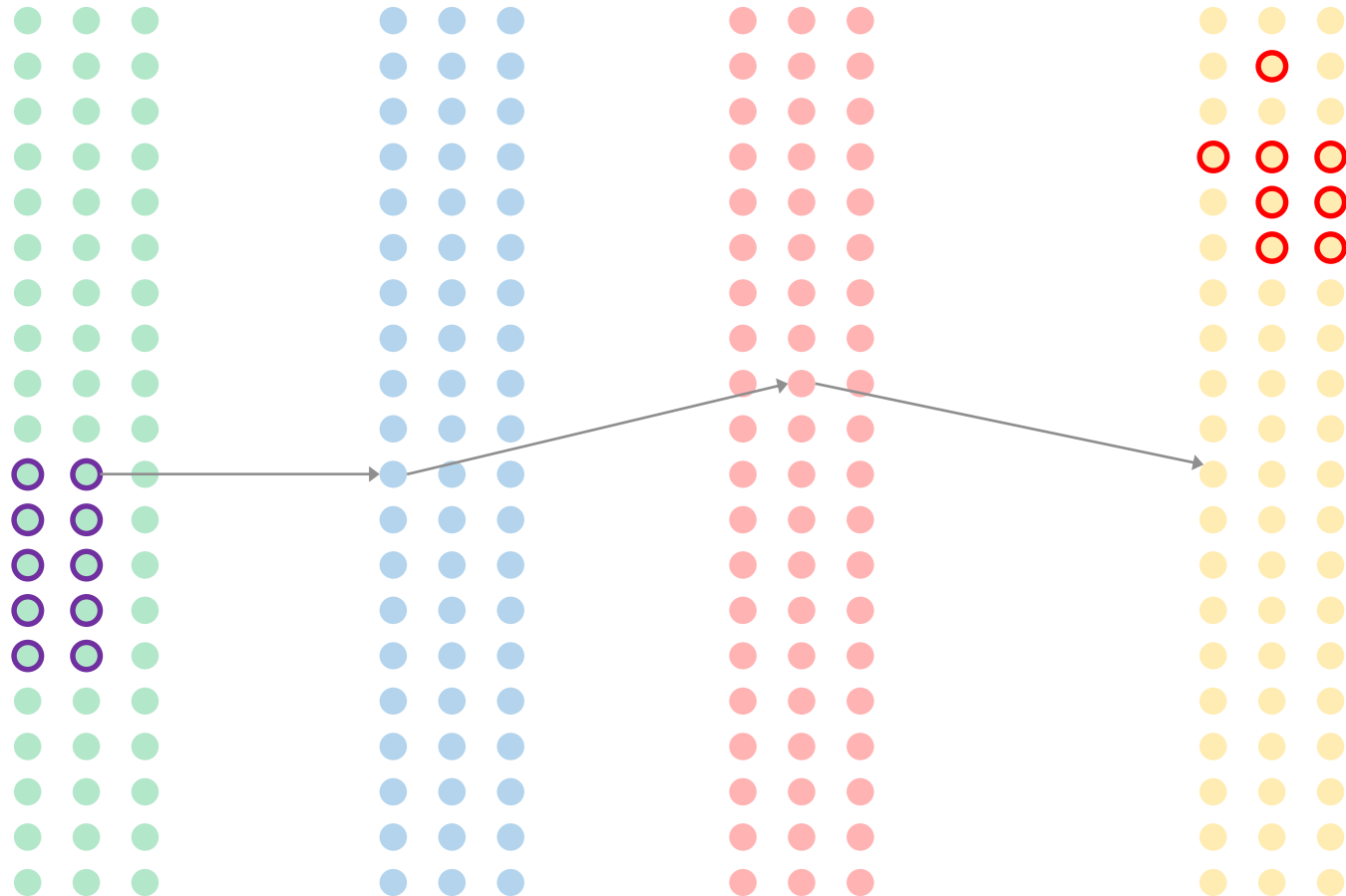
Formalization of the **concrete** semantics of the program

Traces semantics

```

1  // x is an unknown parameter
2  float runTheLoop(int x)
3  {
4    if (x<0 || x>100)
5      return 1.0;
6    float in = input();
7    float current_out = 0.0;
8    float prev_out = 0.0;
9    int i;
10   for (i=0;i<=x;i++)
11   {
12     current_out = in*0.1 + prev_out*0.9;
13     prev_out = current_out;
14     if (random())
15       prev_out = 0.0;
16     in = input();
17   }
18   current_out = current_out + 3.0;
19   return 1/current_out;
20 }

```



Infinite state machine
that simulates program
execution

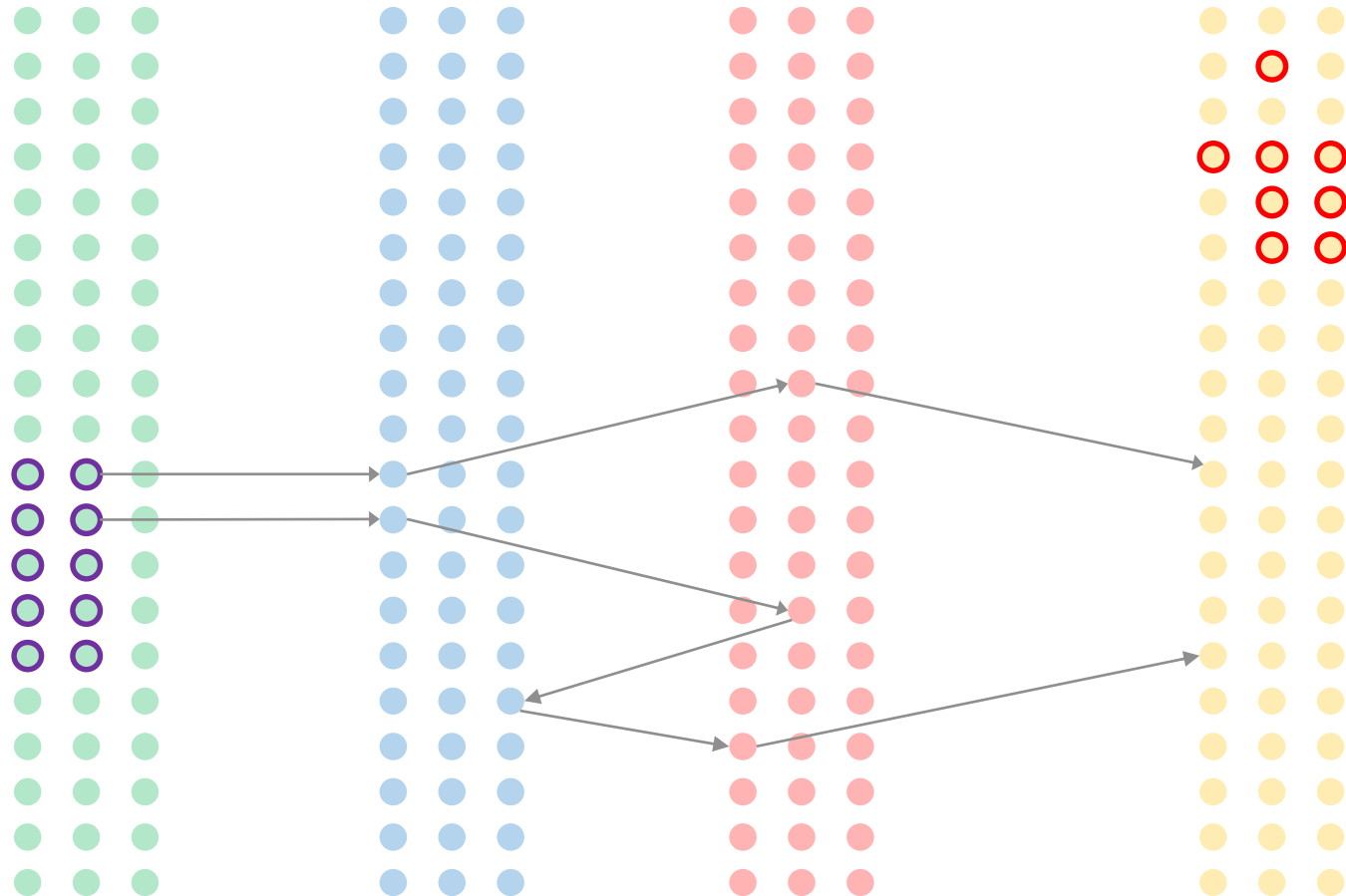
Formalization of the **concrete** semantics of the program

Traces semantics

```

1  // x is an unknown parameter
2  float runTheLoop(int x)
3  {
4    if (x<0 || x>100)
5      return 1.0;
6    float in = input();
7    float current_out = 0.0;
8    float prev_out = 0.0;
9    int i;
10   for (i=0;i<=x;i++)
11   {
12     current_out = in*0.1 + prev_out*0.9;
13     prev_out = current_out;
14     if (random())
15       prev_out = 0.0;
16     in = input();
17   }
18   current_out = current_out + 3.0;
19   return 1/current_out;
20 }

```



Infinite state machine
that simulates program
execution

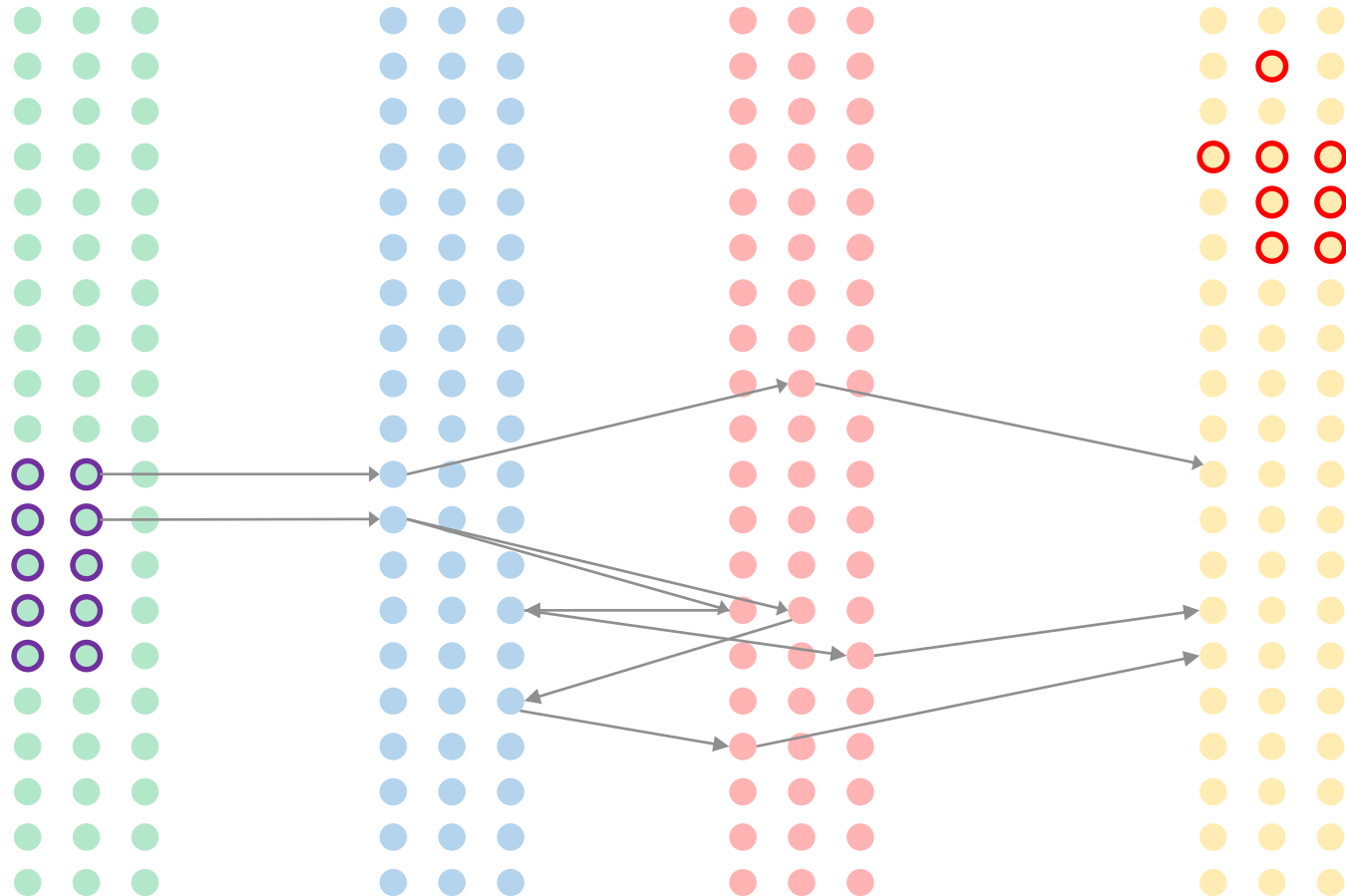
Formalization of the **concrete** semantics of the program

Traces semantics

```

1  // x is an unknown parameter
2  float runTheLoop(int x)
3  {
4  if (x<0 || x>100)
5    return 1.0;
6  float in = input();
7  float current_out = 0.0;
8  float prev_out = 0.0;
9  int i;
10 for (i=0;i<=x;i++)
11 {
12   current_out = in*0.1 + prev_out*0.9;
13   prev_out = current_out;
14   if (random())
15     prev_out = 0.0;
16   in = input();
17 }
18 current_out = current_out + 3.0;
19
20
21 return 1/current_out;
22 }

```



Infinite state machine
that simulates program
execution

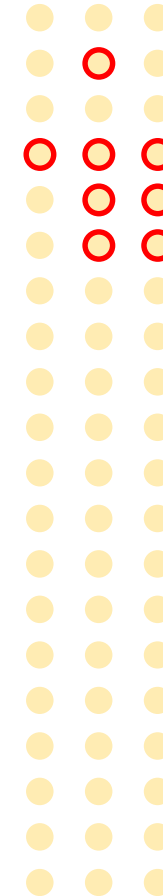
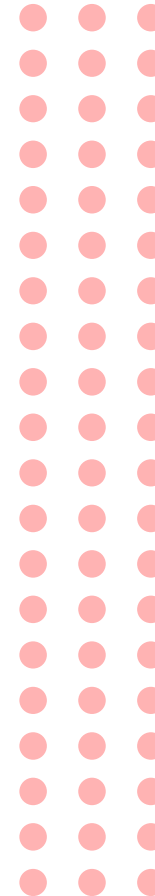
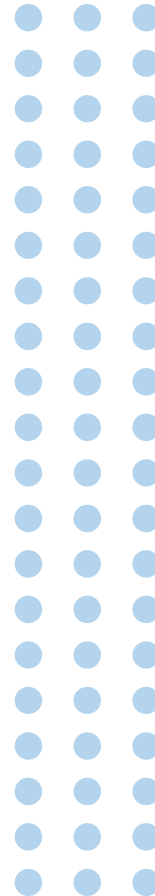
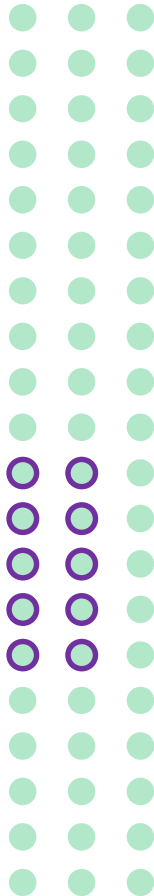
Formalization of the **collecting** semantics of the program

Invariants semantics

```

1  // x is an unknown parameter
2  float runTheLoop(int x)
3  {
4    if (x<0 || x>100)
5      return 1.0;
6    float in = input();
7    float current_out = 0.0;
8    float prev_out = 0.0;
9    int i;
10   for (i=0;i<=x;i++)
11   {
12     current_out = in*0.1 + prev_out*0.9;
13     prev_out = current_out;
14     if (random())
15       prev_out = 0.0;
16     in = input();
17   }
18   current_out = current_out + 3.0;
19   return 1/current_out;
20 }

```



Iterate until no new
states are reached

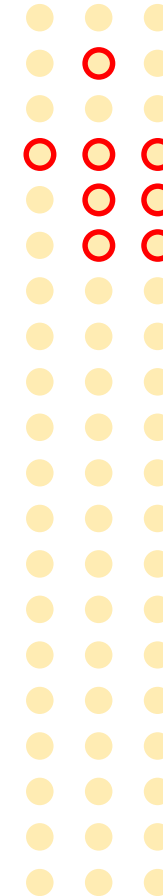
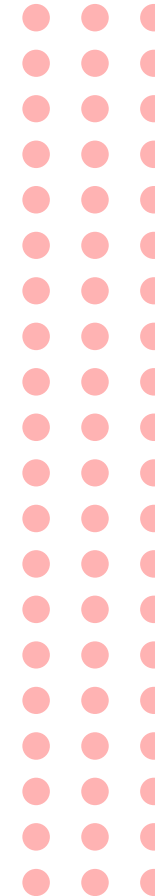
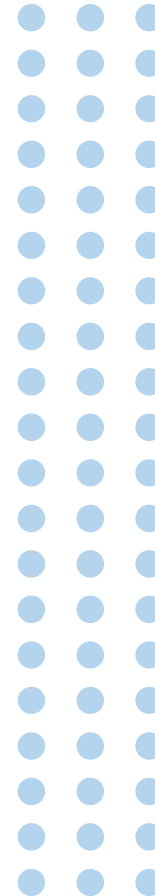
Formalization of the **collecting** semantics of the program

Invariants semantics

```

1  // x is an unknown parameter
2  float runTheLoop(int x)
3  {
4    if (x<0 || x>100)
5      return 1.0;
6    float in = input();
7    float current_out = 0.0;
8    float prev_out = 0.0;
9    int i;
10   for (i=0;i<=x;i++)
11   {
12     current_out = in*0.1 + prev_out*0.9;
13     prev_out = current_out;
14     if (random())
15       prev_out = 0.0;
16     in = input();
17   }
18   current_out = current_out + 3.0;
19   return 1/current_out;
20 }

```



Iterate until no new
states are reached

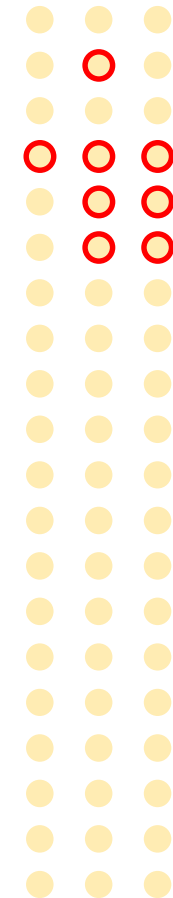
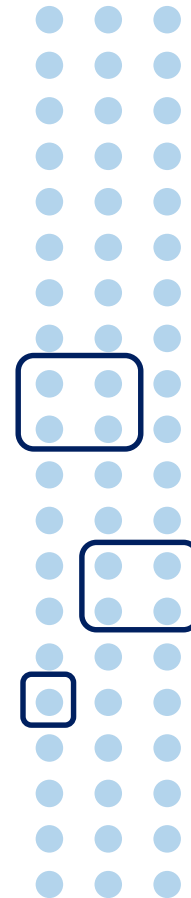
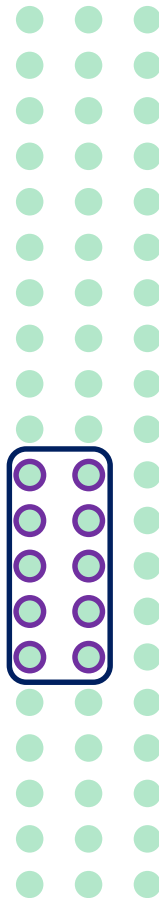
Formalization of the **collecting** semantics of the program

Invariants semantics

```

1  // x is an unknown parameter
2  float runTheLoop(int x)
3  {
4    if (x<0 || x>100)
5      return 1.0;
6    float in = input();
7    float current_out = 0.0;
8    float prev_out = 0.0;
9    int i;
10   for (i=0;i<=x;i++)
11   {
12     current_out = in*0.1 + prev_out*0.9;
13     prev_out = current_out;
14     if (random())
15       prev_out = 0.0;
16     in = input();
17   }
18   current_out = current_out + 3.0;
19   return 1/current_out;
20 }

```



Iterate until no new
states are reached

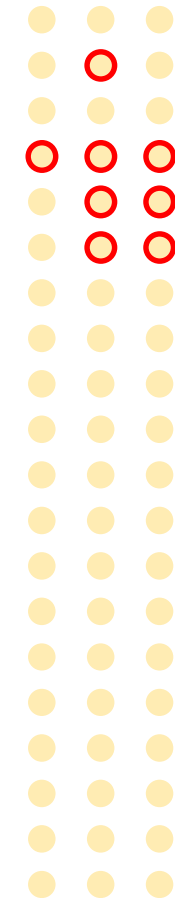
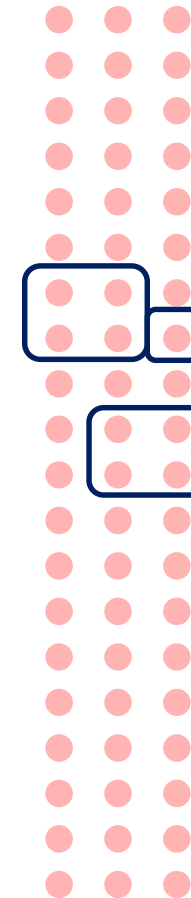
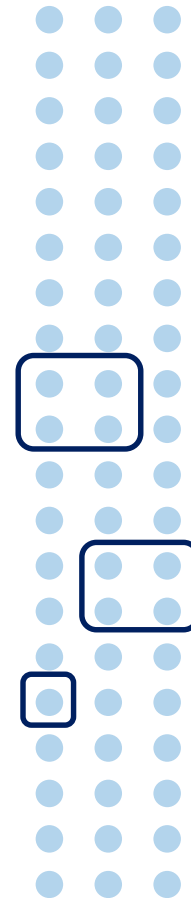
Formalization of the **collecting** semantics of the program

Invariants semantics

```

1  // x is an unknown parameter
2  float runTheLoop(int x)
3  {
4    if (x<0 || x>100)
5      return 1.0;
6    float in = input();
7    float current_out = 0.0;
8    float prev_out = 0.0;
9    int i;
10   for (i=0;i<=x;i++)
11   {
12     current_out = in*0.1 + prev_out*0.9;
13     prev_out = current_out;
14     if (random())
15       prev_out = 0.0;
16     in = input();
17   }
18   current_out = current_out + 3.0;
19   return 1/current_out;
20 }

```



Iterate until no new
states are reached

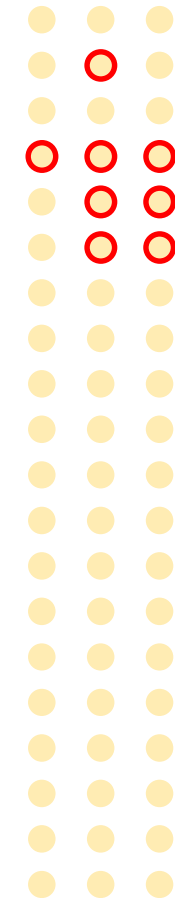
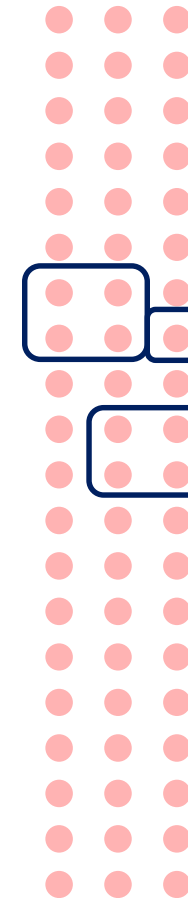
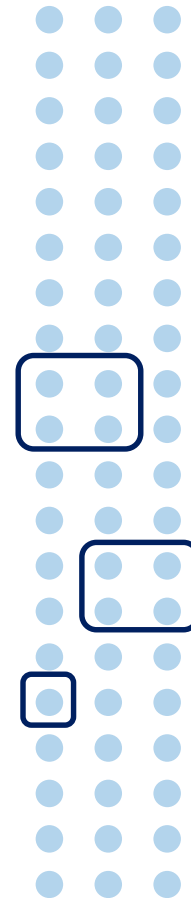
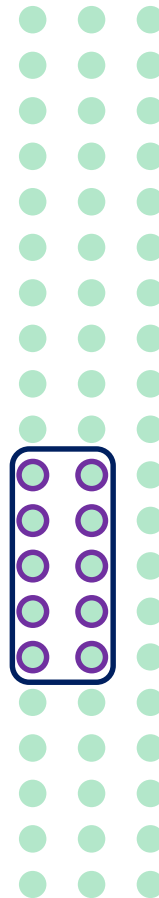
Formalization of the **collecting** semantics of the program

Invariants semantics

```

1 // x is an unknown parameter
2 float runTheLoop(int x)
3 {
4   if (x<0 || x>100)
5     return 1.0;
6   float in = input();
7   float current_out = 0.0;
8   float prev_out = 0.0;
9   int i;
10  for (i=0;i<=x;i++)
11  {
12    current_out = in*0.1 + prev_out*0.9;
13    prev_out = current_out;
14    if (random())
15      prev_out = 0.0;
16    in = input();
17  }
18  current_out = current_out + 3.0;
19  return 1/current_out;
20 }

```



Iterate until no new
states are reached

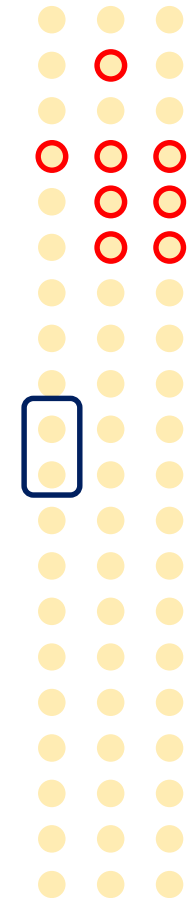
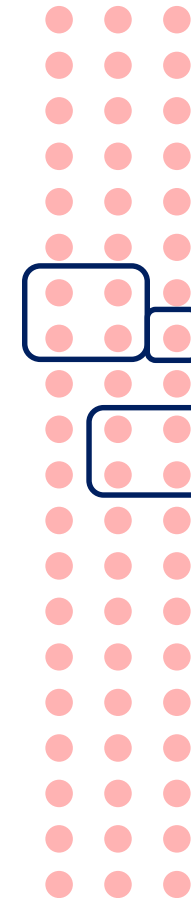
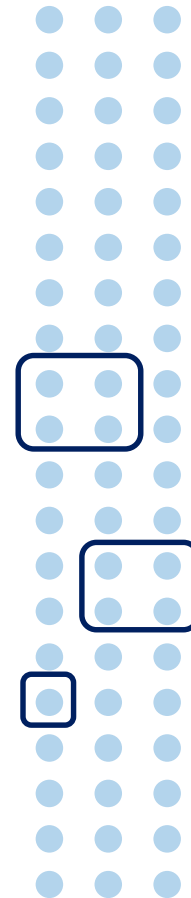
Formalization of the **collecting** semantics of the program

Invariants semantics

```

1  // x is an unknown parameter
2  float runTheLoop(int x)
3  {
4    if (x<0 || x>100)
5      return 1.0;
6    float in = input();
7    float current_out = 0.0;
8    float prev_out = 0.0;
9    int i;
10   for (i=0;i<=x;i++)
11   {
12     current_out = in*0.1 + prev_out*0.9;
13     prev_out = current_out;
14     if (random())
15       prev_out = 0.0;
16     in = input();
17   }
18   current_out = current_out + 3.0;
19   return 1/current_out;
20 }

```

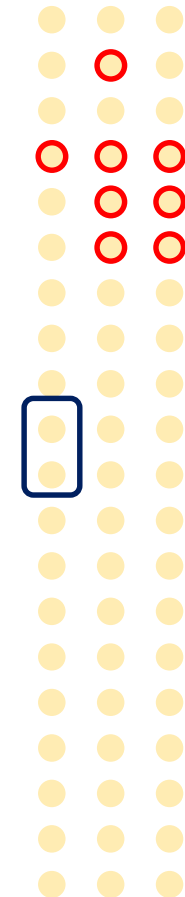
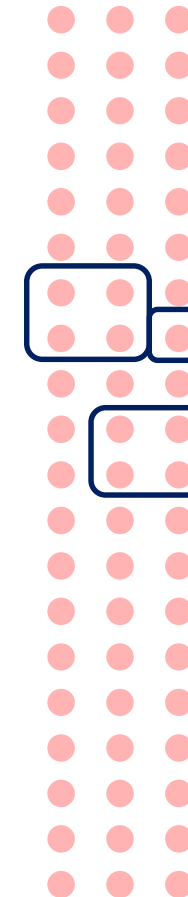
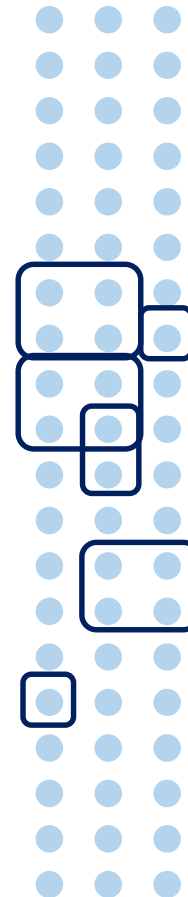
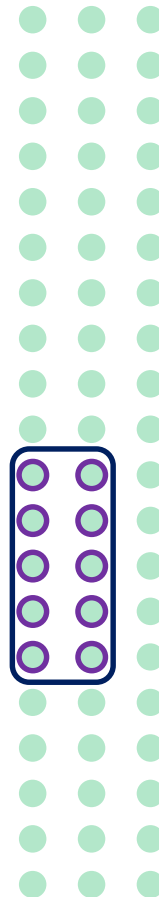


Iterate until no new
states are reached

Formalization of the **collecting** semantics of the program

Invariants semantics

```
1 // x is an unknown parameter
2 float runTheLoop(int x)
3 {
4   if (x<0 || x>100)
5     return 1.0;
6   float in = input();
7   float current_out = 0.0;
8   float prev_out = 0.0;
9   int i;
10  for (i=0;i<=x;i++)
11  {
12    current_out = in*0.1 + prev_out*0.9;
13    prev_out = current_out;
14    if (random())
15      prev_out = 0.0;
16    in = input();
17  }
18  current_out = current_out + 3.0;
19  return 1/current_out;
20 }
21
22 }
```



Iterate until no new
states are reached

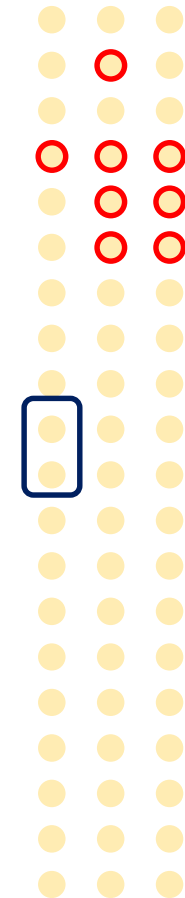
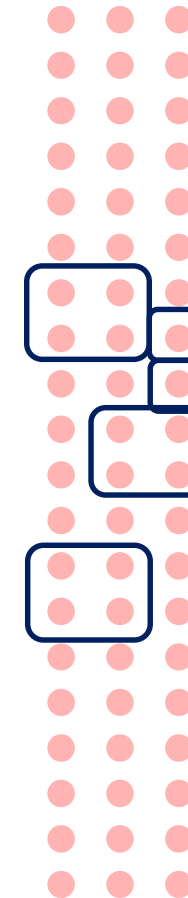
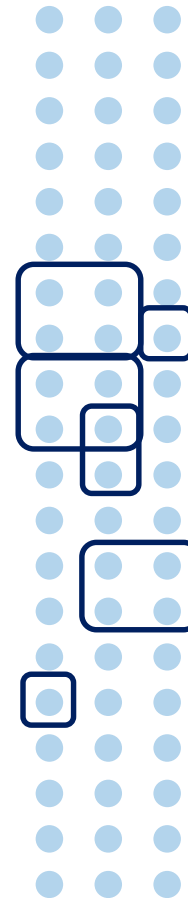
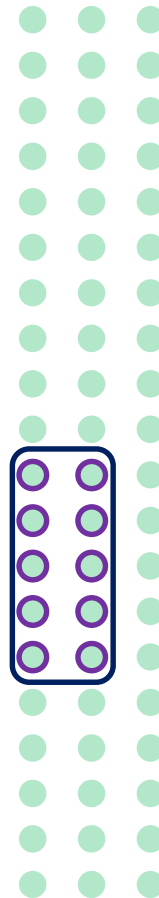
Formalization of the **collecting** semantics of the program

Invariants semantics

```

1  // x is an unknown parameter
2  float runTheLoop(int x)
3  {
4    if (x<0 || x>100)
5      return 1.0;
6    float in = input();
7    float current_out = 0.0;
8    float prev_out = 0.0;
9    int i;
10   for (i=0;i<=x;i++)
11   {
12     current_out = in*0.1 + prev_out*0.9;
13     prev_out = current_out;
14     if (random())
15       prev_out = 0.0;
16     in = input();
17   }
18   current_out = current_out + 3.0;
19   return 1/current_out;
20 }

```

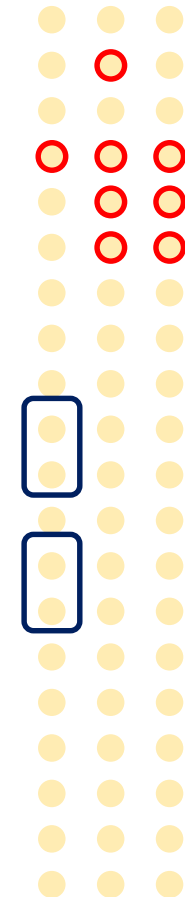
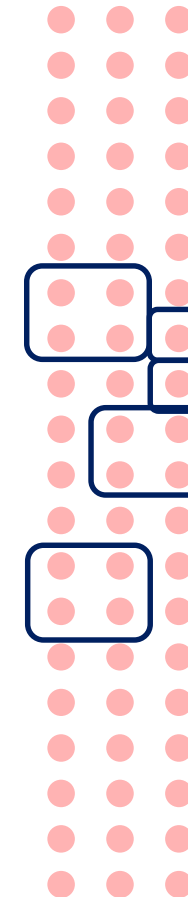
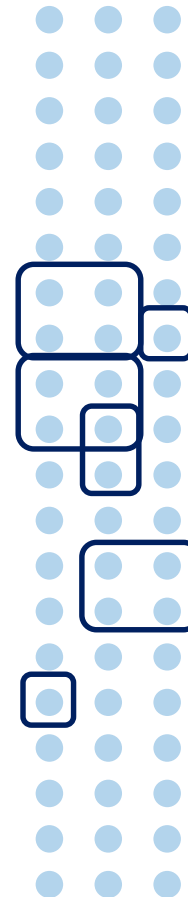
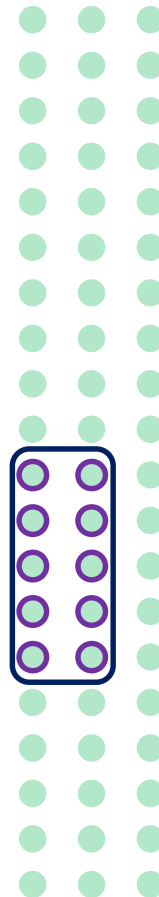


Iterate until no new
states are reached

Formalization of the **collecting** semantics of the program

Invariants semantics

```
1 // x is an unknown parameter
2 float runTheLoop(int x)
3 {
4   if (x<0 || x>100)
5     return 1.0;
6   float in = input();
7   float current_out = 0.0;
8   float prev_out = 0.0;
9   int i;
10  for (i=0;i<=x;i++)
11  {
12    current_out = in*0.1 + prev_out*0.9;
13    prev_out = current_out;
14    if (random())
15      prev_out = 0.0;
16    in = input();
17  }
18  current_out = current_out + 3.0;
19  return 1/current_out;
20 }
21
22 }
```



Iterate until no new
states are reached

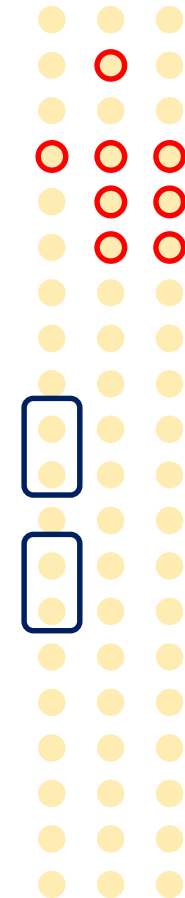
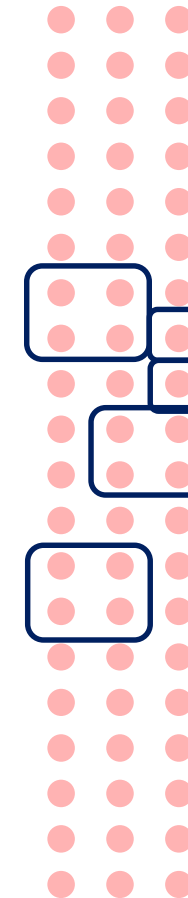
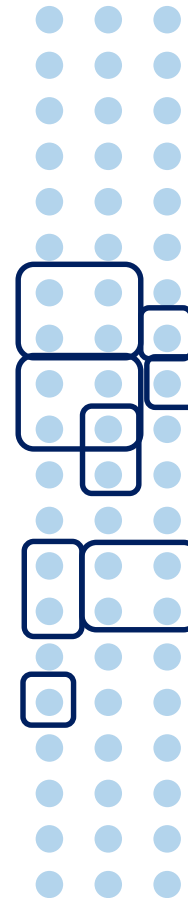
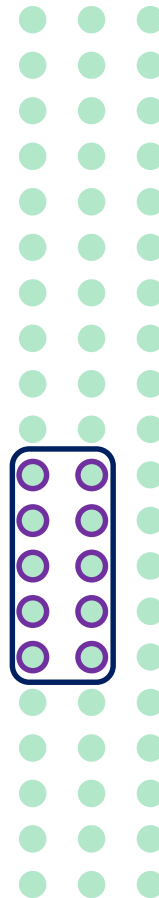
Formalization of the **collecting** semantics of the program

Invariants semantics

```

1  // x is an unknown parameter
2  float runTheLoop(int x)
3  {
4    if (x<0 || x>100)
5      return 1.0;
6    float in = input();
7    float current_out = 0.0;
8    float prev_out = 0.0;
9    int i;
10   for (i=0;i<=x;i++)
11   {
12     current_out = in*0.1 + prev_out*0.9;
13     prev_out = current_out;
14     if (random())
15       prev_out = 0.0;
16     in = input();
17   }
18   current_out = current_out + 3.0;
19   return 1/current_out;
20 }

```



Iterate until no new
states are reached

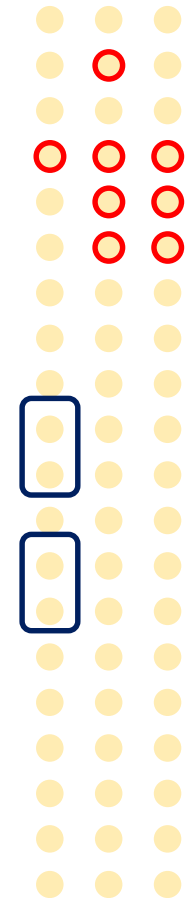
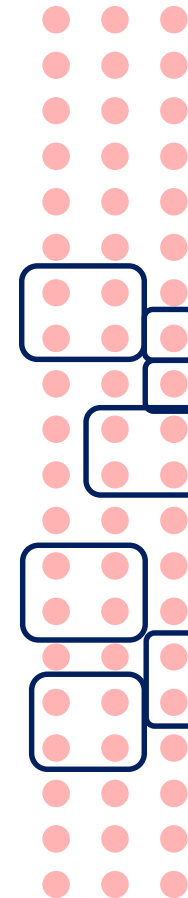
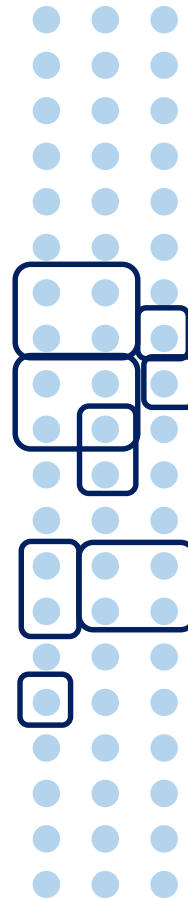
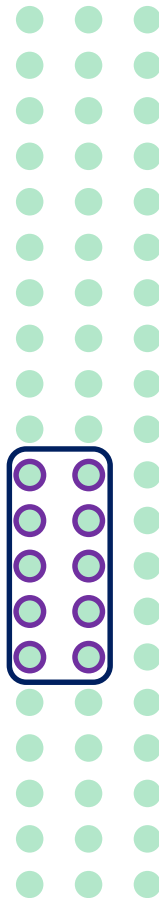
Formalization of the **collecting** semantics of the program

Invariants semantics

```

1  // x is an unknown parameter
2  float runTheLoop(int x)
3  {
4    if (x<0 || x>100)
5      return 1.0;
6    float in = input();
7    float current_out = 0.0;
8    float prev_out = 0.0;
9    int i;
10   for (i=0;i<=x;i++)
11   {
12     current_out = in*0.1 + prev_out*0.9;
13     prev_out = current_out;
14     if (random())
15       prev_out = 0.0;
16     in = input();
17   }
18   current_out = current_out + 3.0;
19   return 1/current_out;
20 }

```



Iterate until no new
states are reached

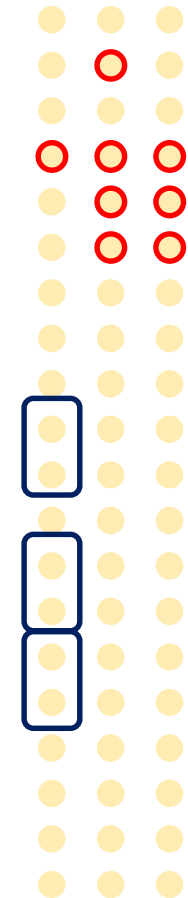
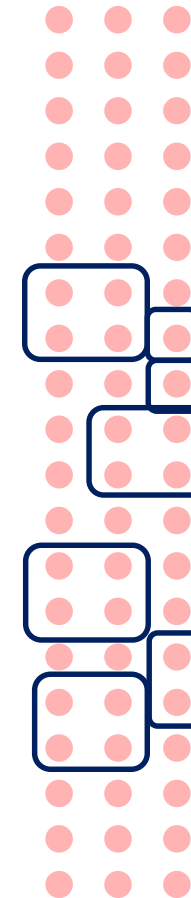
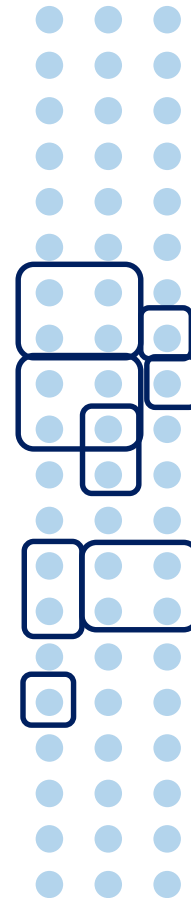
Formalization of the **collecting** semantics of the program

Invariants semantics

```

1  // x is an unknown parameter
2  float runTheLoop(int x)
3  {
4    if (x<0 || x>100)
5      return 1.0;
6    float in = input();
7    float current_out = 0.0;
8    float prev_out = 0.0;
9    int i;
10   for (i=0;i<=x;i++)
11   {
12     current_out = in*0.1 + prev_out*0.9;
13     prev_out = current_out;
14     if (random())
15       prev_out = 0.0;
16     in = input();
17   }
18   current_out = current_out + 3.0;
19   return 1/current_out;
20 }

```



Iterate until no new
states are reached

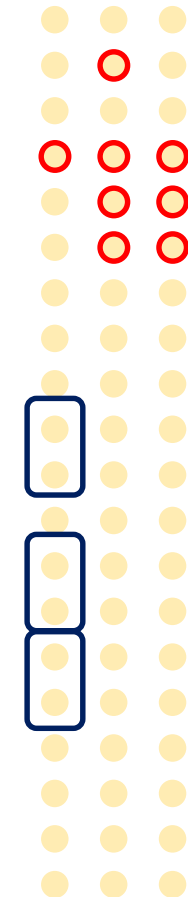
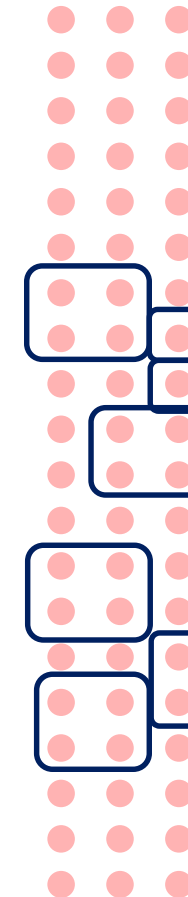
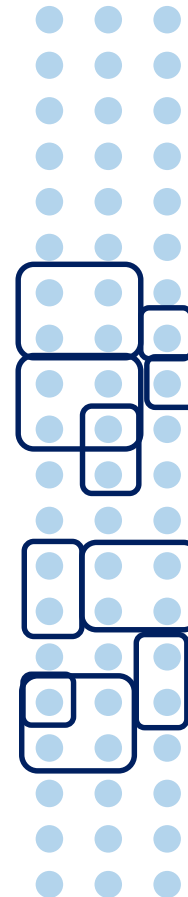
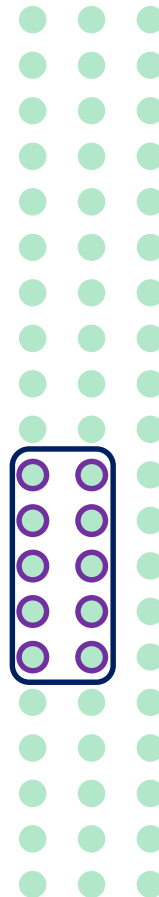
Formalization of the **collecting** semantics of the program

Invariants semantics

```

1  // x is an unknown parameter
2  float runTheLoop(int x)
3  {
4    if (x<0 || x>100)
5      return 1.0;
6    float in = input();
7    float current_out = 0.0;
8    float prev_out = 0.0;
9    int i;
10   for (i=0;i<=x;i++)
11   {
12     current_out = in*0.1 + prev_out*0.9;
13     prev_out = current_out;
14     if (random())
15       prev_out = 0.0;
16     in = input();
17   }
18   current_out = current_out + 3.0;
19   return 1/current_out;
20 }

```



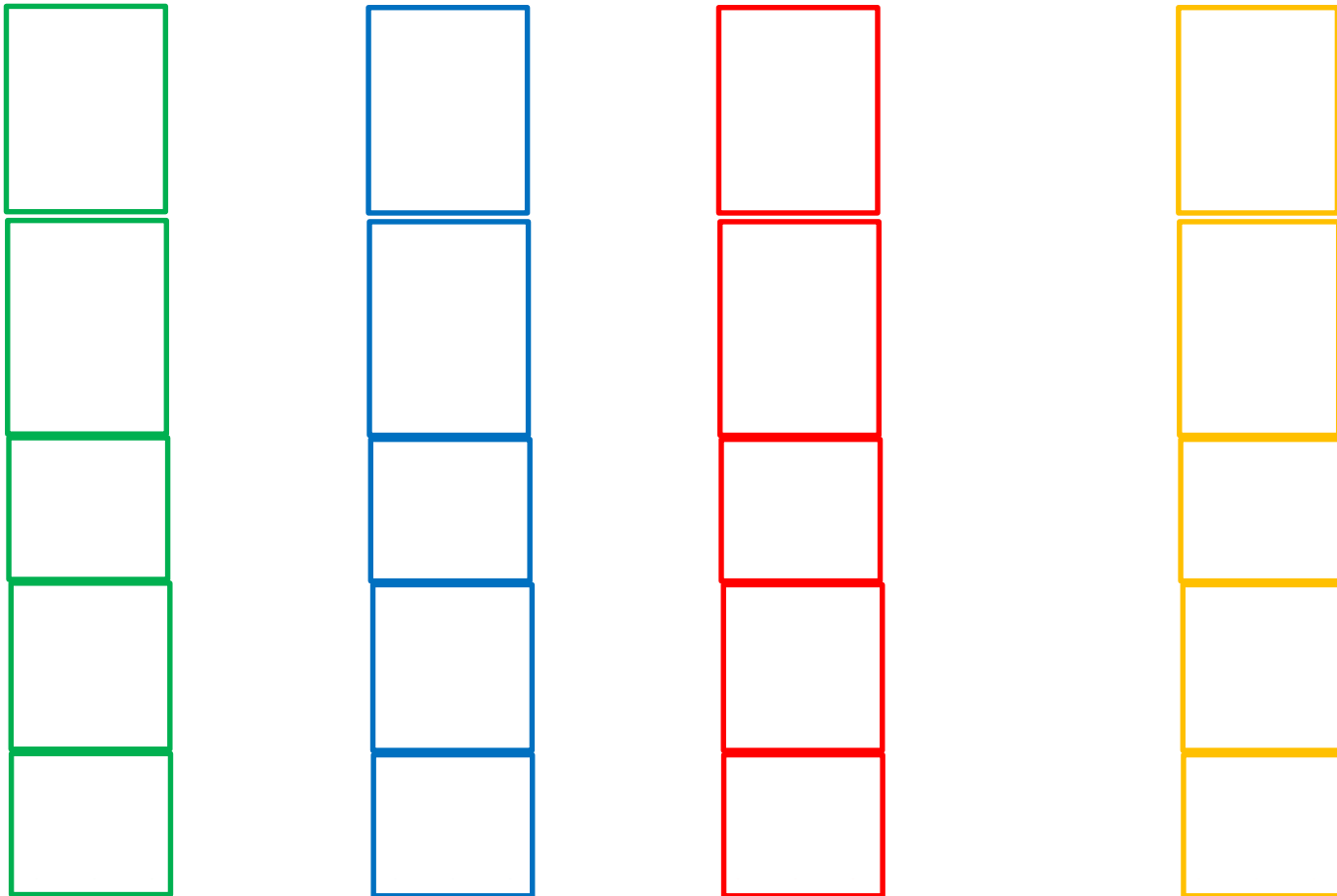
Iterate until no new
states are reached

Formalization of the **abstract** semantics of the program

Properties semantics – Loss of expressivity

```
1 // x is an unknown parameter
2 float runTheLoop(int x)
3 {
4   if (x<0 || x>100)
5     return 1.0;
6   float in = input();
7   float current_out = 0.0;
8   float prev_out = 0.0;
9   int i;
10  for (i=0;i<=x;i++)
11  {
12    current_out = in*0.1 + prev_out*0.9;
13    prev_out = current_out;
14    if (random())
15      prev_out = 0.0;
16    in = input();
17  }
18
19  current_out = current_out + 3.0;
20
21  return 1/current_out;
22 }
```

Finite state machine
that **over-approximates**
program execution

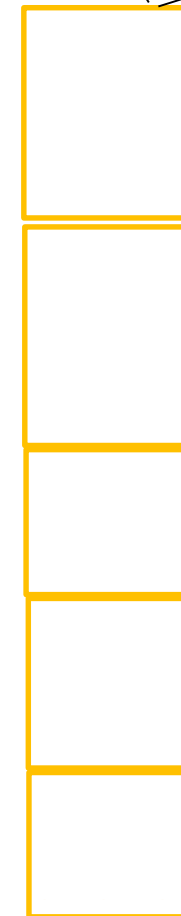
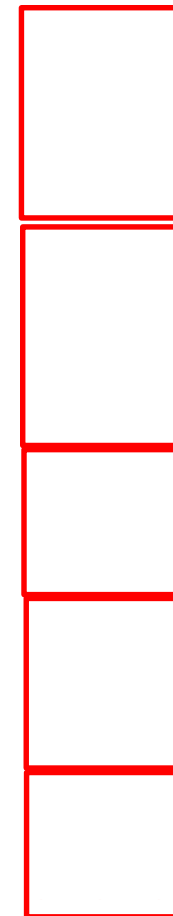
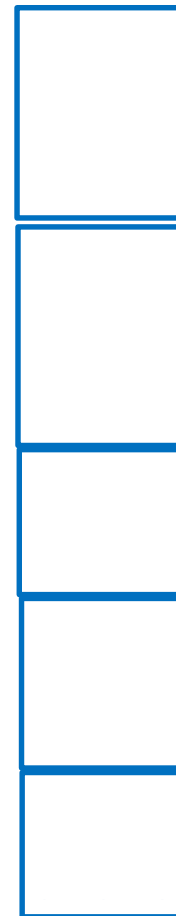
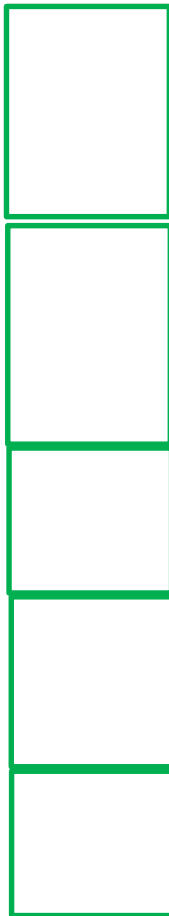


Formalization of the **abstract** semantics of the program

Properties semantics – Loss of expressivity

```
1 // x is an unknown parameter
2 float runTheLoop(int x)
3 {
4   if (x<0 || x>100)
5     return 1.0;
6   float in = input();
7   float current_out = 0.0;
8   float prev_out = 0.0;
9   int i;
10  for (i=0;i<=x;i++)
11  {
12    current_out = in*0.1 + prev_out*0.9;
13    prev_out = current_out;
14    if (random())
15      prev_out = 0.0;
16    in = input();
17  }
18
19  current_out = current_out + 3.0;
20
21  return 1/current_out;
22 }
```

Finite state machine
that **over-approximates**
program execution



$X \geq 0$
 $\text{current_out} \leq 0$
 $\text{in} \geq 0$
 $i \geq 0$

Formalization of the **abstract** semantics of the program

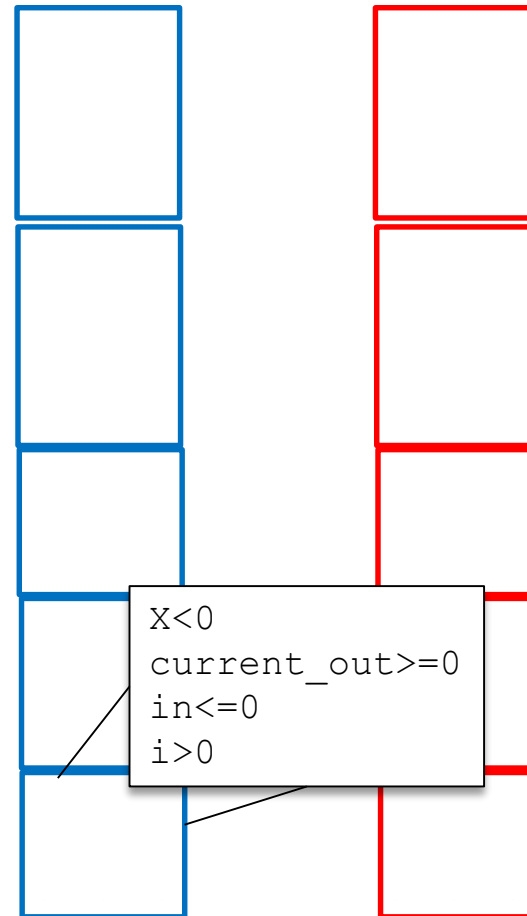
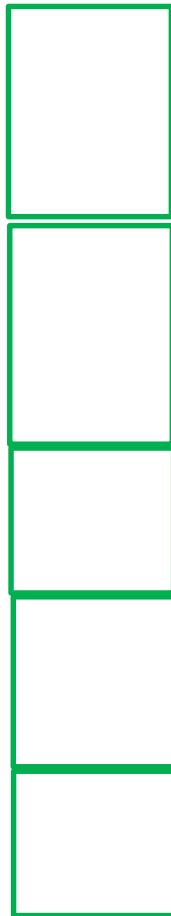
Properties semantics – Loss of expressivity

```

1 // x is an unknown parameter
2 float runTheLoop(int x)
3 {
4   if (x<0 || x>100)
5     return 1.0;
6   float in = input();
7   float current_out = 0.0;
8   float prev_out = 0.0;
9   int i;
10  for (i=0;i<=x;i++)
11  {
12    current_out = in*0.1 + prev_out*0.9;
13    prev_out = current_out;
14    if (random())
15      prev_out = 0.0;
16    in = input();
17  }
18  current_out = current_out + 3.0;
19  return 1/current_out;
20 }

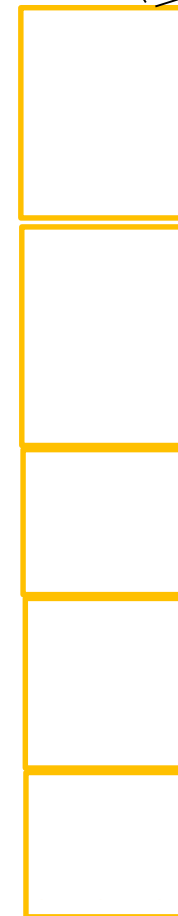
```

Finite state machine
that **over-approximates**
program execution



$X < 0$
 $current_out \geq 0$
 $in \leq 0$
 $i > 0$

$X \geq 0$
 $current_out \leq 0$
 $in \geq 0$
 $i \geq 0$

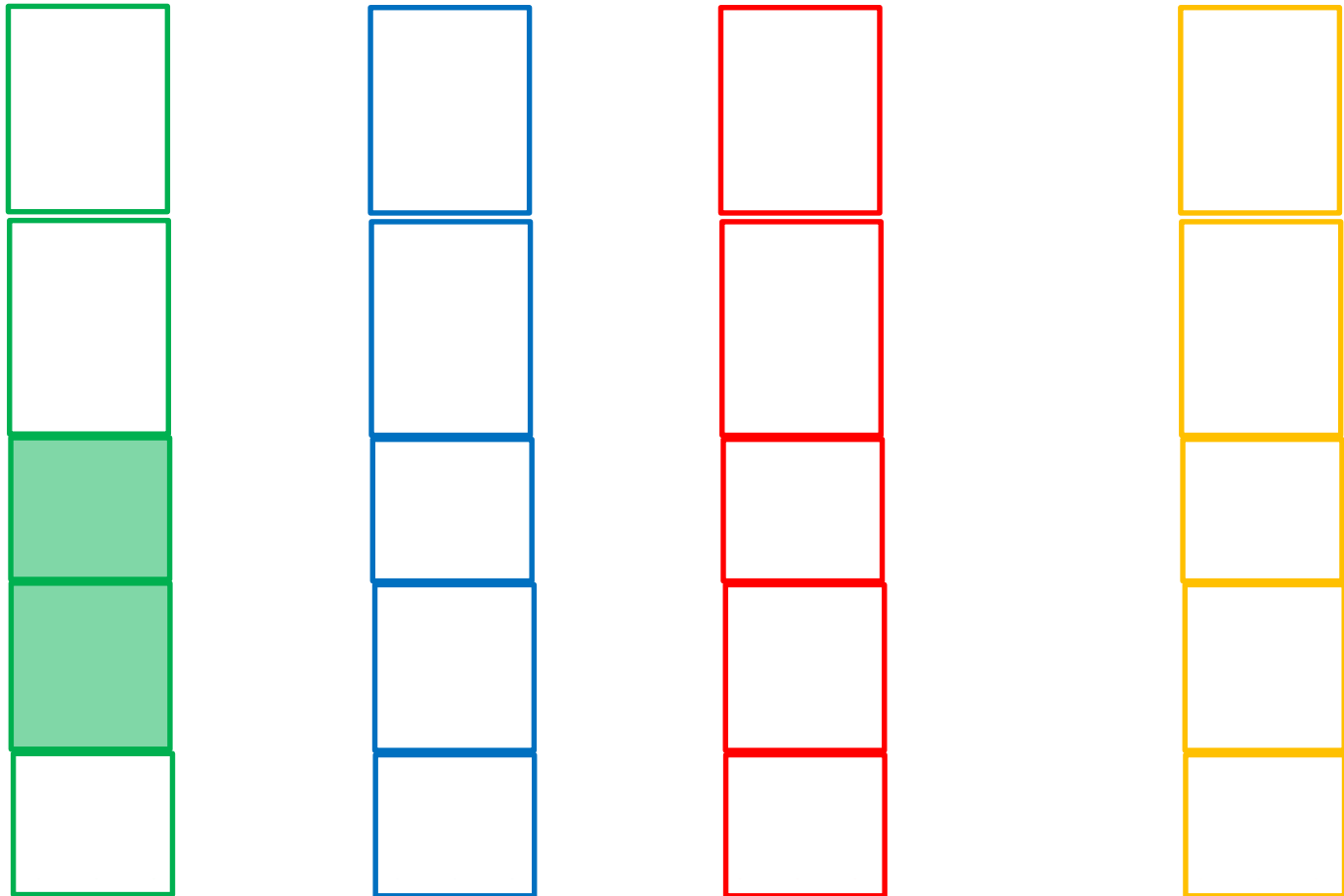


Formalization of the **abstract** semantics of the program

Properties semantics – Fixpoint computation

```
1 // x is an unknown parameter
2 float runTheLoop(int x)
3 {
4   if (x<0 || x>100)
5     return 1.0;
6   float in = input();
7   float current_out = 0.0;
8   float prev_out = 0.0;
9   int i;
10  for (i=0;i<=x;i++)
11  {
12    current_out = in*0.1 + prev_out*0.9;
13    prev_out = current_out;
14    if (random())
15      prev_out = 0.0;
16    in = input();
17  }
18  current_out = current_out + 3.0;
19  return 1/current_out;
20 }
21
22 }
```

Finite state machine
that **over-approximates**
program execution

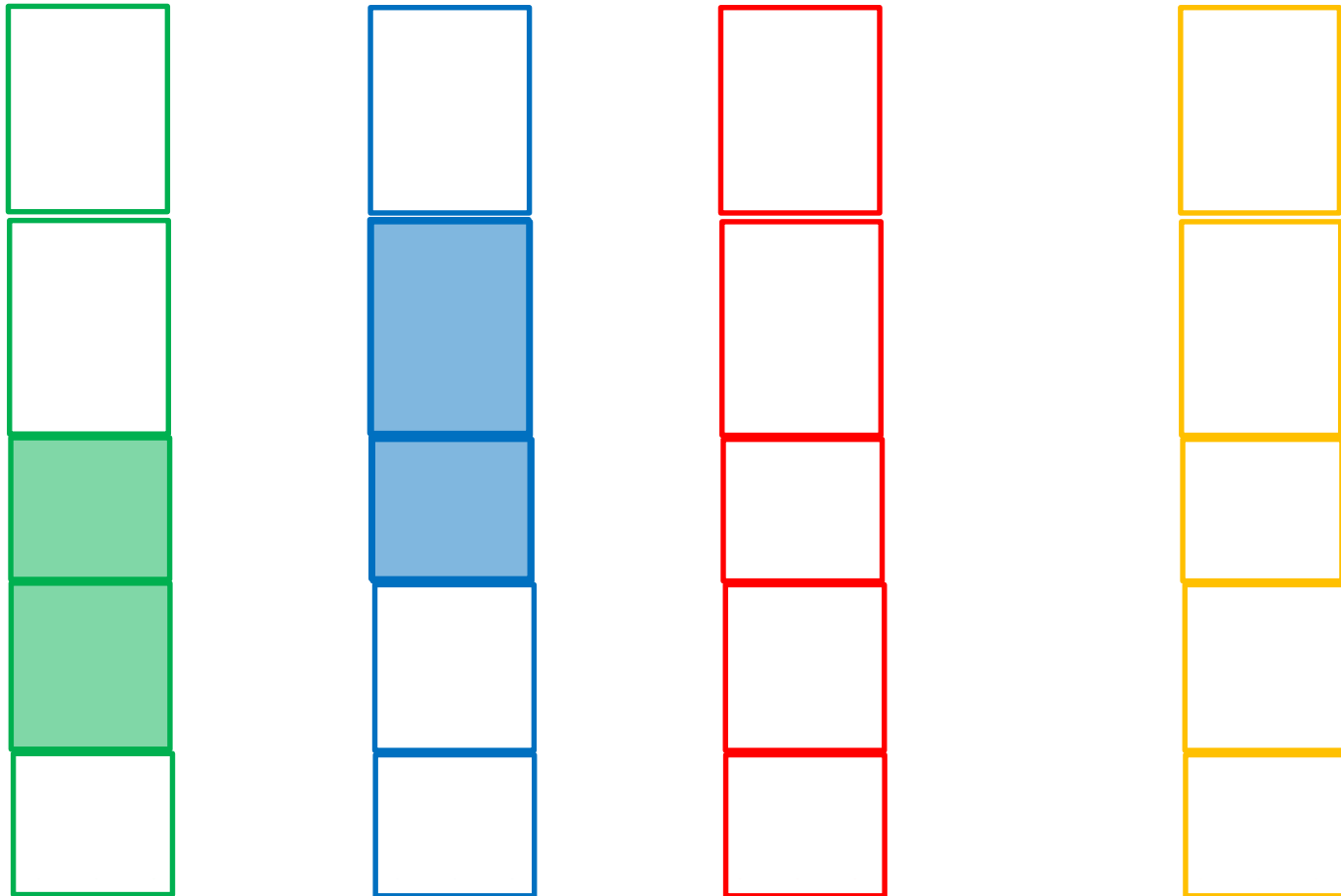


Formalization of the **abstract** semantics of the program

Properties semantics – Fixpoint computation

```
1 // x is an unknown parameter
2 float runTheLoop(int x)
3 {
4   if (x<0 || x>100)
5     return 1.0;
6   float in = input();
7   float current_out = 0.0;
8   float prev_out = 0.0;
9   int i;
10  for (i=0;i<=x;i++)
11  {
12    current_out = in*0.1 + prev_out*0.9;
13    prev_out = current_out;
14    if (random())
15      prev_out = 0.0;
16    in = input();
17  }
18  current_out = current_out + 3.0;
19  return 1/current_out;
20 }
21
22 }
```

Finite state machine
that **over-approximates**
program execution

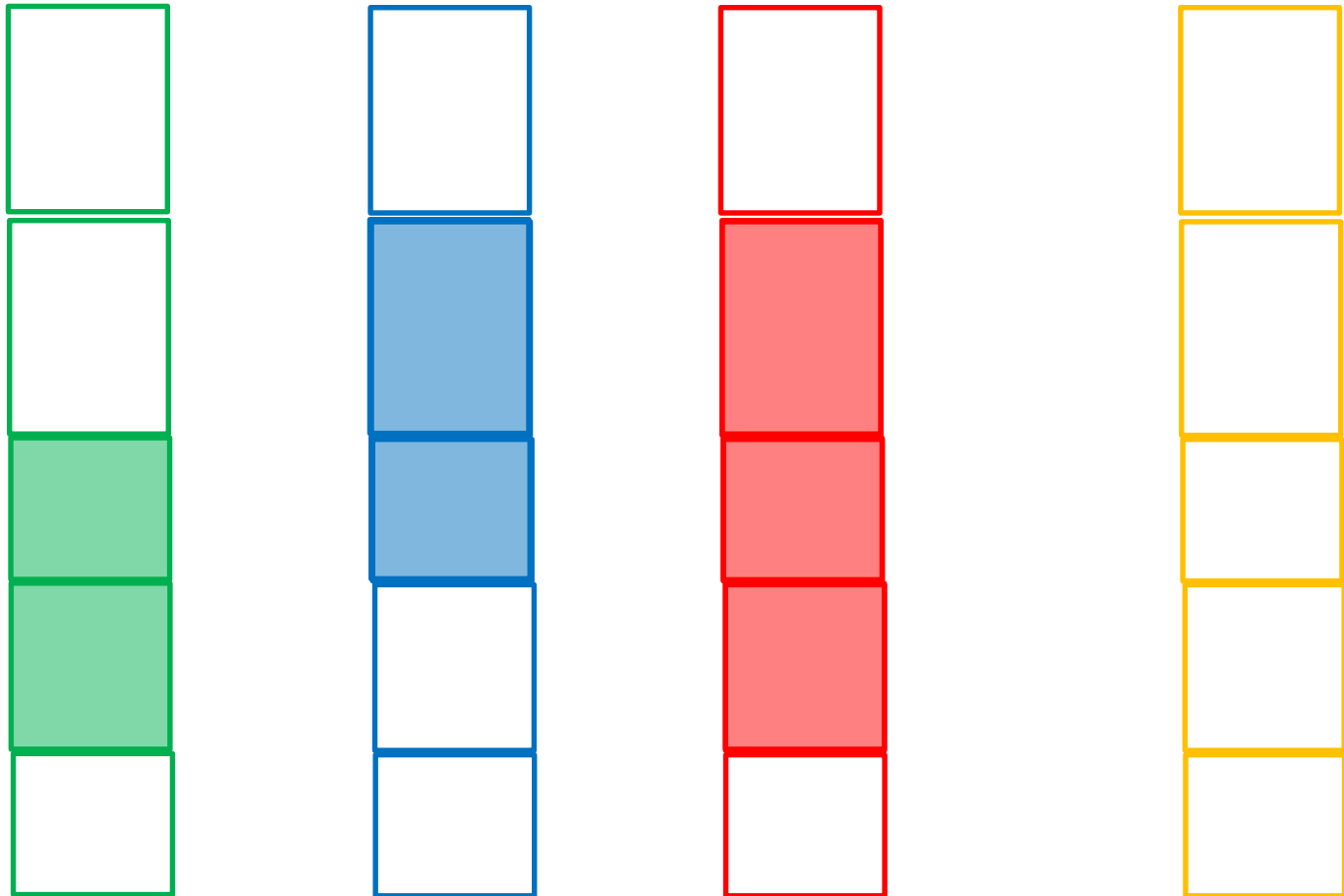


Formalization of the **abstract** semantics of the program

Properties semantics – Fixpoint computation

```
1 // x is an unknown parameter
2 float runTheLoop(int x)
3 {
4   if (x<0 || x>100)
5     return 1.0;
6   float in = input();
7   float current_out = 0.0;
8   float prev_out = 0.0;
9   int i;
10  for (i=0;i<=x;i++)
11  {
12    current_out = in*0.1 + prev_out*0.9;
13    prev_out = current_out;
14    if (random())
15      prev_out = 0.0;
16    in = input();
17  }
18
19  current_out = current_out + 3.0;
20
21  return 1/current_out;
22 }
```

Finite state machine
that **over-approximates**
program execution

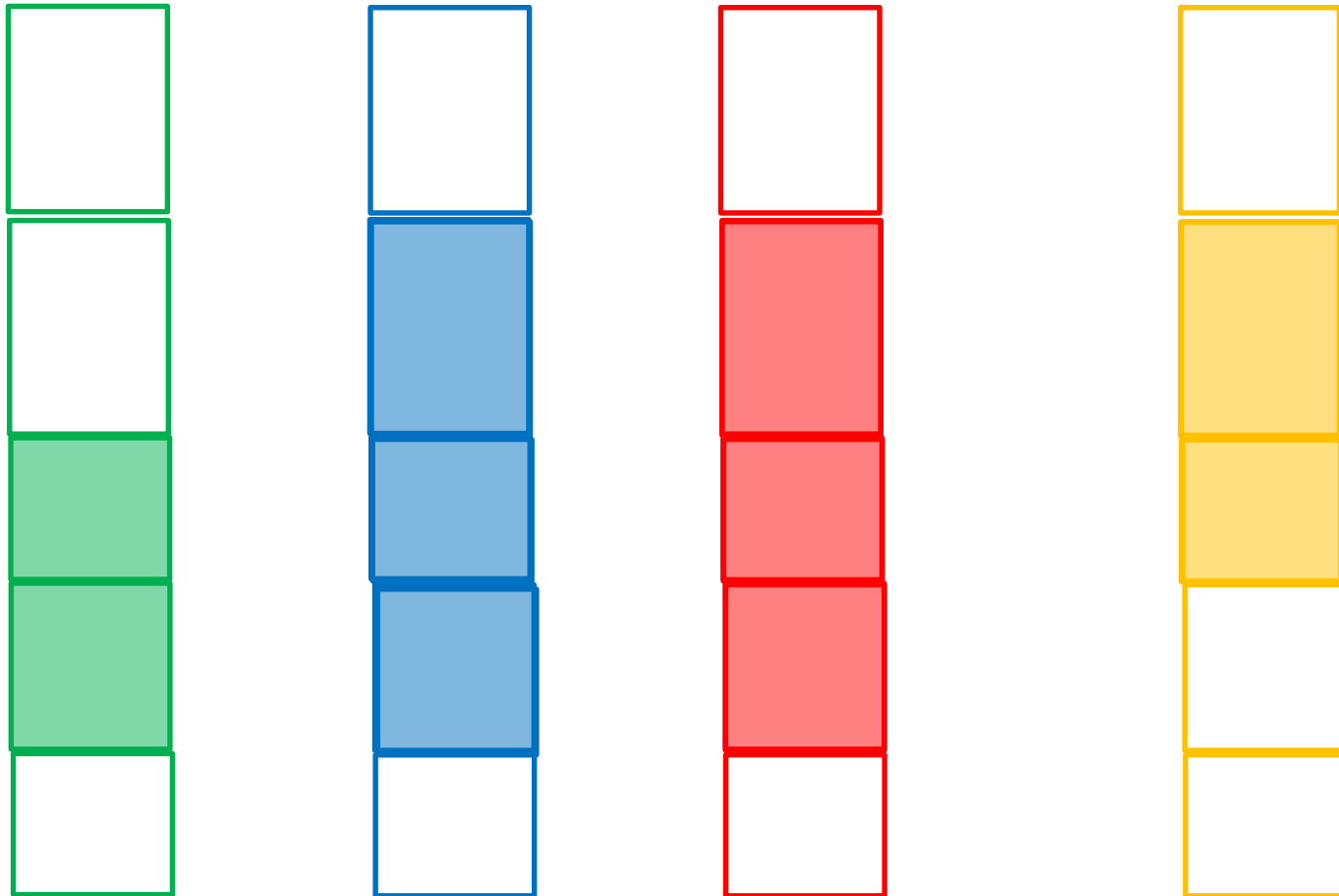


Formalization of the **abstract** semantics of the program

Properties semantics – Fixpoint computation

```
1 // x is an unknown parameter
2 float runTheLoop(int x)
3 {
4   if (x<0 || x>100)
5     return 1.0;
6   float in = input();
7   float current_out = 0.0;
8   float prev_out = 0.0;
9   int i;
10  for (i=0;i<=x;i++)
11  {
12    current_out = in*0.1 + prev_out*0.9;
13    prev_out = current_out;
14    if (random())
15      prev_out = 0.0;
16    in = input();
17  }
18  current_out = current_out + 3.0;
19  return 1/current_out;
20 }
21
22 }
```

Finite state machine
that **over-approximates**
program execution



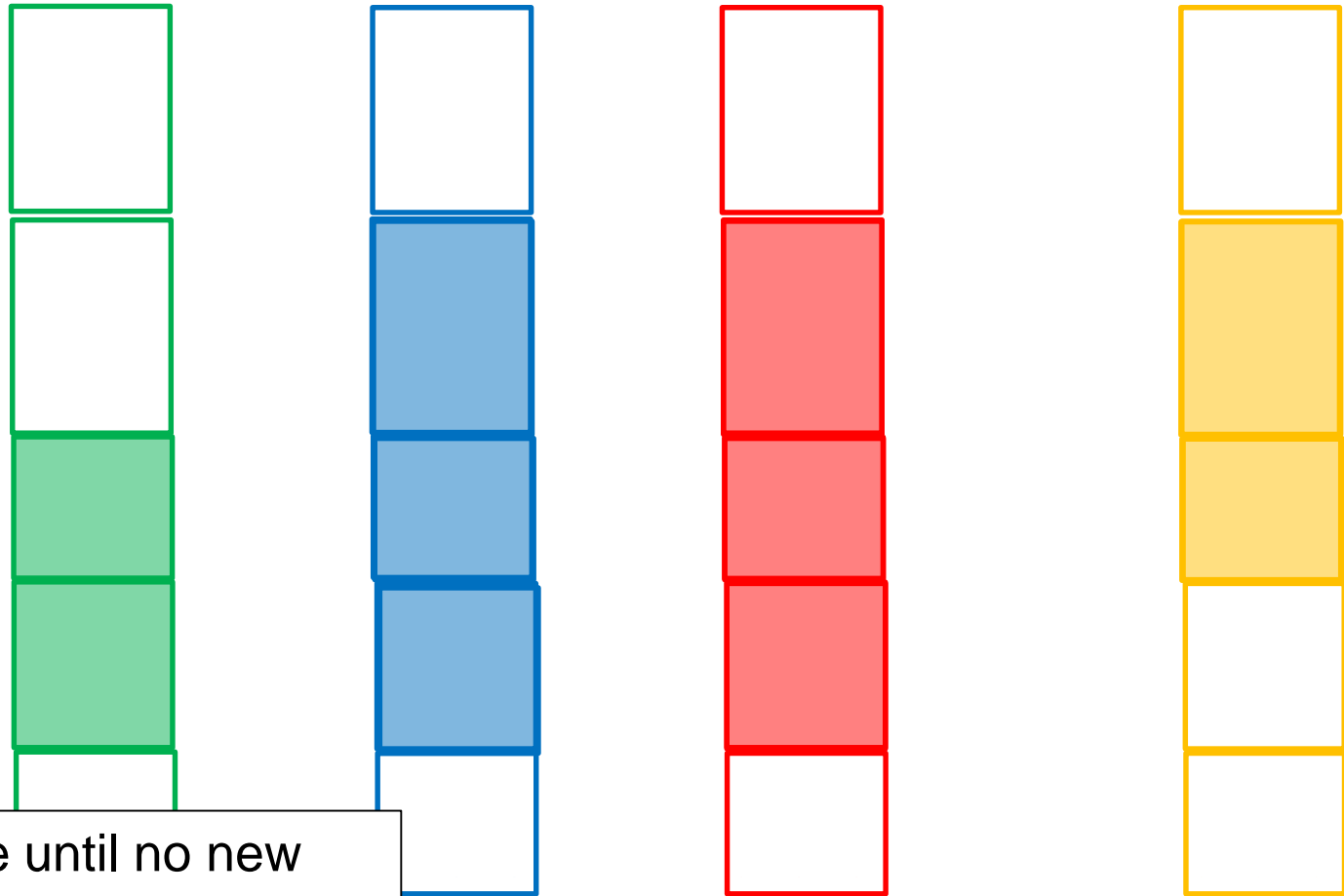
Formalization of the **abstract** semantics of the program

Properties semantics – Fixpoint computation

```
1 // x is an unknown parameter
2 float runTheLoop(int x)
3 {
4     if (x<0 || x>100)
5         return 1.0;
6     float in = input();
7     float current_out = 0.0;
8     float prev_out = 0.0;
9     int i;
10    for (i=0;i<=x;i++)
11    {
12        current_out = in*0.1 + prev_out*0.9;
13        prev_out = current_out;
14        if (random())
15            prev_out = 0.0;
16        in = input();
17    }
18    current_out = current_out + 3.0;
19    return 1/current_out;
20 }
21
22 }
```

Finite state machine
that **over-approximates**
program execution

Iterate until no new
states are reached

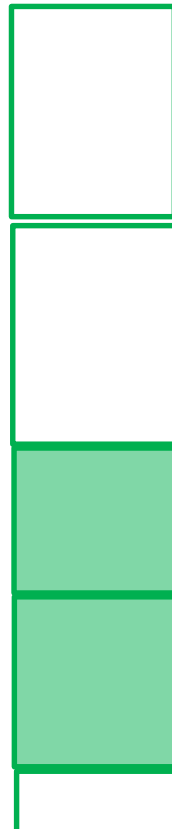


Formalization of the **abstract** semantics of the program

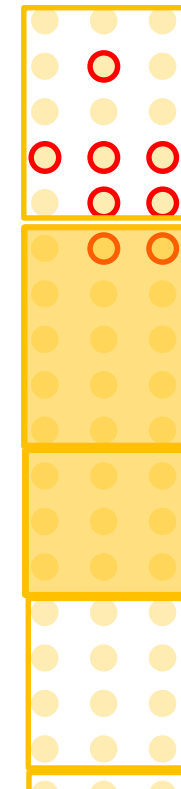
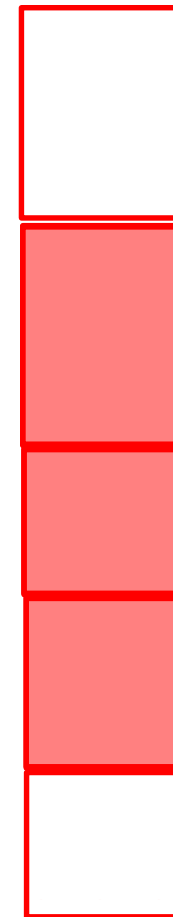
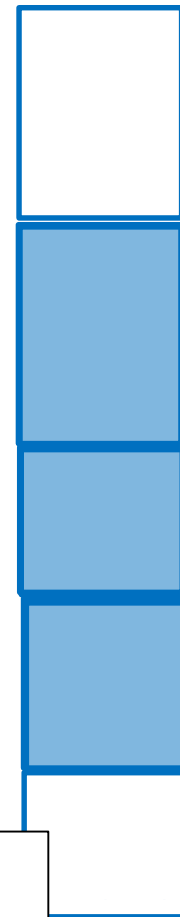
Properties semantics – Fixpoint computation

```
1 // x is an unknown parameter
2 float runTheLoop(int x)
3 {
4   if (x<0 || x>100)
5     return 1.0;
6   float in = input();
7   float current_out = 0.0;
8   float prev_out = 0.0;
9   int i;
10  for (i=0;i<=x;i++)
11  {
12    current_out = in*0.1 + prev_out*0.9;
13    prev_out = current_out;
14    if (random())
15      prev_out = 0.0;
16    in = input();
17  }
18  current_out = current_out + 3.0;
19  return 1/current_out;
20 }
21
22 }
```

Finite state machine
that **over-approximates**
program execution



Iterate until no new
states are reached



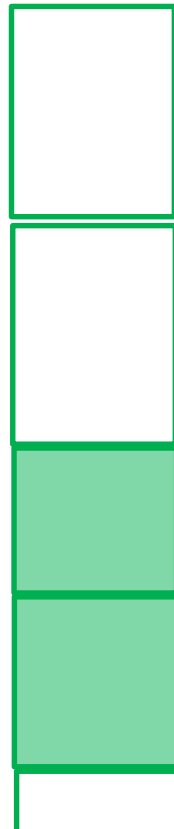
Abstract properties
represent infinitely
many states

Formalization of the **abstract** semantics of the program

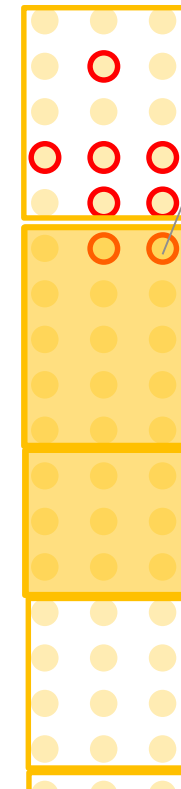
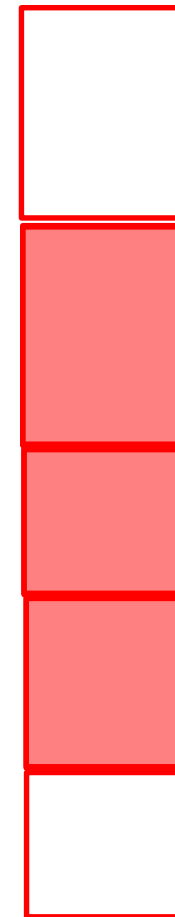
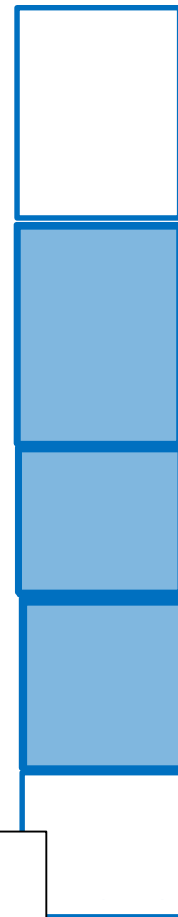
Properties semantics – Fixpoint computation

```
1 // x is an unknown parameter
2 float runTheLoop(int x)
3 {
4   if (x<0 || x>100)
5     return 1.0;
6   float in = input();
7   float current_out = 0.0;
8   float prev_out = 0.0;
9   int i;
10  for (i=0;i<=x;i++)
11  {
12    current_out = in*0.1 + prev_out*0.9;
13    prev_out = current_out;
14    if (random())
15      prev_out = 0.0;
16    in = input();
17  }
18  current_out = current_out + 3.0;
19  return 1/current_out;
20 }
21
22 }
```

Finite state machine
that **over-approximates**
program execution



Iterate until no new
states are reached



False alarm

Abstract properties
represent infinitely
many states

Applications of Abstract Interpretation

- Analysis of C Code
 - Code prover
 - Bug finder

Code Prover

- Performs abstract invariants computation on C/C++/ADA programs
 - With pointers
 - With floating points
 - With function calls
 - With multi-tasking
- Zero False Positives
- Computes many properties on program variables
 - Ranges (interval domains) with holes
 - Multiplicity (both integer and floating point)
 - Linear relations between program variables
 - e.g. $3x - 2y + 4z \leq 32 \wedge y + 3z - t \leq 1$

Local variable 'z' (unsigned int 32): multiples of 16 in [0 .. 15984] or [32000 .. 65520]

Simple color code for presenting the information

Green: reliable
safe pointer access

Red: faulty
out of bounds error

Gray: dead
unreachable code

Orange: unproven
may be unsafe for some
conditions

Purple: violation
MISRA-C/C++ or JSF++
code rules

Range data
tool tip

```
static void pointer_arithmetic (void) {
    int array[100];
    int *p = array;
    int i;

    for (i = 0; i < 100; i++) {
        *p = 0;
        p++;
    }

    if (get_bus_status() > 0) {
        if (get_oil_pressure() > 0) {
            *p = 5;
        } else {
            i++;
        }
    }

    i = get_bus_status();

    if (i >= 0) {
        *(p - i) = 10;
    }
}
```

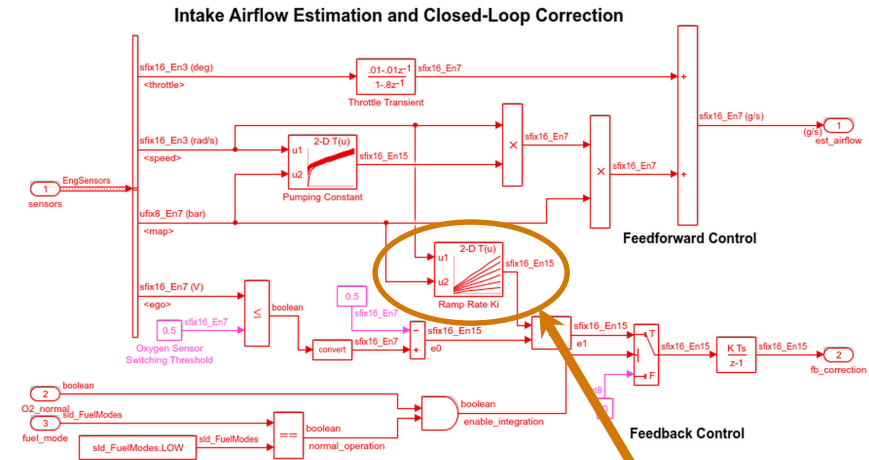
variable 'i' (int32): [0 .. 99]
assignment of 'i' (int32): [1 .. 100]

Bug Finder

- Bug Finder uses abstract interpretation techniques to find bugs
 - Does not claim to be exhaustive
Cannot prove absence of bugs
 - May find a broader set of bugs
Some are hard to deal with in Code Prover
- Few False Positives

Applications of Abstract Interpretation

- Analysis of C Code
 - Code prover
 - Bug finder
- Analysis of code generated from models
 - Opportunities for using model semantics to improve precision of code-based analysis



Code Generation Report

Find: Match Case

Highlight code for block: 'fuel_rate_control' 1 of 16 Remove Highlights

Contents

- Summary
- Subsystem Report
- Code Interface Report
- Traceability Report
- Static Code Metrics Report
- Code Replacements Report
- Generated Code
 - [-] Main file
 - ert_main.c
 - [-] Model files
 - fuel_rate_control.c (8)
 - fuel_rate_control.h (4)
 - [-] Data files
 - fuel_rate_control_data.c (4)
 - [+] Utility files (1)

```

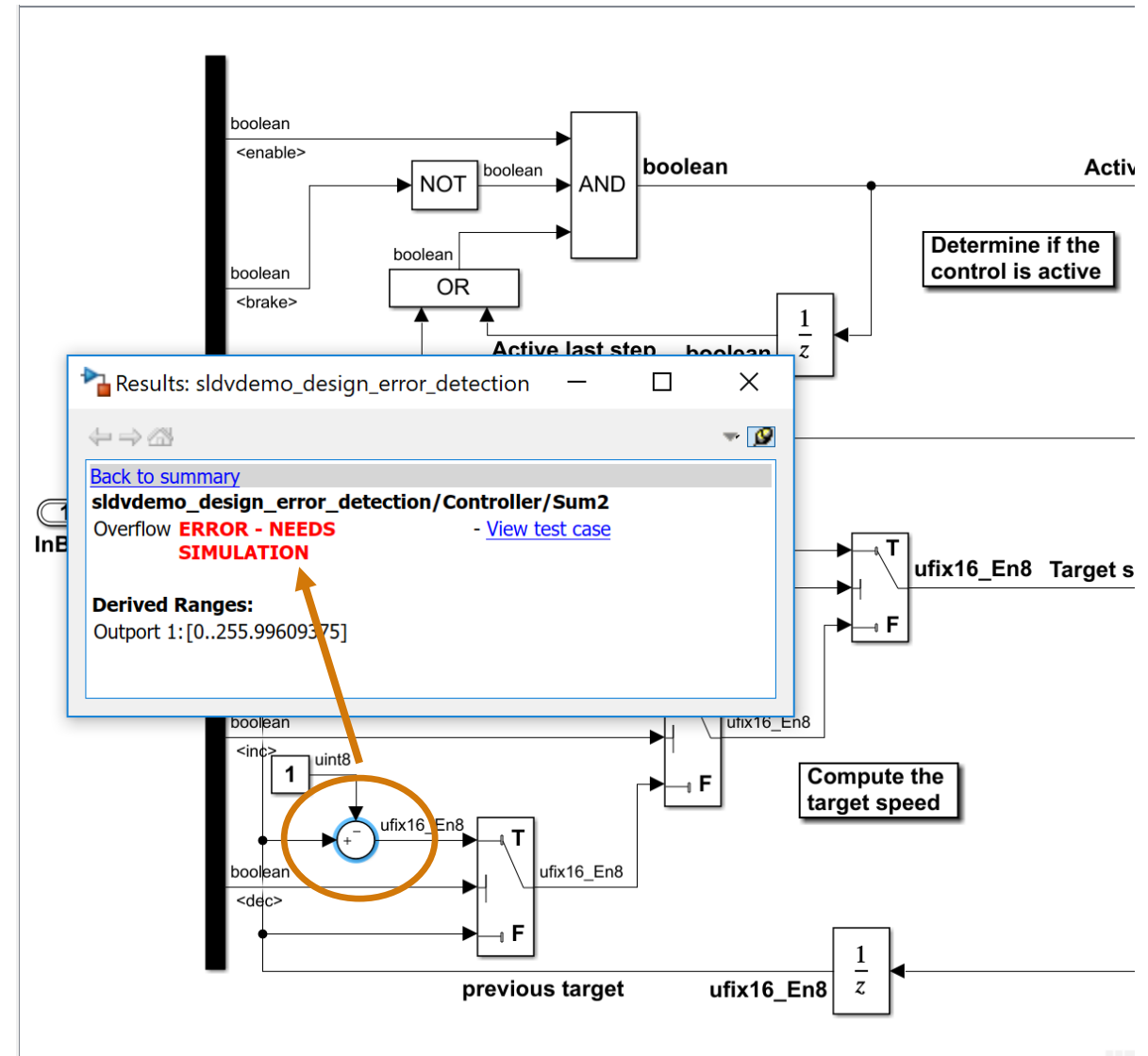
886  /*
887  /* DiscreteFilter: 'S13/Discrete_Filter' incorporates:
888  * DiscreteIntegrator: 'S2/Discrete_Integrator'
889  */
890  DiscreteFilter_tmp = (int16_T)((rtDWork.DiscreteIntegrator_DSTATE << 14) -
891  -12137 * rtDWork.DiscreteFilter_states_i) >> 14);
892
893  /* Output: '<Root>/fuel_rate' incorporates:
894  * DiscreteFilter: 'S13/Discrete_Filter'
895  * DiscreteFilter: 'S2/Throttle_Transient'
896  * Lookup_n-D: 'S2/Pumping_Constant'
897  * Product: 'S10/Product'
898  * Product: 'S2/Product'
899  * Product: 'S2/Product2'
900  * Sum: 'S13/Sum3'
901  * Sum: 'S2/Sum'
902  */
903  rtv.fuel_rate = (int16_T)((int16_T)((int16_T)((int16_T)((int16_T)
904  (rtDWork.es_o.speed * look2_is16u8l3n32ts1_K77emhOK(rtDWork.es_o.speed,
905  rtDWork.es_o.map, rtConstP.pooled2, rtConstP.pooled8,
906  rtConstP.PumpingConstant_tableData, rtConstP.pooled1, 180)) >> 11) *
907  (rtDWork.es_o.map) >> 7) + (int16_T)((20972 * rtb_SpeedEstimation - 20972 *
908  rtDWork.Throttle_Transient_states) >> 17)) * rtb_ThrottleTransientSwitch) >> 7) <<
909  8) + (int16_T)((17960 * DiscreteFilter_tmp + -17429 *
910  rtDWork.DiscreteFilter_states_i) >> 11)) >> 8);
911
912  /* Update for DiscreteFilter: 'S13/Discrete_Filter' */
913  rtDWork.DiscreteFilter_states_i = DiscreteFilter_tmp;

```

-
- The diagram illustrates the control logic for a speed control system. It features several inputs and outputs, including a bus system (InBus) and various control signals. The logic is implemented using a combination of gates (AND, OR, NOT), a delay block (1/z), and a summing junction (+/-). The system determines if control is active based on the 'Actual_speed' (ufix16_En8) and the 'previous target' speed. If active, the target speed is updated based on the difference between the target and actual speed, filtered by a 1/z block. The logic involves AND/OR gates, NOT gates, and a 1/z block for the target speed.

Applications of Abstract Interpretation

- Analysis of C Code
 - Code prover
 - Bug finder
- Analysis of code generated from models
 - Opportunities for using model semantics to improve precision of code-based analysis
- Analysis of Models



Detailed design – Data-type design

From Specification to Code

- High level design focuses on maximum flexibility in design space exploration
 - Use doubles
 - Let Simulink choose appropriate data-types
- Code generation focuses on designs that are quite precisely specified and optimised
 - Cost of computation
 - Cost of memory
 - Precision requirements

Data type design

- Start with MATLAB algorithm (uses doubles)

```
% Compute Output
if (u_state > limit_upper)
    y = limit_upper;
    clip_status = -2;
elseif (u_state >= limit_upper)
    y = limit_upper;
    clip_status = -1;
elseif (u_state < limit_lower)
    y = limit_lower;
    clip_status = 2;
elseif (u_state <= limit_lower)
    y = limit_lower;
    clip_status = 1;
else
    y = u_state;
    clip_status = 0;
end
```

Data type design

- Start with MATLAB algorithm (uses doubles)
- Convert all computations to use fixed-point calculations

```
% Compute Output
if (u_state > limit_upper)
    %F2F: No information found for converting the following block of code
    %F2F: Start block
    y = fi(limit_upper, 0, 16, 7, fm);
    clip_status = fi(-2, 0, 1, 0, fm);
    %F2F: End block
elseif (u_state >= limit_upper)
    %F2F: No information found for converting the following block of code
    %F2F: Start block
    y = fi(limit_upper, 0, 16, 7, fm);
    clip_status = fi(-1, 0, 1, 0, fm);
    %F2F: End block
elseif (u_state < limit_lower)
    %F2F: No information found for converting the following block of code
    %F2F: Start block
    y = fi(limit_lower, 0, 16, 7, fm);
    clip_status = fi(2, 0, 1, 0, fm);
    %F2F: End block
elseif (u_state <= limit_lower)
    %F2F: No information found for converting the following block of code
    %F2F: Start block
    y = fi(limit_lower, 0, 16, 7, fm);
    clip_status = fi(1, 0, 1, 0, fm);
    %F2F: End block
else
    y = fi(u_state, 0, 16, 7, fm);
    clip_status = fi(0, 0, 1, 0, fm);
end
```

Data type design

- Start with MATLAB algorithm (uses doubles)

```
% Compute Output
if (u_state > limit_upper)
    y = limit_upper;
    clip_status = -2;
elseif (u_state >= limit_upper)
    y = limit_upper;
    clip_status = -1;
elseif (u_state < limit_lower)
    y = limit_lower;
    clip_status = 2;
elseif (u_state <= limit_lower)
    y = limit_lower;
    clip_status = 1;
else
    y = u_state;
    clip_status = 0;
end
```

Data type design

- Start with MATLAB algorithm (uses doubles)
- Execute test-bench with sample inputs

```
% Compute Output
if (u_state > limit_upper)
    y = limit_upper;
    clip_status = -2;
elseif (u_state >= limit_upper)
    y = limit_upper;
    clip_status = -1;
elseif (u_state < limit_lower)
    y = limit_lower;
    clip_status = 2;
elseif (u_state <= limit_lower)
    y = limit_lower;
    clip_status = 1;
else
    y = u_state;
    clip_status = 0;
end
```

Data type design

- Start with MATLAB algorithm (uses doubles)
- Execute test-bench with sample inputs
- Instrument MATLAB to monitor min/max for each variable

```
% Compute Output
if (u_state > limit_upper)
    y = limit_upper;
    clip_status = -2;
elseif (u_state >= limit_upper)
    y = limit_upper;
    clip_status = -1;
elseif (u_state < limit_lower)
    y = limit_lower;
    clip_status = 2;
elseif (u_state <= limit_lower)
    y = limit_lower;
    clip_status = 1;
else
    y = u_state;
    clip_status = 0;
end
```

Data type design

- Start with MATLAB algorithm (uses doubles)
- Execute test-bench with sample inputs
- Instrument MATLAB to monitor min/max for each variable
- Propose fixpoint types based on observed ranges

```
% Compute Output
if (u_state > limit_upper)
    y = limit_upper;
    clip_status = -2;
elseif (u_state >= limit_upper)
    y = limit_upper;
    clip_status = -1;
elseif (u_state < limit_lower)
    y = limit_lower;
    clip_status = 2;
elseif (u_state <= limit_lower)
    y = limit_lower;
    clip_status = 1;
else
    y = u_state;
    clip_status = 0;
end
```

Data type design

- Start with MATLAB algorithm (uses doubles)
- Execute test-bench with sample inputs
- Instrument MATLAB to monitor min/max for each variable
- Propose fixpoint types based on observed ranges

```
% Compute Output
if (u_state > limit_upper)
    y = limit_upper;
    clip_status = -2;
elseif (u_state >= limit_upper)
    y = limit_upper;
    clip_status = -1;
elseif (u_state < limit_lower)
    y = limit_lower;
    clip_status = 2;
elseif (u_state <= limit_lower)
    y = limit_lower;
    clip_status = 1;
else
    y = u_state;
    clip_status = 0;
```

Variable	Type	Sim Min	Sim Max	Whole Num...	Proposed Type
Input					
u_in	double		-1	1 No	numerictype(1, 16, 14)
Output					
y	double		2	320.31 No	numerictype(0, 16, 7)
clip_status	double		0	0 Yes	numerictype(0, 1, 0)
Persistent					
u_state	double		2	320.31 No	numerictype(0, 16, 7)
Local					
init_val	double		1	1 Yes	numerictype(0, 1, 0)
gain_val	double		1	1 Yes	numerictype(0, 1, 0)
limit_upper	double		500	500 Yes	numerictype(0, 9, 0)
limit_lower	double		-500	-500 Yes	numerictype(1, 10, 0)
tprod	double		-1	1 No	numerictype(1, 16, 14)

Data type design

- Start with MATLAB algorithm (uses doubles)
- Execute test-bench with sample inputs
- Instrument MATLAB to monitor min/max for each variable
- Propose fixpoint types based on observed ranges
- Change algorithm to use the proposed types

```
% Compute Output
if (u_state > limit_upper)
    y = limit_upper;
    clip_status = -2;
elseif (u_state >= limit_upper)
    y = limit_upper;
    clip_status = -1;
elseif (u_state < limit_lower)
    y = limit_lower;
    clip_status = 2;
elseif (u_state <= limit_lower)
    y = limit_lower;
    clip_status = 1;
else
    y = u_state;
    clip_status = 0;
```

Variable	Type	Sim Min	Sim Max	Whole Num...	Proposed Type
Input					
u_in	double		-1	1 No	numerictype(1, 16, 14)
Output					
y	double		2	320.31 No	numerictype(0, 16, 7)
clip_status	double		0	0 Yes	numerictype(0, 1, 0)
Persistent					
u_state	double		2	320.31 No	numerictype(0, 16, 7)
Local					
init_val	double		1	1 Yes	numerictype(0, 1, 0)
gain_val	double		1	1 Yes	numerictype(0, 1, 0)
limit_upper	double		500	500 Yes	numerictype(0, 9, 0)
limit_lower	double		-500	-500 Yes	numerictype(1, 10, 0)
tprod	double		-1	1 No	numerictype(1, 16, 14)

Data type design

- Start with MATLAB algorithm (uses doubles)
- Execute test-bench with sample inputs
- Instrument MATLAB to monitor min/max for each variable
- Propose fixpoint types based on observed ranges
- Change algorithm to use the proposed types
- Verify behaviour on test-bench

```
% Compute Output
if (u_state > limit_upper)
    y = limit_upper;
    clip_status = -2;
elseif (u_state >= limit_upper)
    y = limit_upper;
    clip_status = -1;
elseif (u_state < limit_lower)
    y = limit_lower;
    clip_status = 2;
elseif (u_state <= limit_lower)
    y = limit_lower;
    clip_status = 1;
else
    y = u_state;
    clip_status = 0;
```

Variable	Type	Sim Min	Sim Max	Whole Num...	Proposed Type
Input					
u_in	double		-1	1 No	numerictype(1, 16, 14)
Output					
y	double		2	320.31 No	numerictype(0, 16, 7)
clip_status	double		0	0 Yes	numerictype(0, 1, 0)
Persistent					
u_state	double		2	320.31 No	numerictype(0, 16, 7)
Local					
init_val	double		1	1 Yes	numerictype(0, 1, 0)
gain_val	double		1	1 Yes	numerictype(0, 1, 0)
limit_upper	double		500	500 Yes	numerictype(0, 9, 0)
limit_lower	double		-500	-500 Yes	numerictype(1, 10, 0)
tprod	double		-1	1 No	numerictype(1, 16, 14)

Data type design

- Start with MATLAB algorithm (uses doubles)

```
% Compute Output
if (u_state > limit_upper)
    y = limit_upper;
    clip_status = -2;
elseif (u_state >= limit_upper)
    y = limit_upper;
    clip_status = -1;
elseif (u_state < limit_lower)
    y = limit_lower;
    clip_status = 2;
elseif (u_state <= limit_lower)
    y = limit_lower;
    clip_status = 1;
else
    y = u_state;
    clip_status = 0;
end
```

Data type design

- Start with MATLAB algorithm (uses doubles)
- Propose fixed-point types based on Abstract Interpretation

```
% Compute Output
if (u_state > limit_upper)
    y = limit_upper;
    clip_status = -2;
elseif (u_state >= limit_upper)
    y = limit_upper;
    clip_status = -1;
elseif (u_state < limit_lower)
    y = limit_lower;
    clip_status = 2;
elseif (u_state <= limit_lower)
    y = limit_lower;
    clip_status = 1;
else
    y = u_state;
    clip_status = 0;
end
```

Data type design

- Start with MATLAB algorithm (uses doubles)
- Propose fixed-point types based on Abstract Interpretation

```
% Compute Output
if (u_state > limit_upper)
    y = limit_upper;
    clip_status = -2;
elseif (u_state >= limit_upper)
    y = limit_upper;
    clip_status = -1;
elseif (u_state < limit_lower)
    y = limit_lower;
    clip_status = 2;
elseif (u_state <= limit_lower)
    y = limit_lower;
```

Variable	Type	Sim Min	Sim Max	Static Min	Static Max	Whole Num...
Input						
u_in	double	-1	1	-1	1	No
Output						
y	double	2	320.31	-500	500	No
clip_status	double	0	0	-2	2	Yes
Persistent						
u_state	double	2	320.31	-501	501	No
Local						
init_val	double	1	1	1	1	Yes
gain_val	double	1	1	1	1	Yes
limit_upper	double	500	500	500	500	Yes
limit_lower	double	-500	-500	-500	-500	Yes
tprod	double	-1	1	-1	1	No

Data type design

- Start with MATLAB algorithm (uses doubles)
- Propose fixed-point types based on Abstract Interpretation
- Change algorithm to use the proposed types

```
% Compute Output
if (u_state > limit_upper)
    y = limit_upper;
    clip_status = -2;
elseif (u_state >= limit_upper)
    y = limit_upper;
    clip_status = -1;
elseif (u_state < limit_lower)
    y = limit_lower;
    clip_status = 2;
elseif (u_state <= limit_lower)
    y = limit_lower;
```

Variable	Type	Sim Min	Sim Max	Static Min	Static Max	Whole Num...
Input						
u_in	double	-1	1	-1	1	No
Output						
y	double	2	320.31	-500	500	No
clip_status	double	0	0	-2	2	Yes
Persistent						
u_state	double	2	320.31	-501	501	No
Local						
init_val	double	1	1	1	1	Yes
gain_val	double	1	1	1	1	Yes
limit_upper	double	500	500	500	500	Yes
limit_lower	double	-500	-500	-500	-500	Yes
tprod	double	-1	1	-1	1	No

Data type design

- Start with MATLAB algorithm (uses doubles)
- Propose fixed-point types based on Abstract Interpretation
- Change algorithm to use the proposed types
- Verify behaviour on test-bench

```
% Compute Output
if (u_state > limit_upper)
    y = limit_upper;
    clip_status = -2;
elseif (u_state >= limit_upper)
    y = limit_upper;
    clip_status = -1;
elseif (u_state < limit_lower)
    y = limit_lower;
    clip_status = 2;
elseif (u_state <= limit_lower)
    y = limit_lower;
```

Variable	Type	Sim Min	Sim Max	Static Min	Static Max	Whole Num...
Input						
u_in	double	-1	1	-1	1	No
Output						
y	double	2	320.31	-500	500	No
clip_status	double	0	0	-2	2	Yes
Persistent						
u_state	double	2	320.31	-501	501	No
Local						
init_val	double	1	1	1	1	Yes
gain_val	double	1	1	1	1	Yes
limit_upper	double	500	500	500	500	Yes
limit_lower	double	-500	-500	-500	-500	Yes
tprod	double	-1	1	-1	1	No

Software Model Checking

Model Checking

- Start with a model
 - Various forms of state machines/transition systems
 - Represents a system
- Add a property
 - Various forms of logic – temporal, modal etc.
 - Represents a property
- Algorithms to check the satisfiability/validity of the property on the model

Software Model Checking

- Model is software
 - Complex control-flow
 - Functions/procedures
 - Data-structures, pointers
 - Libraries
 - ...

Software Model Checking

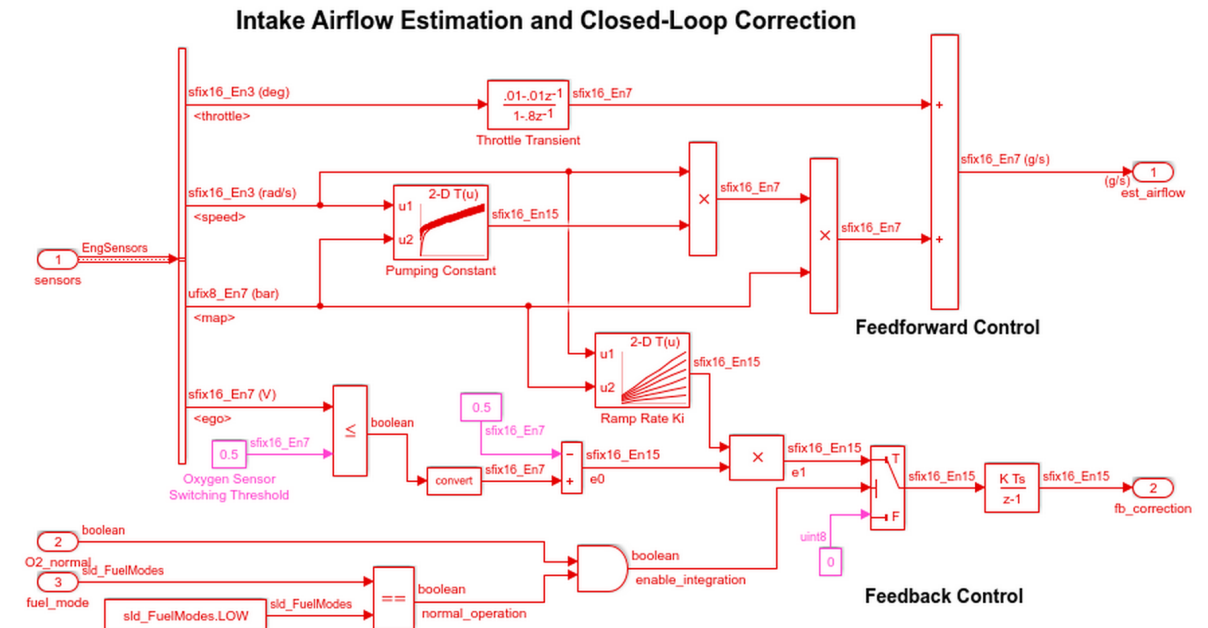
- Model is software
 - Complex control-flow
 - Functions/procedures
 - Data-structures, pointers
 - Libraries
 - ...
- Some properties can also be encoded in the software
 - ASSERT statements

Software Model Checking

- Model is software
 - Complex control-flow
 - Functions/procedures
 - Data-structures, pointers
 - Libraries
 - ...
- Some properties can also be encoded in the software
 - ASSERT statements
- Prove that the ASSERT statements cannot be hit or provide evidence of reachability

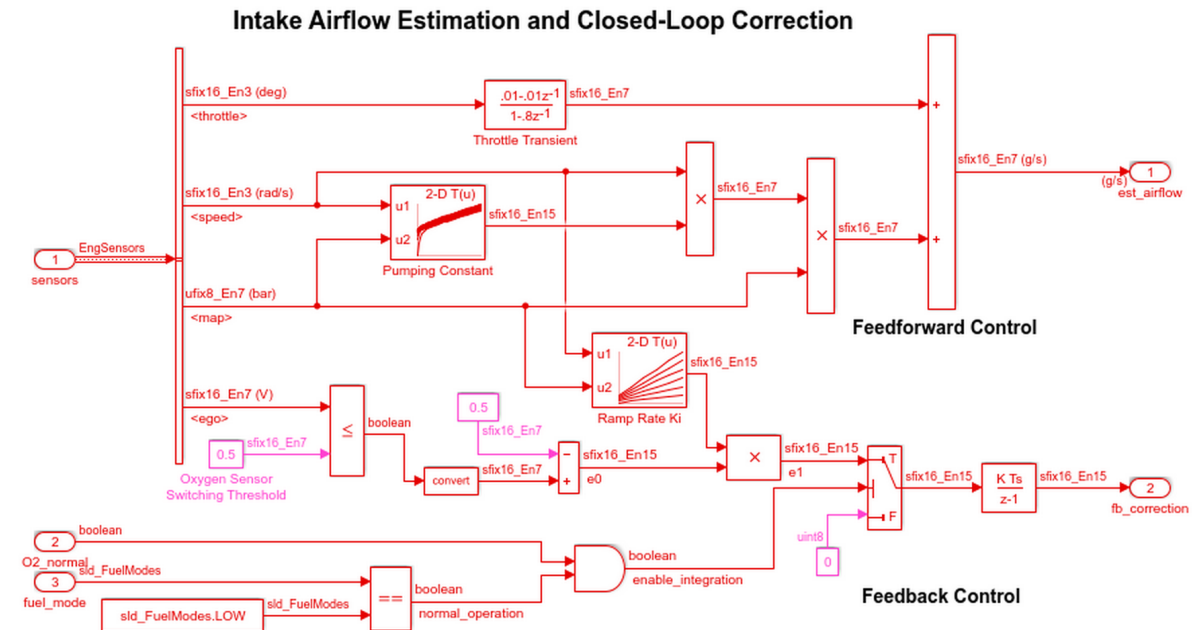
Model-Based Software Model Checking

- Model is the software (is the model)!



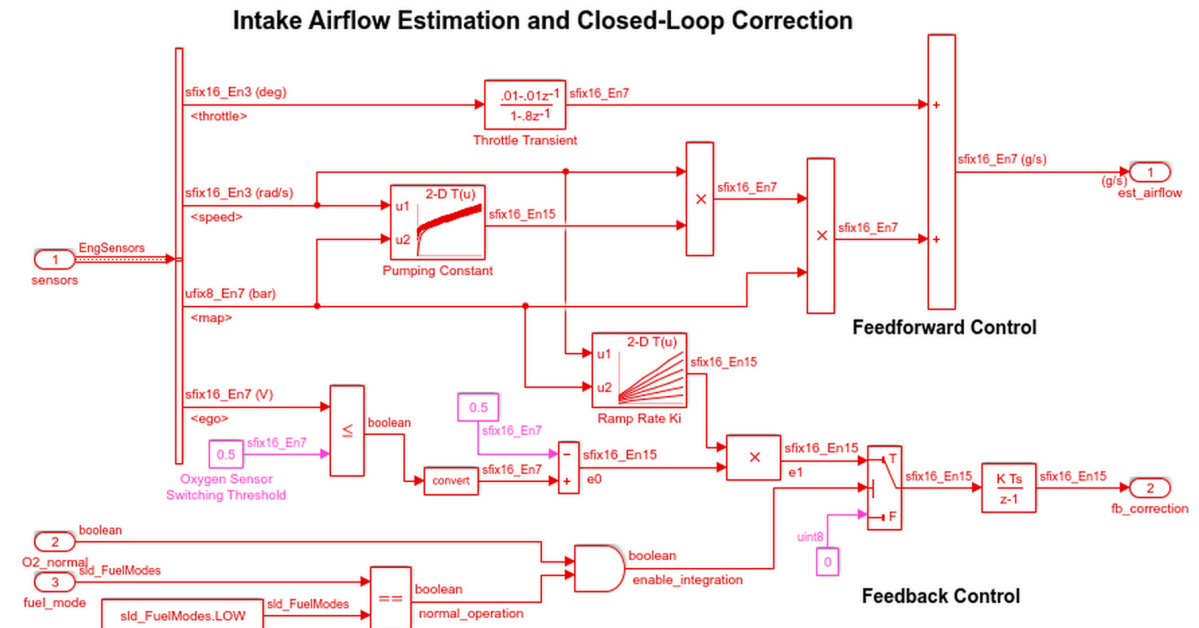
Model-Based Software Model Checking

- Model is the software (is the model)!
 - Discrete state Simulink diagrams



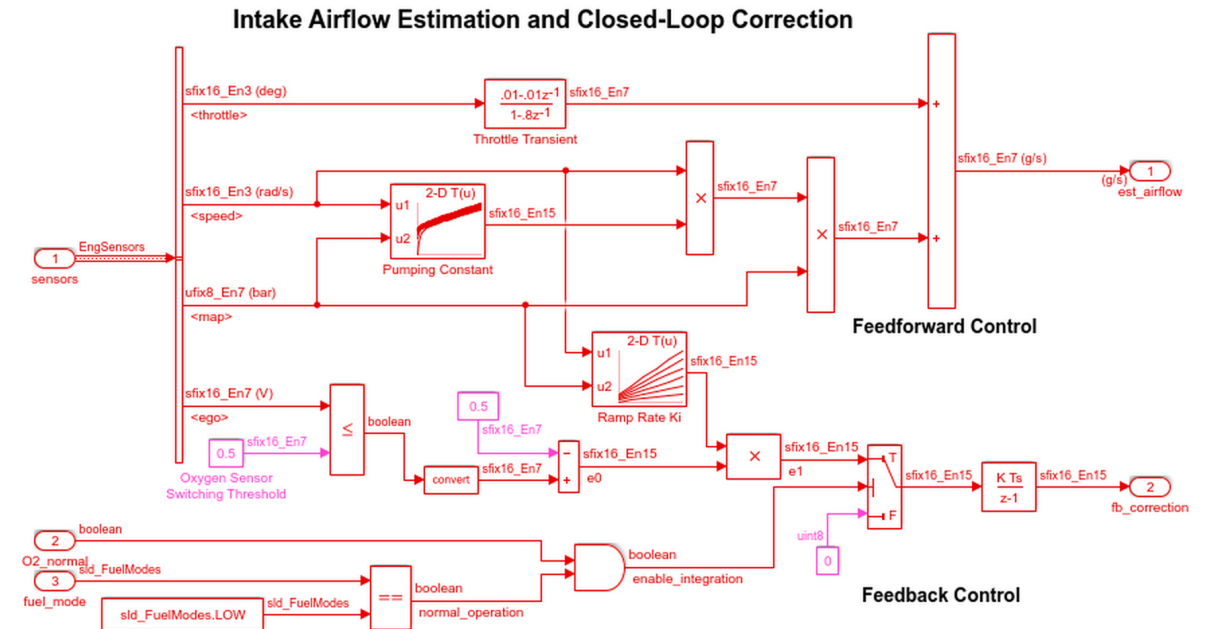
Model-Based Software Model Checking

- Model is the software (is the model)!
 - Discrete state Simulink diagrams
 - Stateflow



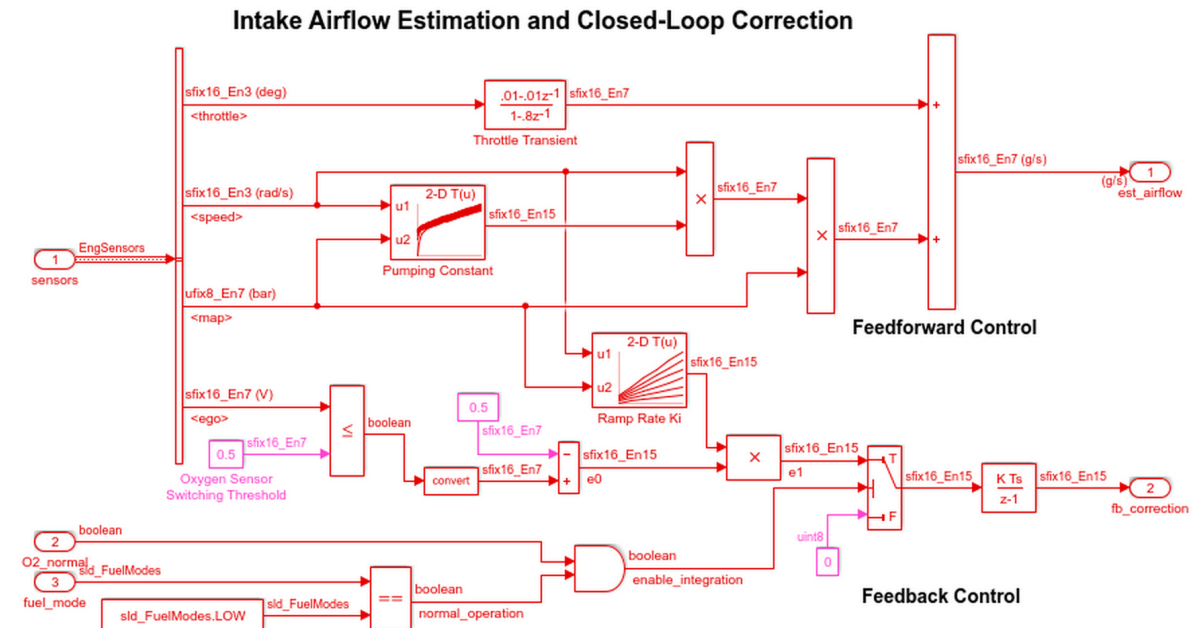
Model-Based Software Model Checking

- Model is the software (is the model)!
 - Discrete state Simulink diagrams
 - Stateflow
 - Tabular notations



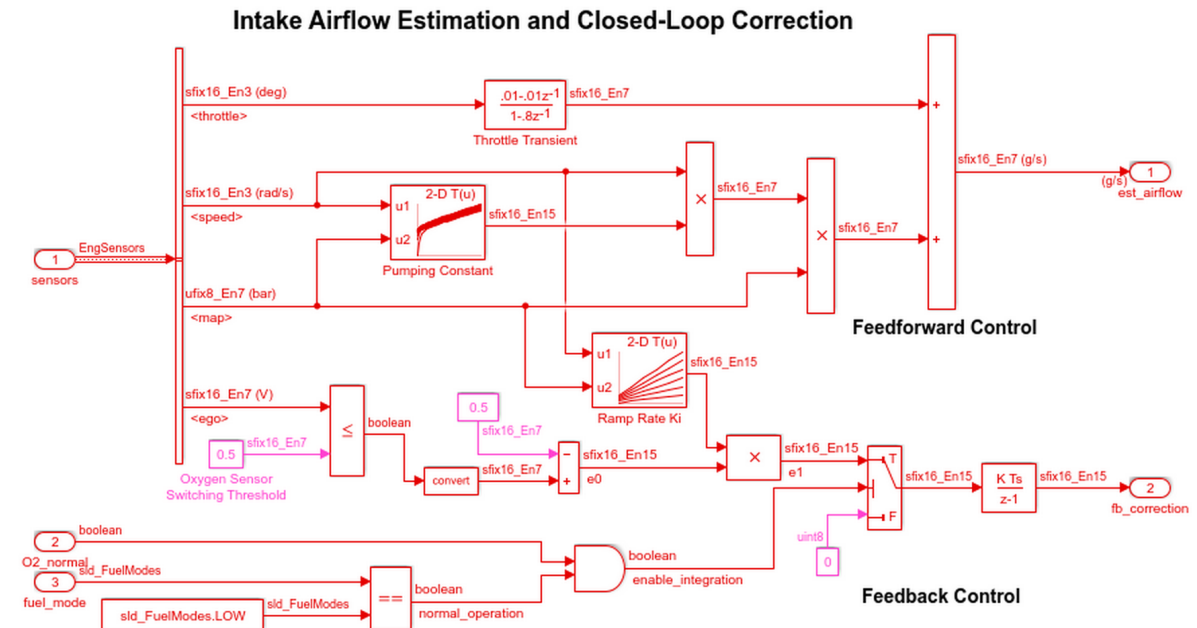
Model-Based Software Model Checking

- Model is the software (is the model)!
 - Discrete state Simulink diagrams
 - Stateflow
 - Tabular notations
 - Integrators, Lookup tables,



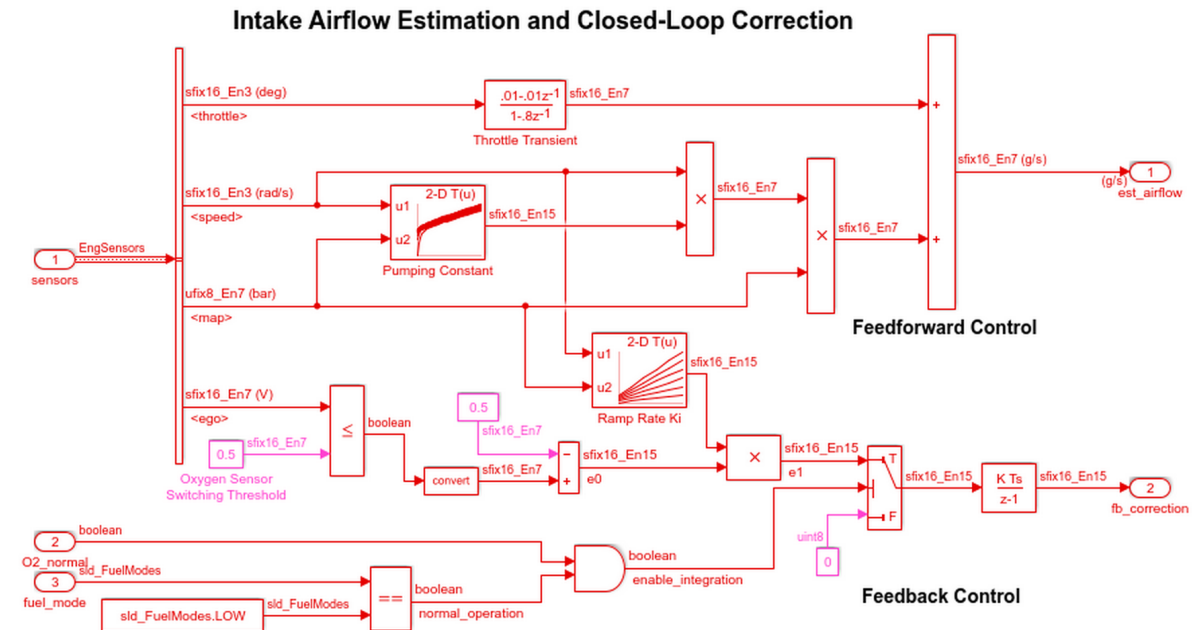
Model-Based Software Model Checking

- Model is the software (is the model)!
 - Discrete state Simulink diagrams
 - Stateflow
 - Tabular notations
 - Integrators, Lookup tables,
 - Signals, Triggers, Function calls, Subsystems, Reference Models



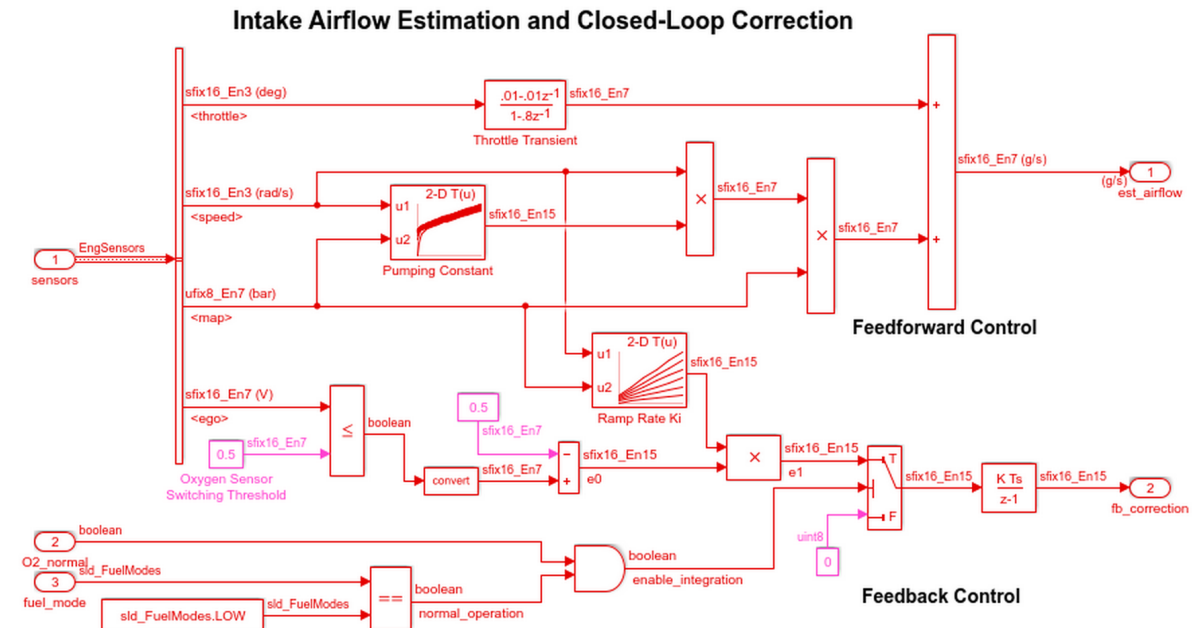
Model-Based Software Model Checking

- Model is the software (is the model)!
 - Discrete state Simulink diagrams
 - Stateflow
 - Tabular notations
 - Integrators, Lookup tables,
 - Signals, Triggers, Function calls, Subsystems, Reference Models
 - Arrays of signals, Buses (structures) of signals



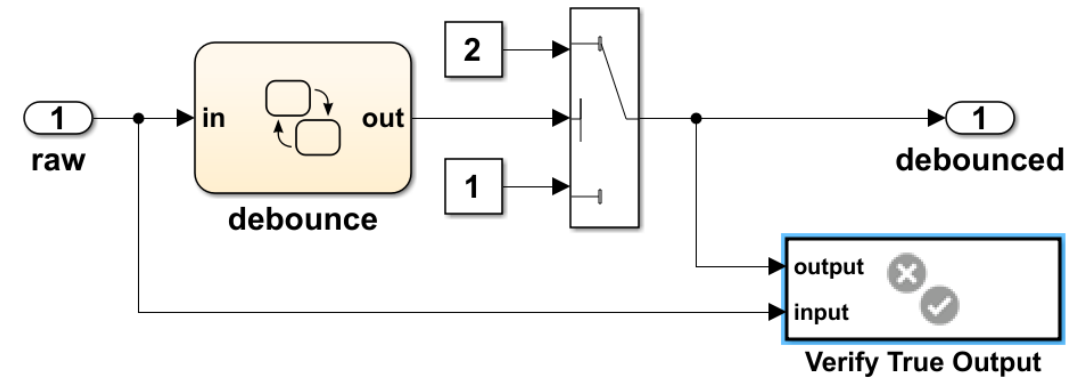
Model-Based Software Model Checking

- Model is the software (is the model)!
 - Discrete state Simulink diagrams
 - Stateflow
 - Tabular notations
 - Integrators, Lookup tables,
 - Signals, Triggers, Function calls, Subsystems, Reference Models
 - Arrays of signals, Buses (structures) of signals
 - MATLAB code



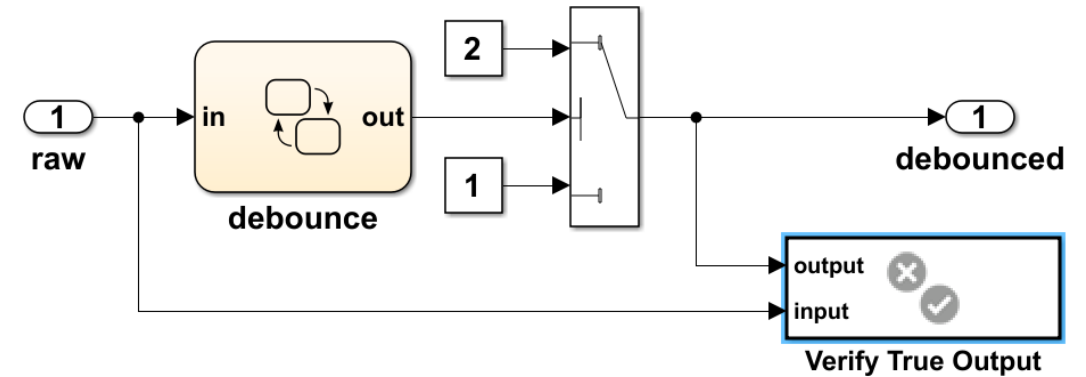
Model-Based Software Model Checking

- Model is the software (is the model!)



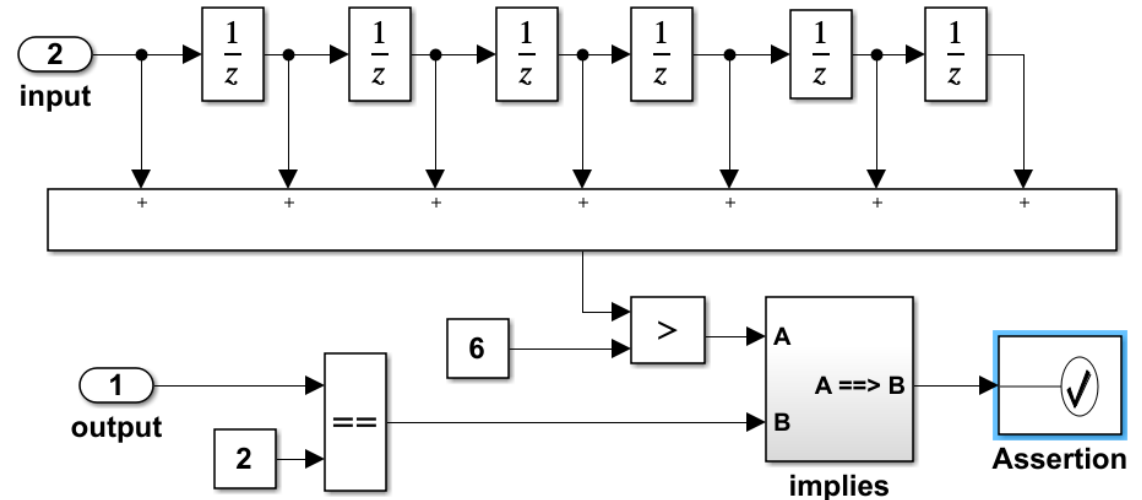
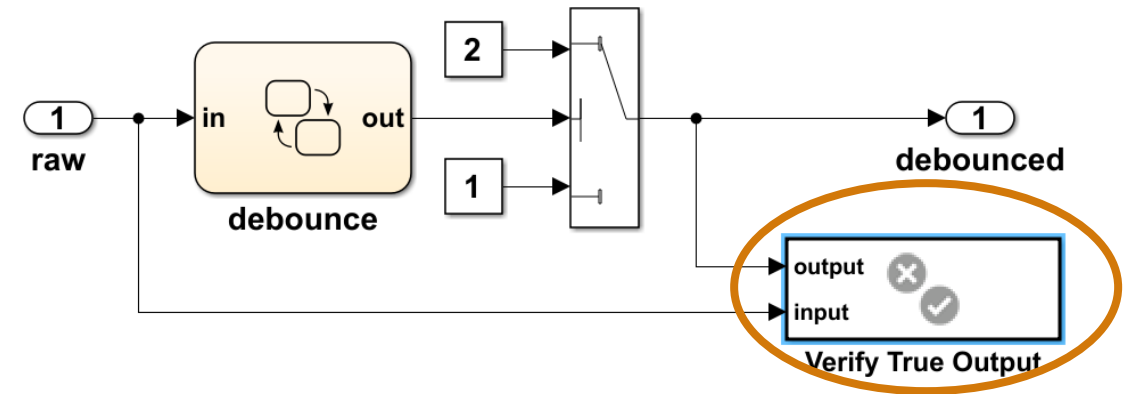
Model-Based Software Model Checking

- Model is the software (is the model!)
- Property is also modelled
 - Assertion blocks
 - Assumption blocks



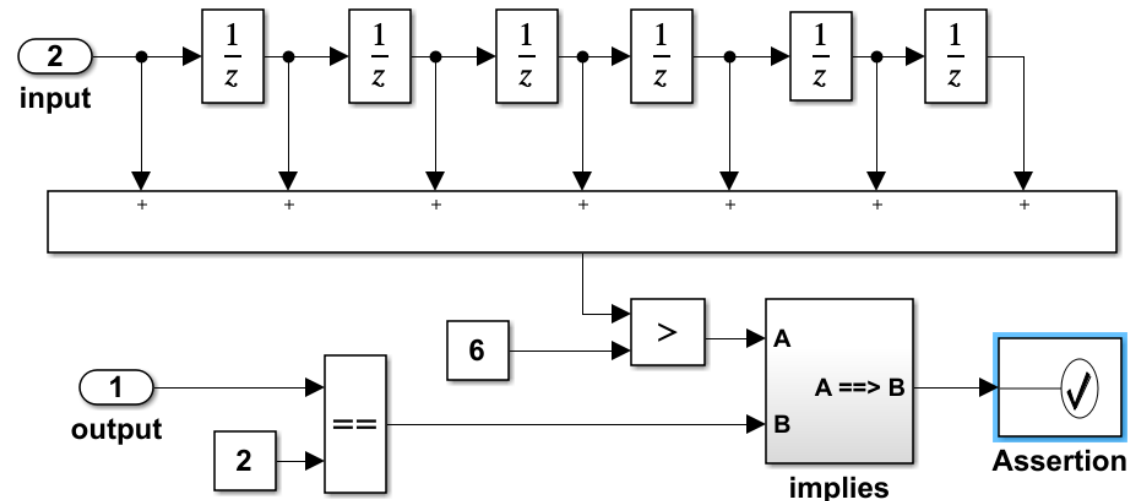
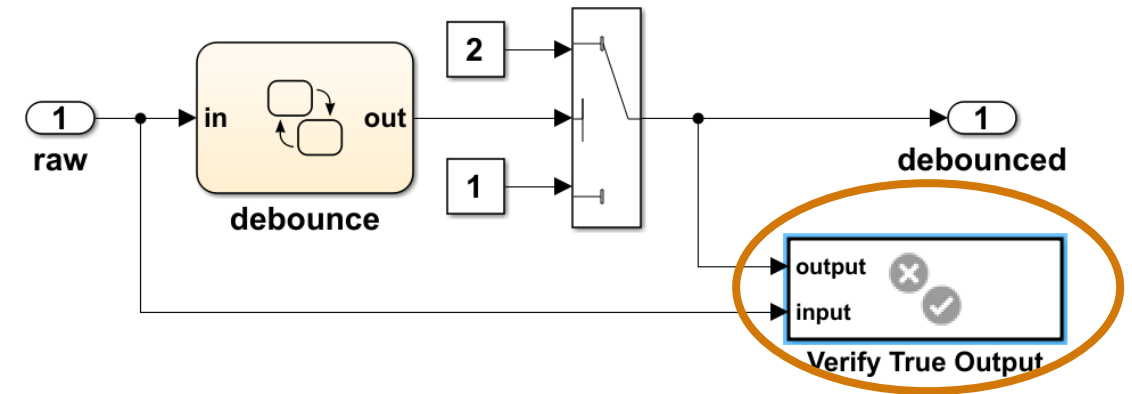
Model-Based Software Model Checking

- Model is the software (is the model!)
- Property is also modelled
 - Assertion blocks
 - Assumption blocks



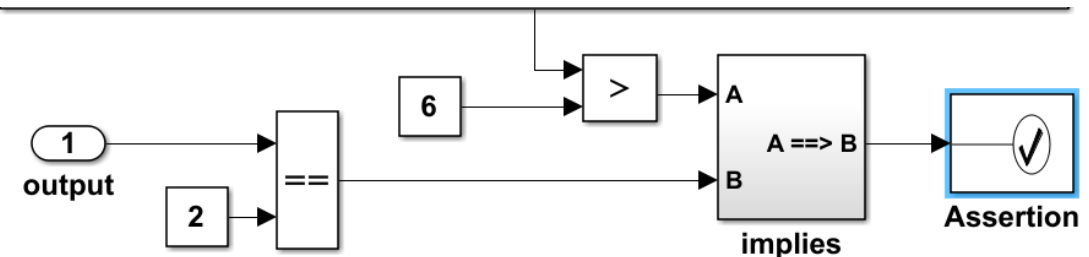
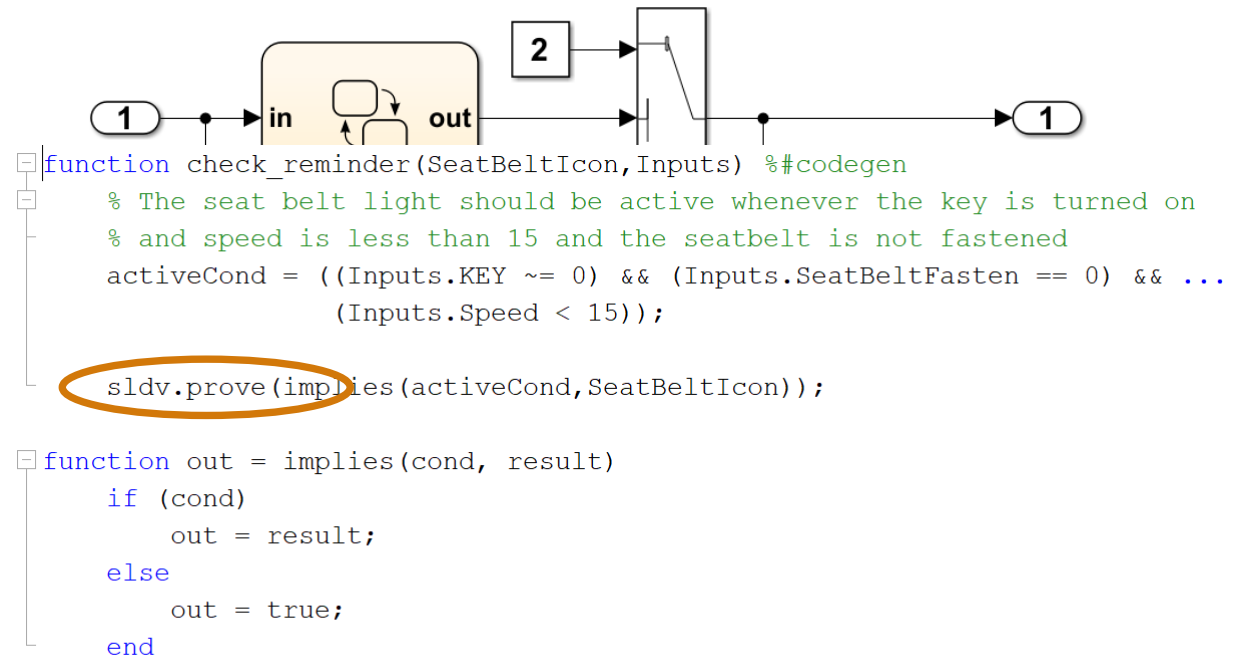
Model-Based Software Model Checking

- Model is the software (is the model!)
- Property is also modelled
 - Assertion blocks
 - Assumption blocks
 - Special statements in MATLAB code



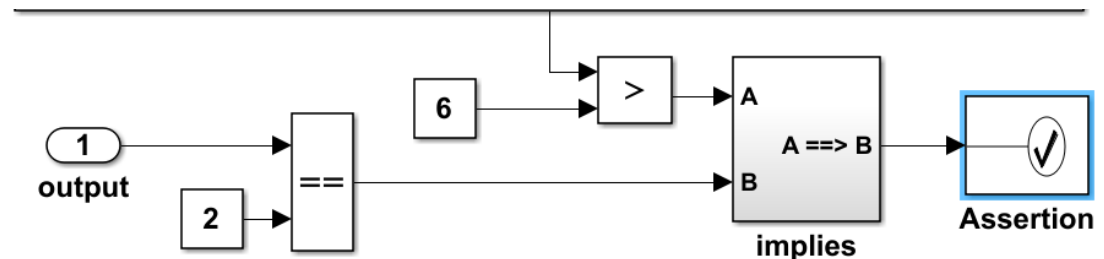
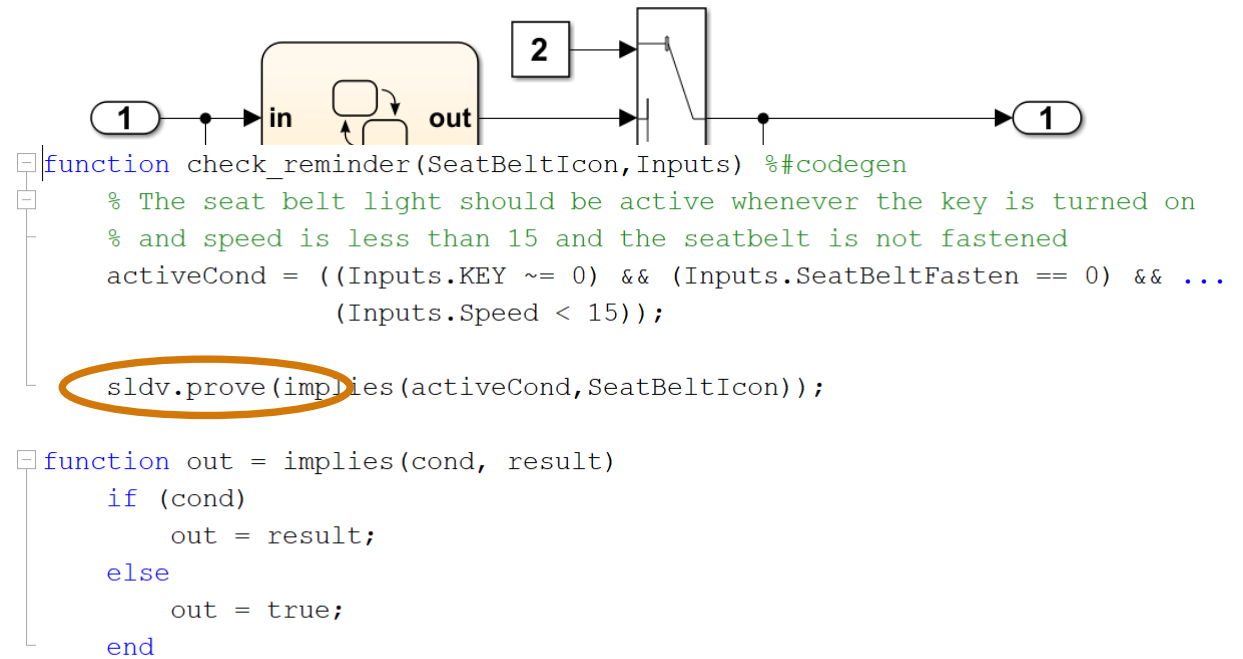
Model-Based Software Model Checking

- Model is the software (is the model!)
- Property is also modelled
 - Assertion blocks
 - Assumption blocks
 - Special statements in MATLAB code



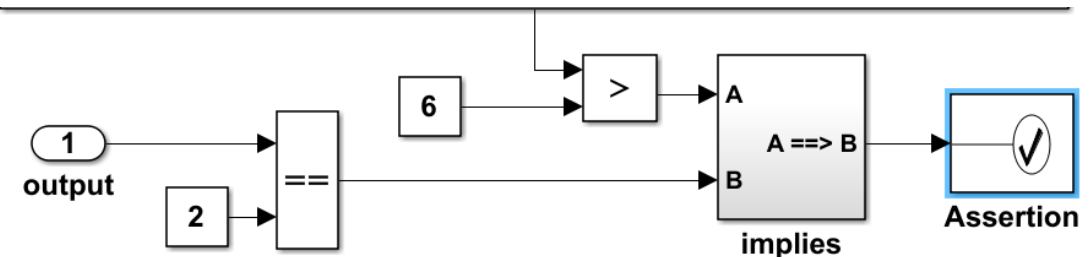
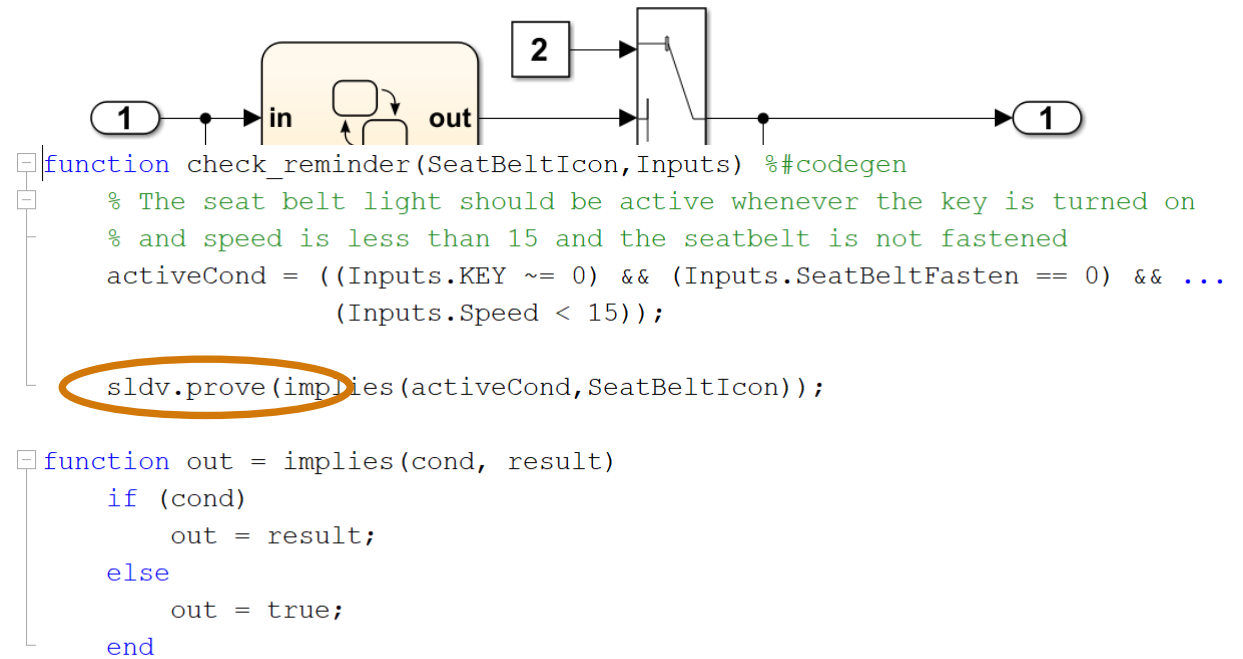
Model-Based Software Model Checking

- Model is the software (is the model!)
- Property is also modelled
 - Assertion blocks
 - Assumption blocks
 - Special statements in MATLAB code
- Prove properties valid

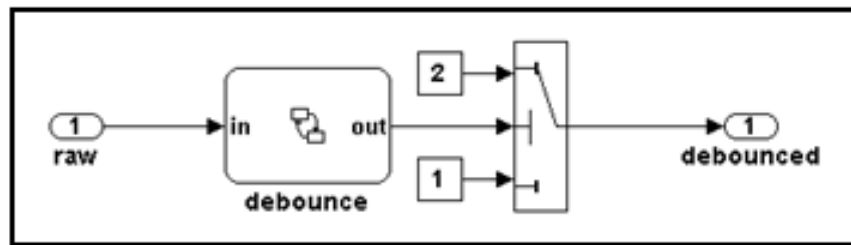


Model-Based Software Model Checking

- Model is the software (is the model!)
- Property is also modelled
 - Assertion blocks
 - Assumption blocks
 - Special statements in MATLAB code
- Prove properties valid
- Provide evidence of invalidity

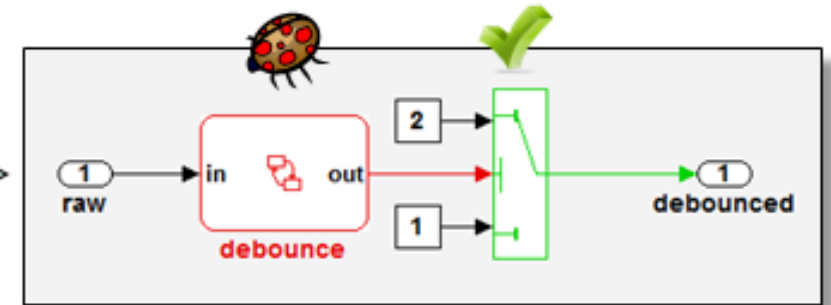


Run-time Error Check



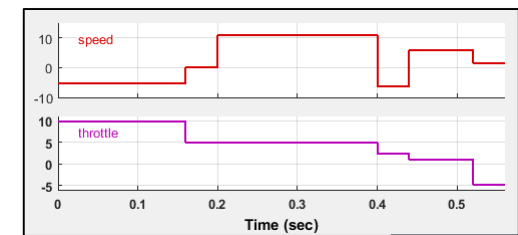
Design Model

Design error detection



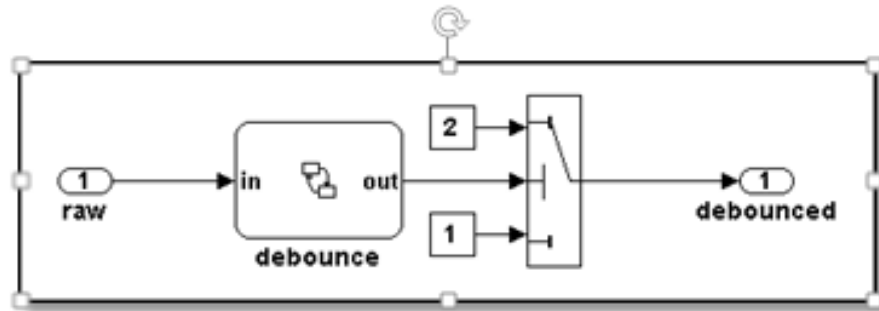
Highlighted Model

- Integer overflow
- Division by zero
- Array out-of-bounds
- Range violations
- Dead Logic



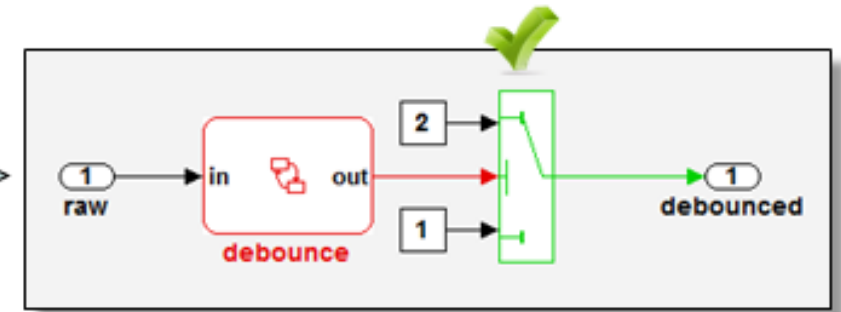
Counter-example

Test Generation Analysis



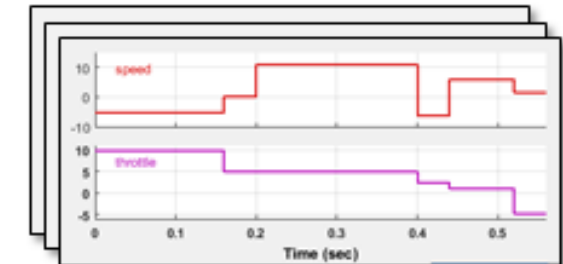
Design Model

Test Generation Analysis



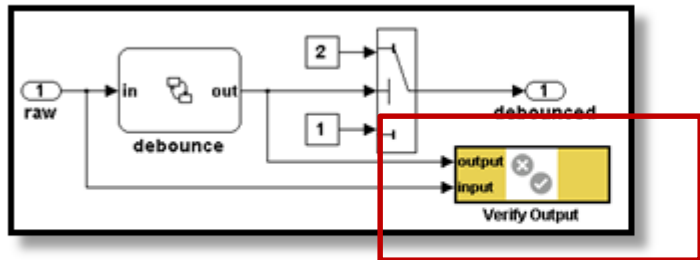
Highlighted Model

- Time varying inputs to satisfy coverage conditions
- Generate tests for missing coverage
- Special heuristics to satisfy timers/counters

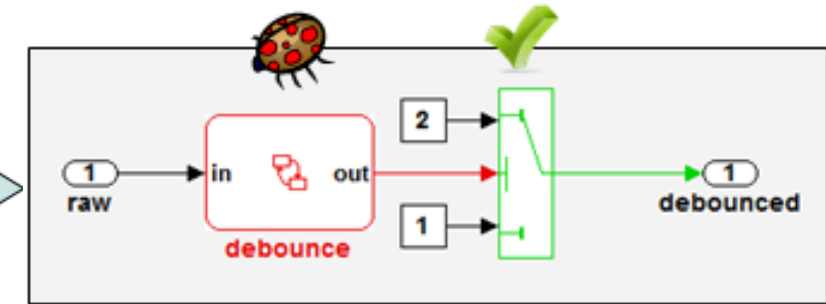


Test Cases

Property Proving

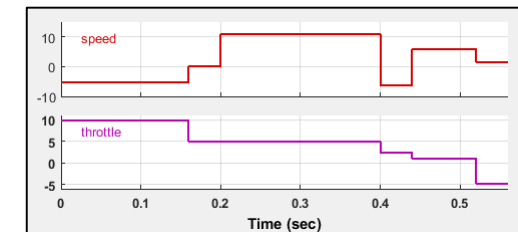


Design augmented with
modeled requirements



Highlighted Model

- Functional verification and witness generation for safety properties
- Check for bounded depth validation or for full guarantees



Counter-example

It's all about the Workflow!

Opportunities at MathWorks Bangalore

- Teams working on multiple areas of MATLAB and Simulink
- IN-SLVnV team
 - Simulink Design Verifier
 - Simulink Solvers and Execution Engine
 - HDL Code Generation and optimizations
 - Simulink Model Advisor
 - Simulink Test
 - Simulink Coder and optimizations
 - MATLAB Coder and optimizations
- Masters with some industry experience or PhD
 - Passion for developing software